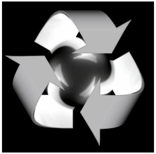
COMP4161
Advanced Topics in Software Verification



Thomas Sewell, Miki Tanaka, Rob Sison

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Content

- Foundations & Principles
  - Intro, Lambda calculus, natural deduction [1,2]
  - Higher Order Logic, Isar (part 1) [2,3[a]]
  - Term rewriting [3,4]
- Proof & Specification Techniques
  - Inductively defined sets, rule induction [4,5]
  - Datatype induction, primitive recursion [5,7]
  - General recursive functions, termination proofs [7]
  - Proof automation, Isar (part 2) [8[b]]
  - Hoare logic, proofs about programs, invariants [8,9]
  - C verification [9,10]
  - Practice, questions, exam prep [10[c]]

[a] a1 due; [b] a2 due; [c] a3 due

Datatypes

Example:

datatype 'a list = Nil | Cons 'a "'a list"

Properties:

- Constructors:

Nil	::	'a list
Cons	::	'a $\Rightarrow$ 'a list $\Rightarrow$ 'a list

- Distinctness: Nil $\neq$ Cons x xs
- Injectivity: (Cons x xs $=$ Cons y ys) $=$ (x $=$ y $\wedge$ xs $=$ ys)

More Examples

Enumeration:

datatype answer = Yes | No | Maybe

Polymorphic:

datatype 'a option = None | Some 'a
datatype ('a,'b,'c) triple = Triple 'a 'b 'c

Recursion:

datatype 'a list = Nil | Cons 'a "'a list"
datatype 'a tree = Tip | Node 'a "'a tree" "'a tree"

Mutual Recursion:

datatype even = EvenZero | EvenSucc odd
and odd = OddSucc even

Nested

Nested recursion:

datatype 'a tree = Tip | Node 'a "'a tree list"

datatype 'a tree = Tip | Node 'a "'a tree option" "'a tree option"

➜ Recursive call is under a type constructor.

The General Case

$$\textbf{datatype}\ (\alpha_1, \ldots, \alpha_n)\ \tau \quad = \quad C_1\ \tau_{1,1}\ \ldots\ \tau_{1,n_1} \;\; | \;\; \ldots \;\; | \;\; C_k\ \tau_{k,1}\ \ldots\ \tau_{k,n_k}$$

- ➜ Constructors: $C_i :: \tau_{i,1} \Rightarrow \ldots \Rightarrow \tau_{i,n_i} \Rightarrow (\alpha_1, \ldots, \alpha_n)\ \tau$
- ➜ Distinctness: $C_i\ \ldots \neq C_j\ \ldots$ if $i \neq j$
- ➜ Injectivity: $(C_i\ x_1 \ldots x_{n_i} = C_i\ y_1 \ldots y_{n_i}) = (x_1 = y_1 \wedge \ldots \wedge x_{n_i} = y_{n_i})$

Distinctness and Injectivity applied automatically

How is this Type Defined?

datatype 'a list = Nil | Cons 'a "'a list"

- ➜ internally reduced to a single constructor, using product and sum
- ➜ constructor defined as an inductive set (like typedef)
- ➜ recursion: least fixpoint

More detail: Tutorial on (Co-)datatypes Definitions at isabelle.in.tum.de

Datatype Limitations

Must be definable as a (non-empty) set.

- ➜ Infinitely branching ok.
- ➜ Mutually recursive ok.
- ➜ Strictly positive (right of function arrow) occurrence ok.

Not ok:

$$\begin{array}{lcl} \textbf{datatype}\ t & = & C\ (t \Rightarrow bool) \\ & | & D\ ((bool \Rightarrow t) \Rightarrow bool) \\ & | & E\ ((t \Rightarrow bool) \Rightarrow bool) \end{array}$$

Because: Cantor's theorem (α set is larger than α)

Datatype Limitations

Not ok (nested recursion):

datatype ('a, 'b) fun_copy = Fun "'a ⇒ 'b"

datatype 'a t = F "('a t, 'a) fun_copy"

- ➜ recursion in ('a1, ...,'an) t is only allowed on a subset of 'a1 ... 'an
- ➜ these arguments are called *live* arguments
- ➜ Mainly: in "'a ⇒ 'b", 'a is dead and 'b is live
- ➜ Thus: in ('a, 'b) fun_copy, 'a is dead and 'b is live
- ➜ type constructors must be registered as *BNFs** to have live arguments
- ➜ BNF defines well-behaved type constructors, ie where recursion is allowed
- ➜ datatypes automatically are BNFs (that's how they are constructed)
- ➜ can register other type constructors as BNFs — not covered here**

* BNF = Bounded Natural Functors.

** *Defining (Co)datatypes and Primitively (Co)recursive Functions in Isabelle/HOL*

Case

Every datatype introduces a **case** construct, e.g.

$$(\text{case } xs \text{ of } [] \Rightarrow \ldots \mid y \# ys \Rightarrow \ldots\ y \ldots\ ys \ldots)$$

In general: one case per constructor

- ➜ Nested patterns allowed: $x \# y \# zs$
- ➜ Dummy and default patterns with _
- ➜ Binds weakly, needs () in context

Cases

apply (case_tac t)

creates k subgoals

$$\llbracket t = C_i \; x_1 \ldots x_p; \ldots \rrbracket \Longrightarrow \ldots$$

one for each constructor C_i

DEMO

RECURSION

Why nontermination can be harmful

How about $f\ x = f\ x + 1$?

Subtract $f\ x$ on both sides.

$$\Longrightarrow$$
$$0 = 1$$

! All functions in HOL must be total !

Primitive Recursion

primrec guarantees termination structurally

Example primrec def:

primrec app :: "'a list ⇒ 'a list ⇒ 'a list"
where
"app Nil ys = ys" |
"app (Cons x xs) ys = Cons x (app xs ys)"

The General Case

If τ is a datatype (with constructors $C_1, \ldots, C_k$) then $f :: \tau \Rightarrow \tau'$ can be defined by **primitive recursion**:

$$
\begin{array}{lcl}
f\ (C_1\ y_{1,1}\ \ldots\ y_{1,n_1}) & = & r_1 \\
\vdots & & \\
f\ (C_k\ y_{k,1}\ \ldots\ y_{k,n_k}) & = & r_k
\end{array}
$$

The recursive calls in r_i must be **structurally smaller**
(of the form $f\ a_1\ \ldots\ y_{i,j}\ \ldots\ a_p$)

How does this Work?

primrec just fancy syntax for a **recursion operator**

Example:

rec_list :: "'a ⇒ ('b ⇒ 'b list ⇒ 'a ⇒ 'a) ⇒ 'b list ⇒ 'a"

$$\text{rec\_list}\ f_1\ f_2\ \text{Nil} = f_1$$

$$\text{rec\_list}\ f_1\ f_2\ (\text{Cons}\ x\ xs) = f_2\ x\ xs\ (\text{rec\_list}\ f_1\ f_2\ xs)$$

$$\text{app} \equiv \text{rec\_list}\ (\lambda ys.\ ys)\ (\lambda x\ xs\ xs'.\ \lambda ys.\ \text{Cons}\ x\ (xs'\ ys))$$

primrec app :: "'a list ⇒ 'a list ⇒ 'a list"
where
"app Nil ys = ys" |
"app (Cons x xs) ys = Cons x (app xs ys)"

rec_list

Defined: automatically, first inductively (set), then by epsilon

$$\frac{}{(\text{Nil}, f_1) \in \text{list\_rel } f_1 \ f_2} \qquad \frac{(xs, xs') \in \text{list\_rel } f_1 \ f_2}{(\text{Cons } x \ xs, f_2 \ x \ xs \ xs') \in \text{list\_rel } f_1 \ f_2}$$

$$\text{rec\_list } f_1 \ f_2 \ xs \equiv \text{THE } y.\ (xs, y) \in \text{list\_rel } f_1 \ f_2$$

Automatic proof that set def indeed is total function

(the equations for rec_list are lemmas!)

Predefined Datatypes

nat is a datatype

datatype nat $=$ 0 $|$ Suc nat

Functions on nat definable by primrec!

primrec

$$f\ 0 \quad = \quad \ldots$$

$$f\ (\text{Suc}\ n) \quad = \quad \ldots\ f\ n\ \ldots$$

Option

datatype 'a option = None | Some 'a

Important application:

'b ⇒ 'a option	∼	partial function:
None	∼	no result
Some *a*	∼	result *a*

Example:
primrec lookup :: 'k ⇒ ('k × 'v) list ⇒ 'v option
where
lookup k [] = None |
lookup k (x #xs) = (if fst x = k then Some (snd x) else lookup k xs)

DEMO

PRIMREC

INDUCTION

Structural induction

$P\ xs$ holds for all lists xs if

- ➜ P Nil
- ➜ and for arbitrary x and xs, $P\ xs \Longrightarrow P\ (x \# xs)$
  Induction theorem **list.induct:**
  $[\![P\ [];\ \bigwedge a\ list.\ P\ list \Longrightarrow P\ (a \# list)]\!] \Longrightarrow P\ list$
- ➜ General proof method for induction: **(induct x)**
  - x must be a free variable in the first subgoal.
  - type of x must be a datatype.

Basic heuristics

Theorems about recursive functions are proved by induction

Induction on argument number i of f
if f is defined by recursion on argument number i

Example

A tail recursive list reverse:

primrec itrev :: 'a list $\Rightarrow$ 'a list $\Rightarrow$ 'a list
where
itrev [] $\quad ys = ys$ |
itrev $(x\#xs) \quad ys =$ itrev xs $(x\#ys)$

lemma itrev xs [] $=$ rev xs

Demo

Proof Attempt

Generalisation

Replace constants by variables

lemma itrev *xs ys* = rev *xs*@*ys*

Quantify free variables by $\forall$
(except the induction variable)

lemma $\forall$*ys*. itrev *xs ys* = rev *xs*@*ys*

Or: **apply (induct xs arbitrary: ys)**

We have seen today ...

- ➜ Datatypes
- ➜ Primitive recursion
- ➜ Case distinction
- ➜ Structural Induction

Exercises

- define a primitive recursive function **lsum** :: nat list ⇒ nat that returns the sum of the elements in a list.
- show "$2 * \text{lsum } [0..< \mathit{Suc}\ n] = n * (n + 1)$"
- show "lsum (replicate n a) $= n * a$"
- define a function **lsumT** using a tail recursive version of listsum.
- show that the two functions are equivalent: lsum *xs* = lsumT *xs*

