

1a)

Not well-formed, violation of grammar rule [14]: the symbol "<" is not allowed  
Inside of CharData

1b)

Not well-formed, violation of grammar rule [5]: the symbols "<" and ">" may not appear  
inside of tag name.

1c)

Not well-formed, violation of grammar rule [39] (there are four b-Start-tags but  
only tree b-End-tags)

1d)

Not well-formed, violation of grammar rule [10]: the symbol "<" may not appear inside of  
an attribute value.

1e)

Well-formed.

1f)

Well-formed.

1g)

Well-formed.

1h)

Not well-formed. Grammar violation of rule [39] just as for c).

2)

```
n=root;
repeat {
  while(lastChild(n)!=NIL)
  { n=lastChild(n);
    If(nodeType(n)==TEXT_NODE) print(nodeValue(n));
  }
  while(previousSibling(n)=NIL)
  { n=parent(n);}
  n=nextSibling(n);
  if(nodeType(n)==TEXT_NODE) print(nodeValue(n));
}
```

3)

```
id=1
while (lab(id)!="")
{
  if (lab(id)=="a") count[id]=1 else count[id]=0;
  for each child in dag(id) do
  {
    count[id] = count[id] + count[child]
  }
  id = id + 1
}
```

4)

When computing the minimal DAG, we need to determine whether a given subtree has  
occurred already. If we keep a table of pointers to subtrees that have already occurred,  
then to check for a given subtree if it is in the table takes worst case time

(# of trees in table) \* (# nodes in the subtree)

Which in the worst case is quadratic in the size of the input tree!

With hashing, we only need  
 $(\# \text{trees in the hash bucket}) * (\# \text{nodes in the subtree}).$

For the example, take  $\text{hash}(\text{tree}) = 1$  if tree is a leaf and  
 $\text{hash}(\text{tree}) = 2$  if not a leaf and contains no "f"  
 $\text{hash}(\text{tree}) = 3$  in all other cases.

Then

$\text{hash}(c) = \text{bucket } 1$

$\text{hash}(b(c, c)) = \text{bucket } 2$

$\text{hash}(f) = \text{bucket } 1$

$\text{hash}(b(f, c)) = \text{bucket } 3$

Etc.

Without hashing: check up to 6 nodes each time.

With hash: check only up to 3 nodes each time!

5)

`Descendants(Node p){`

`for(i=1; i<size(p); i++) print( p + i )`

`}`

`Children(Node p){`

`c = p+1;`

`while( c < p+size(p) ) { print( c ); c = c+size(c) }`

`}`

`Parent(Node p){`

`for(i=1; i<p; i++) if p is in Children(i) then print(i)`

`}`

`Following-Siblings(Node p){`

`f = p + size(p);`

`while( f < Parent(p) + p ) { print( f ); f = f+size(f) }`

`}`

`Preceding(Node p){`

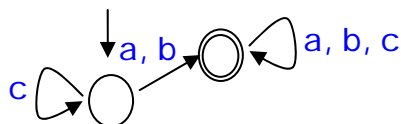
`for(i=1; i<p; i++) if(p not in Descendants (i)) then print(i)`

`}`

6a)

The string "a" is accepted; the string "c" is not accepted.

It is not deterministic (the initial state has two outgoing a-edges)

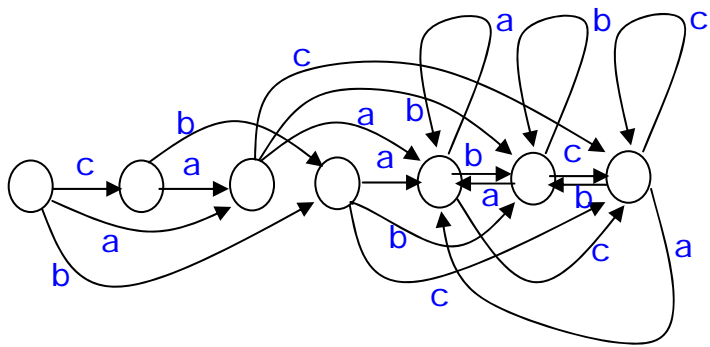


6b)

$c^*(a+b)(a+b+c)^*$

6c)

Not 1-unambiguous: Glushkov automaton is non-deterministic.



6d)

$(b^*(ab)^*)^*$