XML and Databases

Exam Preperation

Discuss Answers to last year's exam

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- a) <author></author><title></title>
- b) <author><title></author></title>
- c) <info temp='25C' >content</info>

```
g) <a><b><c><c/></c><c/>ab&e; </b></a>
```

```
a) <author></author><title></title>
                                             \rightarrow XML grammar
                                              (cannot be derived by grammar!)
b) <author><title></author></title>
c) <i nfo temp=' 25C' >content</i nfo>
d) <! DOCTYPE greeting [
    <! ELEMENT greeting (#PCDATA)>
    <! ENTI TY e1 "&e2; e3">
    <! ENTI TY e2 "&e3; ">
    <! ENTITY e3 "&e2; ">
1>
<greeting> &e1; </greeting>
e) <a at1="blah" at&lt; 2="foo" > 1 &lt; 5 </a>
f) <a b3="a" b2="b" b1="a" b2="5"/>
q) <a><b><c><c/><c/>>ab&e; </b></a>
```

```
a) <author></author><title> \rightarrow XML grammar
```

```
b) <author><title></author></title> \rightarrow WFC
```

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c) <info temp='25C' >content</info>
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g) <a><b><c><c/></c><c/>ab&e; </b></a>
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e) <a at1="blah" at&lt; 2="foo" > 1 &lt; 5 </a>
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f) <a b3="a" b2="b" b1="a" b2="5"/>
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```
g) <a><b><c><c/<>a>&e; </b></a>
```

 \rightarrow XML grammar

```
\rightarrow WFC
```

→ XML grammar (should be "25c"; -is actually OK..!)

a) <author></author> <title></title>	→ XML grammar		
b) <author><title></title></author>	\rightarrow WFC		
c) <info temp="25C">content</info>	\rightarrow XML grammar (should be "25c"; -is actually 0K 1)		
d) greeting [</td <td>→ WFC</td>	→ WFC		
<pre><!-- ELEMENT greeting (#PCDATA) --> Well-fou <!-- ENTITY e1 "&e2; e3"--> <!-- ENTITY e2 "&e3; "--> <!-- ENTITY e3 "&e2; "--> <!-- ENTITY e3 "&e2; "--> 1></pre>	rmedness constraint: No Recursion d entity <i>MUST NOT</i> contain a recursive ce to itself, either directly or indirectly.		
<preeting> &e1 </preeting>			
e) <a at<2="foo" at1="blah"> 1 < 5 			
f) <a b1="a" b2="5" b3="a">			
g) <a><c><c></c></c><c></c>ab&e 			

```
a) <author></author><title></title>
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b) <author><title></author></title></author></title>

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c) <info temp='25C' >content</info>
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```
d) <! DOCTYPE greeting [
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            <! ENTITY e2 "&e3; ">
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]>
```

```
e) <a at1="blah" at&lt; 2="foo" > 1 &lt; 5 </a> \rightarrow XML grammar
```

```
f) <a b3="a" b2="b" b1="a" b2="5"/>
```

```
g) <a><b><c><c/<>ab&e; </b></a>
```

[41] Attribute	::=	<u>Name Eq AttValue</u>	
[5] Name	::=	(<u>Letter</u> '_' ':')	(NameChar)*
[84] Letter	::=	[a-zA-Z]	

 \rightarrow XML grammar

 \rightarrow WFC

- \rightarrow XML grammar (should be "25c";
 - -is actually OK..!)

 \rightarrow WFC

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                                              \rightarrow XML grammar
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f) <a b3="a" b2="b" b1="a" b2="5"/>
                                                    \rightarrow WFC
g) <a><b><c><c/<c><c/><c/>ab&e; </b></a>
                                       An attribute name MUST NOT appear more than
                                        once in the same start-tag or empty-element tag.
```

```
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                                                  \rightarrow XML grammar
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                                                 \rightarrow WFC
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                                                       \rightarrow WFC
q) <a><b><c><c/<>ab&e; </b></a>
                                                       \rightarrow WFC
```

Well-formedness constraint: Entity Declared ... the <u>Name</u> given in the entity reference *MUST* <u>match</u> that in an <u>entity declaration</u> that...

- (2) Show sequences of Unicode characters for which
- a) UTF-8 needs more space than UTF-16

b) UTF-16 needs more space than UTF-8 together with the corresponding UTF codes and their lengths.

c) Explain how to binary sort a sequence of UTF-8 characters. Use pseudo code if appropriate.

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a) U+FFFF

UTF-8: 11101111 10111111 10111111 = 24bits UTF-16: 11111111 1111111 = 16bits (2) Show sequences of Unicode characters for which

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a) U+FFFF

UTF-8: 11101111 10111111 10111111 = 24bits UTF-16: 11111111 1111111 = 16bits

b) U+00

UTF-8: 0000000 = 8bits UTF-16: 0000000 0000000 = 16bits (2) Show sequences of Unicode characters for which

a) UTF-8 needs more space than UTF-16

b) UTF-16 needs more space than UTF-8 together with the corresponding UTF codes and their lengths.

c) Explain how to binary sort a sequence of UTF-8 characters. Use pseudo code if appropriate.

a) U+FFFF

UTF-8: 11101111 10111111 10111111 = 24bits UTF-16: 11111111 1111111 = 16bits

b) U+00

```
UTF-8: 0000000 = 8bits
UTF-16: 0000000 0000000 = 16bits
```

c)

To binary compare two characters, simply start from the highest bit! In this way, for characters with different lengths in UTF-8, after ≤ 4 bits we will be done!

- In case UTF-8 lengths are same \rightarrow normal binary compare..

a) returns all ancestors of the node

b) returns the previous sibling of the node.

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local name(e1) = "el em" children(e1) = [e2, c1] local name(e2) = "el em" children(e2) = [] parent(e2) = e1 code(c1) = U+00 parent(c1) = e1 attributes(e1) = [] attributes(e2) = []

a) returns all ancestors of the node

b) returns the previous sibling of the node.

```
localname(e1)
                   = "elem"
    children(e1) = [e2, c1]
    local name(e2) = "elem"
    children(e2)
                   = []
    parent(e2)
                   = e1
    code(c1)
                   = U + 00
    parent(c1)
                   = e1
    attributes(e1) = []
    attributes(e2) = []
a) getAncestors(Node n): NodeSet
    {
      NodeSet result=NULL;
      if(n.type!=DOC){
        for(; n=n->parent ; n. type! = "DOC") Add(n, result);
        Add(n, result);
      }
      return result;
    }
```

a) returns all ancestors of the node

b) returns the previous sibling of the node.

```
localname(e1)
                    = "elem"
                                                               getPrevSib(Node n): Node
                                                            b)
    children(e1)
                    = [e2, c1]
    local name(e2) = "elem"
                                                                  NodeList I=NULL;
    children(e2)
                    = []
                                                                  if(n.type!="DOC")
    parent(e2)
                    = e1
                                                                  {
    code(c1)
                    = U + 00
                                                                    Node parent=n->parentNode();
    parent(c1)
                    = e1
                                                                    I = n->chi I dren;
    attributes(e1) = []
                                                                    if(I==NULL) return NULL;
    attributes(e2) = []
                                                                    s=first(l);
                                                                    if(s==1) return NULL;
                                                                    while(s->next!=n)
                                                                      s=s->next();
a) getAncestors(Node n): NodeSet
                                                                  }
    {
                                                                  return s;
      NodeSet result=NULL;
                                                                }
      if(n.type!=DOC){
        for(; n=n->parent ; n. type! = "DOC") Add(n, result);
        Add(n, result);
      }
      return result;
    }
```

(4) Using DOM, give pseudo code that determines the average depth of the XML tree. The average depth of <a/>is 1.

```
int total =0;
int count=0;
call calcAverage(root, 1);
return total/count;
void calcAverage(Node n, int depth)
{
    NodeList children = n->childList();
    if(children->isEmpty())
    {
      total += depth;
      count++;
      return;
    }
    else for each Node c in children calcAverage(c, depth+1);
}
```

(5) Explain in detail, using an example, why hashing is useful for finding the minimal DAG of a tree.Why are updates more expensive on a DAG than on a tree?Give an example that clearly explains this.

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When computing the minimal DAG, we need to determine whether a given subtree has occurred already. If we keep a table of pointers to the subtrees that have occurred already, then to check for a given subtree whether or not it occurs in the table takes in the worst case (#of trees in the table) x (# nodes in subtree) which, in the worst case, is quadratic to the size of the input tree!

With hashing, we only need (# trees in the hash bucket) x (# nodes in the subtree).

Thus, if a bucket has only a constant number of trees, on average, then the complexity goes from quadratic to linear!

Example tree:



Seen: =NULL	
Seen:=seen + "a-tree"	
Contains(Seen, "b-tree")?	🗲 needs 1 comparison
Seen:=Seen + "b-tree"	
Contains(Seen, "c-tree")?	🗲 needs 2 comparisons
Seen:=Seen + "c-tree"	
Contains(Seen, "d-tree")?	← needs 3 comparisons
Seen:=Seen + "d-tree"	

Assume hash("a-tree")=1 hash("b-tree")=2 hash("c-tree")=3 hash("d-tree")=4

Then we need no (tree) comparisons whatsover!

(5) Explain in detail, using an example, why hashing is useful for finding the minimal DAG of a tree.Why are updates more expensive on a DAG than on a tree?Give an example that clearly explains this.



Inserting a single new child required adding two rows and changing one existing one.

(the shared 2nd child "b-subtree" of the a-node must be duplicated first, before the c-child can be added.

b) Give pseudo code that computes the POST order of a tree in an iterative way,
 i.e., without any recursive calls(!). You can use firstChild(n),
 nextSibling(n), and parent(n) for a node n.

- c) the ancestors of a node
- d) the last child of a node
- e) the maximal depth of the subtree at a node.

(6) Give the PRE/POST table for the tree <a><c/><c><d/><d></d></c><d/><d></c><d/></c><d/></c><d/></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c></d></c></d></d>

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- e) the maximal depth of the subtree at a node.

```
b) int i=1;
Node n=root;
repeat{
    while(firstChild(n)!=NULL) n=firstChild(n);
    post(i)=n;
    i++;
    while(nextSibling(n)==NIL){
        n=parent(n);
        if(n==NULL) break;
        post(i)=n;
        i++;
    }
    n=nextSibling(n);
}
```

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 - c) Given (pre, post) of a node, its ancestors are all nodes with pre-value < pre and post-value > post.

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 - c) Given (pre, post) of a node, its ancestors are all nodes with pre-value < pre and post-value > post.
 - d) If there is a node with pre-value > pre and with post-value=post-1, then that is the last child of (pre, post)

(6) Give the PRE/POST table for the tree <a><c/><c><d/><d></d></c><d/><d></c><d/></c><d/></c><d/></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c><d/>></c></d></c></d>

b) Give pseudo code that computes the POST order of a tree in an iterative way, i.e., without any recursive calls(!). You can use firstChild(n), nextSibling(n), and parent(n) for a node n.

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 - c) Given (pre, post) of a node, its ancestors are all nodes with pre-value < pre and post-value > post.
 - d) If there is a node with pre-value > pre and with post-value=post-1, then that is the last child of (pre, post)

```
e) int maxDepth(int pr){
    size(int p): int{
        int s=0;
        for(int pr2=p+1; post(pr2)<post(p); pr2++) s++
        return s;
    }
    int D=0; int u, L;
    L=pr+size(pr)-post(pr);
    for(int pr2=pr+1; post(pr2)<post(pr); pr2++){
        u=pr2+size(pr2)-post(pr2)-L;
        if(u>D) D=u;
    return D;
    }
```

(8) Show the Glushkov automaton for the regular expression E=(a | b)*a. Is this expression 1-unambiguous? Explain! Give a deterministic automaton for the same expression. Is E2=(b*a(a|b))*a equivalent to E? Is it 1-unambiguous? Show a 1-unambiguous expression that is equivalent to a(a | b)*. (8) Show the Glushkov automaton for the regular expression E=(a | b)*a. Is this expression 1-unambiguous? Explain! Give a deterministic automaton for the same expression. Is E2=(b*a(a|b))*a equivalent to E? Is it 1-unambiguous? Show a 1-unambiguous expression that is equivalent to a(a | b)*.



Deterministic??

(8) Show the GLushkov automaton for the regular expression E=(a | b)*a. Is this expression 1-unambiguous? Explain! Give a deterministic automaton for the same expression. Is E2=(b*a(a|b))*a equivalent to E? Is it 1-unambiguous? Show a 1-unambiguous expression that is equivalent to a(a | b)*.



Deterministic??

 \rightarrow no!

Thus, E is not 1-unambiguous.

(8) Show the GLushkov automaton for the regular expression E=(a | b)*a. Is this expression 1-unambiguous? Explain! Give a deterministic automaton for the same expression. Is E2=(b*a(a|b))*a equivalent to E? Is it 1-unambiguous? Show a 1-unambiguous expression that is equivalent to a(a | b)*.



Deterministic automaton:



(8) Show the Glushkov automaton for the regular expression $E=(a \mid b)*a$. Is this expression 1-unambiguous? Explain! Give a deterministic automaton for the same expression. Is $E2=(b*a(a \mid b))*a$ equivalent to E? Is it 1-unambiguous? Show a 1-unambiguous expression that is equivalent to $a(a \mid b)*$.

NOT equivalent to E! The string "ba" is matched by E, but NOT by E2.

E2 is NOT 1-unambiguous:



(8) Show the GLushkov automaton for the regular expression $E=(a \mid b)*a$. Is this expression 1-unambiguous? Explain! Give a deterministic automaton for the same expression. Is E2=(b*a(a|b))*a equivalent to E? Is it 1-unambiguous? Show a 1-unambiguous expression that is equivalent to $a(a \mid b)*$.

The expression b*a(b*a)* is

- \rightarrow equivalent to E
- \rightarrow 1-unambi guous.

(9) For the tree given in 6, write XPath expressions that

- a) select all b nodes
- b) select all b nodes that have a c-child
- c) select all b nodes that have no c-children
- d) select the right most c-node
- e) select all nodes that have a c-parent



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- d) select the right most c-node
- e) select all nodes that have a c-parent



a) //b

b) //b[c]

c) //b[not(c)]

- a) select all b nodes
- b) select all b nodes that have a c-child
- c) select all b nodes that have no c-children
- d) select the right most c-node (node no 11)
- e) select all nodes that have a c-parent



a) //b

- b) //b[c]
- c) //b[not(c)]

d) //c[b]

or //c[b and d]

- a) select all b nodes
- b) select all b nodes that have a c-child
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- d) select the right most c-node (node no 11)
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- a) //b
- b) //b[c]
- c) //b[not(c)]
- d) //c[b]

or //c[b and d]

e) //*[parent::c]

or //c/*

a) //c//d

- b) //*[a or b]
- c) //b/ancestor::d/following::d
- d) //*[not(.//b | ./ancestor::c)]
- e) //c//d/precedi ng::*//d



a) //c//d

- b) //*[a or b]
- c) //b/ancestor::d/following::d
- d) //*[not(.//b | ./ancestor::c)]
- e) //c//d/preceding::*//d



a) 5, 6, 13

- a) //c//d
- b) //*[a or b]
- c) //b/ancestor::d/following::d
- d) //*[not(.//b | ./ancestor::c)]
- e) //c//d/preceding::*//d



- a) //c//d
- b) //*[a or b]
- c) //b/ancestor::d/following::d
- d) //*[not(.//b | ./ancestor::c)]
- e) //c//d/precedi ng::*//d



- a) 5, 6, 13
- b) 1, 6, 11
- c) 9, 13, 14

Show the sequences of node numbers that are selected by the following queries.

a) //c//d

- b) //*[a or b]
- c) //b/ancestor::d/following::d
- d) //*[not(. //b | . /ancestor::c)]
- e) //c//d/precedi ng::*//d



- a) 5, 6, 13
 b) 1, 6, 11
 c) 9, 13, 14
- d) 2, 3, 9, 14

Show the sequences of node numbers that are selected by the following queries.

- a) //c//d
- b) //*[a or b]
- c) //b/ancestor::d/following::d
- d) //*[not(.//b | ./ancestor::c)]
- e) //c//d/precedi ng::*//d



- b) 1, 6, 11 c) 9, 13, 14
- d) 2, 3, 9, 14
- e) 5, 6

```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
(10) This a tree corresponding to the XML in (6).
```

```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
(10) This a tree corresponding to the XML in (6).
```

```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
(10) This a tree corresponding to the XML in (6).
```

```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```

For query c) show in detail how the **Core-XPath evaluation algorithm** computes the answer to this query. Do the same for **query d**).



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```
c) //b/ancestor::d/following::d
d) //*[not(.//b | ./ancestor::c)]
```



q1 = /b//b q2 = //b

or

q1 = //b q2 = //*



or

q1 = //b q2 = //*





















(12) Construct a DTD such that 10a) is included in 10e), and another DTD such that 10e) is included in 10a). [Very easy!!]

> 10a) //c//d 10e) //c//d/preceding::*//d

(12) Construct a DTD such that 10a) is included in 10e), and another DTD such that 10e) is included in 10a). [Very easy!!]

> 10a) //c//d 10e) //c//d/precedi ng::*//d

For the first part: *Any DTD* so that c-nodes do NOT have d-descendants! ©

For the second part: Any DTD, so that all d-nodes are c-descendants.
- a) /*
- b) /a/b/*
- c) //a/*//b
- d) //a/following-sibling::b

- a) /*
- b) /a/b/*
- c) //a/*//b
- d) //a/following-sibling::b

SELECT DISTINCT r1.pre FROM doc_tbl r1 WHERE r1.pre=1 ORDERED BY r1.pr

- a) /*
- b) /a/b/*
- c) //a/*//b
- d) //a/following-sibling::b

```
SELECT DISTINCT r4. pre FROM doc_tbl r1, r2, r3, r4

WHERE r1. pre=0

AND r2. pre>r1. pre

AND r2. post<r1. post

AND (r2. pre-r2. post+r2. size)=(r1. pre-r1. post+r1. size)+1

AND r2. tag="a"

AND r3. pre>r2. pre

AND r3. post<r2. post

AND (r3. pre-r3. post+r3. size)=(r2. pre-r2. post+r2. size)+1

AND r3. tag="b"

AND r4. pre>r3. pre

AND r4. post<r3. post

AND (r4. pre-r4. post+r4. size)=(r3. pre-r3. post+r3. size)+1

ORDERED BY r4. pre
```

```
Recall: level(n) = pre(n)-post(n)+size(n)
```

a) /*

- b) /a/b/*
- c) //a/*//b
- d) //a/following-sibling::b

```
SELECT DI STI NCT r4. pre FROM doc_tbl r1, r2, r3, r4
WHERE r1. pre=0
AND r2. pre>r1. pre
AND r2. post<r1. post
AND r2. tag="a"
AND r3. pre>r2. pre
AND r3. post<r2. post
AND (r3. pre-r3. post+r3. si ze)=(r2. pre-r2. post+r2. si ze)+1
AND r4. pre>r3. pre
AND r4. post<r3. post
AND r4. tag="b"
ORDERED BY r4. pre</pre>
```

Recall: level(n) = pre(n)-post(n)+size(n)

- a) /*
- b) /a/b/*
- c) //a/*//b
- d) //a/following-sibling::b



Recall: level(n) = pre(n)-post(n)+size(n)