

# SET OPERATIONS – UNION, INTERSECT, DIFFERENCE

Operation compatibility (i.e., compatible domain type)

(a) STUDENT

Fn	Ln
Susan	Yao
Ramesh	Shah
Johnny	Kohler
Barbara	Jones
Amy	Ford
Jimmy	Wang
Ernest	Gilbert

INSTRUCTOR

Fname	Lname
John	Smith
Ricardo	Browne
Susan	Yao
Francis	Johnson
Ramesh	Shah

(b)

Fn	Ln
Susan	Yao
Ramesh	Shah
Johnny	Kohler
Barbara	Jones
Amy	Ford
Jimmy	Wang
Ernest	Gilbert
John	Smith
Ricardo	Browne
Francis	Johnson

(c)

Fn	Ln
Susan	Yao
Ramesh	Shah

(d)

Fn	Ln
Johnny	Kohler
Barbara	Jones
Amy	Ford
Jimmy	Wang
Ernest	Gilbert

(e)

Fname	Lname
John	Smith
Ricardo	Browne
Francis	Johnson

Q. Show the tuples of the results in the following operations

- (b) STUDENT  $\cup$  INSTRUCTOR. (c) STUDENT  $\cap$  INSTRUCTOR. (d) STUDENT – INSTRUCTOR.  
(e) INSTRUCTOR – STUDENT.

**EMPLOYEE**

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

**WORKS\_ON**

<u>Essn</u>	<u>Pno</u>	Hours
123456789	1	32.5
123456789	2	7.5
666884444	3	40.0
453453453	1	20.0
453453453	2	20.0
333445555	2	10.0
333445555	3	10.0
333445555	10	10.0
333445555	20	10.0
999887777	30	30.0
999887777	10	10.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
987654321	20	15.0
888665555	20	NULL

**DEPARTMENT**

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

**DEPT\_LOCATIONS**

<u>Dnumber</u>	<u>Dlocation</u>
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

**PROJECT**

Pname	<u>Pnumber</u>	Plocation	Dnum
ProductX	1	Bellaire	5
ProductY	2	Sugarland	5
ProductZ	3	Houston	5
Computerization	10	Stafford	4
Reorganization	20	Houston	1
Newbenefits	30	Stafford	4

**DEPENDENT**

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
333445555	Alice	F	1986-04-05	Daughter
333445555	Theodore	M	1983-10-25	Son
333445555	Joy	F	1958-05-03	Spouse
987654321	Abner	M	1942-02-28	Spouse
123456789	Michael	M	1988-01-04	Son
123456789	Alice	F	1988-12-30	Daughter
123456789	Elizabeth	F	1967-05-05	Spouse

Figure 5.6 (company DB, textbook)

# SELECT and PROJECT operators

OPERATION	PURPOSE	NOTATION
SELECT	Selects all tuples that satisfy the selection condition from a relation $R$ .	$\sigma_{\langle \text{selection condition} \rangle}(R)$
PROJECT	Produces a new relation with only some of the attributes of $R$ , and removes duplicate tuples.	$\pi_{\langle \text{attribute list} \rangle}(R)$

$\sigma_{\text{Dno}=4 \text{ AND } \text{Salary}>25000}(\text{EMPLOYEE})$

$\sigma_{(\text{Dno}=4 \text{ AND } \text{Salary}>25000) \text{ OR } (\text{Dno}=5 \text{ AND } \text{Salary}>30000)}(\text{EMPLOYEE})$

$\pi_{\text{Lname, Fname, Salary}}(\text{EMPLOYEE})$

$\pi_{\text{Fname, Lname, Salary}}(\sigma_{\text{Dno}=5}(\text{EMPLOYEE}))$

intermediate relation, and using the **assignment operation**, denoted by  $\leftarrow$  (left arrow), as follows:

$\text{DEP5\_EMPS} \leftarrow \sigma_{\text{Dno}=5}(\text{EMPLOYEE})$

$\text{RESULT} \leftarrow \pi_{\text{Fname, Lname, Salary}}(\text{DEP5\_EMPS})$

# Sequence of operations and RENAME

$TEMP \leftarrow \sigma_{Dno=5}(EMPLOYEE)$

$R(First\_name, Last\_name, Salary) \leftarrow \pi_{Fname, Lname, Salary}(TEMP)$

We can also define a formal **RENAME** operation—which can rename either the relation name or the attribute names, or both—as a unary operator. The general RENAME operation when applied to a relation  $R$  of degree  $n$  is denoted by any of the following three forms:

$\rho_{S(B1, B2, \dots, Bn)}(R)$  or  $\rho_S(R)$  or  $\rho_{(B1, B2, \dots, Bn)}(R)$

Q. Retrieve the Social Security numbers of all employees who either work in department 5 or directly supervise an employee who works in department 5,

*Use a sequence of operations (i.e., intermediate outputs)*

$DEP5\_EMPS \leftarrow \sigma_{Dno=5}(EMPLOYEE)$

$RESULT1 \leftarrow \pi_{Ssn}(DEP5\_EMPS)$

$RESULT2(Ssn) \leftarrow \pi_{Super\_ssn}(DEP5\_EMPS)$

$RESULT \leftarrow RESULT1 \cup RESULT2$

# JOIN operators (EQUI-join most common)

The JOIN operation, denoted by  $\bowtie$ , is used to combine related tuples from two relations into single “longer” tuples. This operation is very important for any relational database with more than a single relation because it allows us to process relationships among relations.

Q. retrieve the name of the manager of each department

$$\text{DEPT\_MGR} \leftarrow \text{DEPARTMENT} \bowtie_{\text{Mgr\_ssn}=\text{Ssn}} \text{EMPLOYEE}$$
$$\text{RESULT} \leftarrow \pi_{\text{Dname, Lname, Fname}}(\text{DEPT\_MGR})$$

Q. retrieve the dependents of each employee

$$\text{EMP\_NAMES} \bowtie_{\text{Ssn}=\text{Essn}} \text{DEPENDENT}$$



# Outer JOINS

Left Outer JOIN

$TEMP \leftarrow (EMPLOYEE \bowtie_{Ssn=Mgr\_ssn} DEPARTMENT)$   
 $RESULT \leftarrow \pi_{Fname, Minit, Lname, Dname}(TEMP)$

Q. Tuples of RESULT relation according to the example COMPANY DB?

## RESULT

Fname	Minit	Lname	Dname
John	B	Smith	NULL
Franklin	T	Wong	Research
Alicia	J	Zelaya	NULL
Jennifer	S	Wallace	Administration
Ramesh	K	Narayan	NULL
Joyce	A	English	NULL
Ahmad	V	Jabbar	NULL
James	E	Borg	Headquarters

# Exercises

Query 1. Retrieve the name and address of all employees who work for the 'Research' department.

```
RESEARCH_DEPT ←  $\sigma_{Dname='Research'}$ (DEPARTMENT)
RESEARCH_EMPS ← (RESEARCH_DEPT  $\bowtie_{Dnumber=Dno}$  EMPLOYEE)
RESULT ←  $\pi_{Fname, Lname, Address}$ (RESEARCH_EMPS)
```

As a single in-line expression, this query becomes:

```
 $\pi_{Fname, Lname, Address} (\sigma_{Dname='Research'}(DEPARTMENT \bowtie_{Dnumber=Dno}(EMPLOYEE)))$ 
```

Query 2. For every project located in 'Stafford', list the project number, the controlling department number, and the department manager's last name, address, and birth date.

```
STAFFORD_PROJS ←  $\sigma_{Plocation='Stafford'}$ (PROJECT)
CONTR_DEPTS ← (STAFFORD_PROJS  $\bowtie_{Dnum=Dnumber}$  DEPARTMENT)
PROJ_DEPT_MGRS ← (CONTR_DEPTS  $\bowtie_{Mgr\_ssn=Ssn}$  EMPLOYEE)
RESULT ←  $\pi_{Pnumber, Dnum, Lname, Address, Bdate}$ (PROJ_DEPT_MGRS)
```

# Exercises

Query 3. List the names of all employees with two or more dependents

Strictly speaking, this query cannot be done in the *basic (original) relational algebra*. We have to use the AGGREGATE FUNCTION operation with the COUNT aggregate function. We assume that dependents of the *same* employee have *distinct* Dependent\_name values.

$$\begin{aligned} T1(\text{Ssn}, \text{No\_of\_dependents}) &\leftarrow_{\text{Essn}} \mathfrak{S} \text{ COUNT Dependent\_name}(\text{DEPENDENT}) \\ T2 &\leftarrow \sigma_{\text{No\_of\_dependents} > 2}(T1) \\ \text{RESULT} &\leftarrow \pi_{\text{Lname}, \text{Fname}}(T2 * \text{EMPLOYEE}) \end{aligned}$$

Query 4. Retrieve the names of employees who have no dependents.

This is an example of the type of query that uses the MINUS (SET DIFFERENCE) operation.

$$\begin{aligned} \text{ALL\_EMPS} &\leftarrow \pi_{\text{Ssn}}(\text{EMPLOYEE}) \\ \text{EMPS\_WITH\_DEPS}(\text{Ssn}) &\leftarrow \pi_{\text{Essn}}(\text{DEPENDENT}) \\ \text{EMPS\_WITHOUT\_DEPS} &\leftarrow (\text{ALL\_EMPS} - \text{EMPS\_WITH\_DEPS}) \\ \text{RESULT} &\leftarrow \pi_{\text{Lname}, \text{Fname}}(\text{EMPS\_WITHOUT\_DEPS} * \text{EMPLOYEE}) \end{aligned}$$



# Exercises

Query 5. For each department, retrieve the department name and the average salary of all employees working in that department.

$AVG\_SAL (Dno, avg\_sal) \leftarrow \gamma_{Dno, avg(salary)} (EMPLOYEE)$

$RESULT \leftarrow \pi_{Dname, avg\_sal} (\sigma_{Dno=Dno} (DEPARTMENT * AVG\_SAL))$