COMP4161: Advanced Topics in Software Verification

C

Gerwin Klein, June Andronick, Ramana Kumar, Miki Tanaka

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Last Time

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➜ Deep and shallow embeddings

Last Time

- ➜ Deep and shallow embeddings
- ➜ Isabelle records

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- Deep and shallow embeddings
- Isabelle records
- Nondeterministic State Monad with Failure

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- Isabelle records
- Nondeterministic State Monad with Failure
- Monadic Weakest Predondition Rules

Content

- ➜ Intro & motivation, getting started [1]
- ➜ Foundations & Principles
  - Lambda Calculus, natural deduction [1,2]
  - Higher Order Logic [3[a]]
  - Term rewriting [4]
- ➜ Proof & Specification Techniques
  - Inductively defined sets, rule induction [5]
  - Datatypes, recursion, induction [6, 7]
  - Hoare logic, proofs about programs, C verification [8[b],9]
  - (mid-semester break)
  - Writing Automated Proof Methods [10]
  - Isar, codegen, typeclasses, locales [11[c],12]

[a] a1 due; [b] a2 due; [c] a3 due

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apply (wp *extra_wp_rules*)

Tactic for automatic application of **weakest precondition rules**

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Today we will learn about AutoCorres and C verification.

Demo

Introduction to AutoCorres and wp

A Brief Overview of C and Simpl

C

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- ➜ expressions with side effects
- ➜ more control flow (do/while, for, break, continue, return)
- ➜ local variables and blocks
- ➜ functions & procedures
- ➜ concrete C data types
- ➜ C memory model and C pointers

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AutoCorres will help, later.

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C Parser: parses C, produces Simpl definitions in Isabelle

- ➜ written by Michael Norrish, NICTA and ANU
- ➜ Handles a non-trivial subset of C
- ➜ Originally written to verify seL4's C implementation
- ➜ AutoCorres is built on top of the C Parser

Commands in Simpl

```
datatype ('s, 'p, 'f) com =
      Skip
    | Basic "'s ⇒ 's"
    | Spec "('s * 's) set"
    | Seq "('s, 'p, 'f) com" "('s, 'p, 'f) com"
    | Cond "'s set" "('s, 'p, 'f) com"  "('s, 'p, 'f) com"
    | While "'s set" "('s, 'p, 'f) com"
    | Call 'p
    | DynCom "'s ⇒ ('s, 'p, 'f) com"
    | Guard 'f "'s set" "('s, 'p, 'f) com"
    | Throw
    | Catch "('s, 'p, 'f) com" "('s, 'p, 'f) com"
```

's = state, **'p** = procedure names, **'f** = faults

Expressions with side effects

`a = a * b;`   `x = f(h);`   `i = ++i - i++;`   `x = f(h) + g(x);`

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Alternative:

Explicitly model nondeterministic order of execution in expressions.

Control flow

```
do { c } while (condition);
```

Already can treat normal while-loops! Automatically translate into:

```
c; while (condition) { c }
```

Similarly:

```
for (init; condition; increment) { c }
```

becomes

```
init; while (condition) { c; increment; }
```


More control flow: break/continue

```
while (condition) {
  foo;
  if (Q) continue;
  bar;
  if (P) break;
}
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- ➜ throw exception **'continue'**, catch at end of body.

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Non-local control flow: **continue** goes to condition, **break** goes to end.
Can be modelled with exceptions:

- ➜ throw exception **'continue'**, catch at end of body.
- ➜ throw exception **'break'**, catch after loop.

Break/continue

Break/continue example becomes:

```
try {
    while (condition) {
        try {
            foo;
            if (Q) { exception = 'continue'; throw; }
            bar;
            if (P) { exception = 'break'; throw; }
        } catch { if (exception == 'continue') SKIP else throw; }
    }
} catch { if (exception == 'break') SKIP else throw; }
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This is not C any more. But it models C behaviour!

Need to be careful that only the translation has access to exception state.

Return

```
if (P) return x;
foo;
return y;
```

Similar non-local control flow.

Return

```
if (P) return x;
foo;
return y;
```

Similar non-local control flow. **Similar solution:** use throw/try/catch

```
try {
    if (P) { return_val = x; exception = 'return'; throw; }
    foo;
    return_val = y; exception = 'return'; throw;
} catch {
    SKIP
}
```


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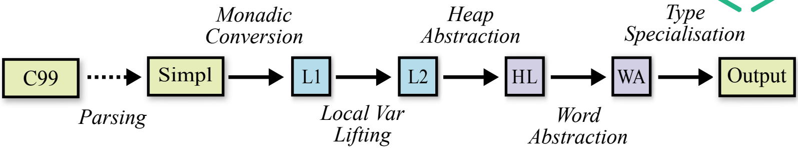
For each Simpl definition C and its generated shallow embedding A:

- ➜ AutoCorres proves an Isabelle theorem stating that C **refines** A
- ➜ Every behaviour of C has a corresponding behaviour of A
- ➜ Refinement guarantees that properties proved about A will also hold for C.
- ➜ (Provided that A never fails. c.f. Total Correctness)

AutoCorres Process



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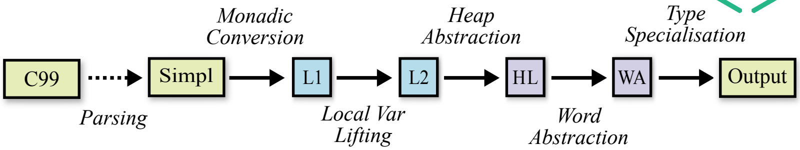

L1: initial monadic shallow embedding

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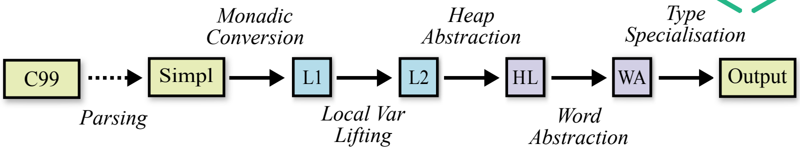

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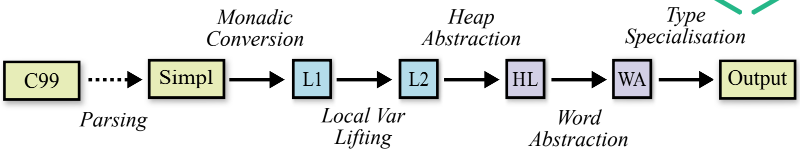

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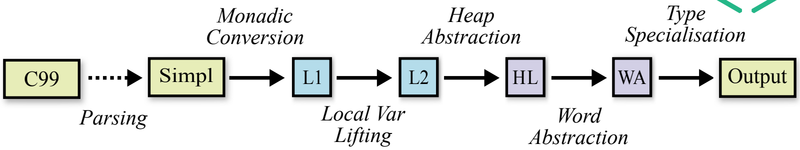

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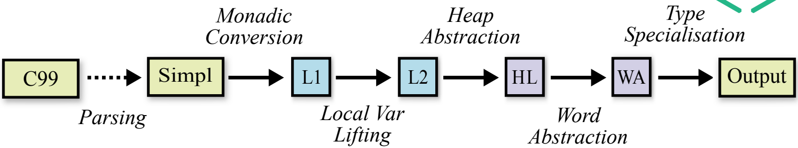

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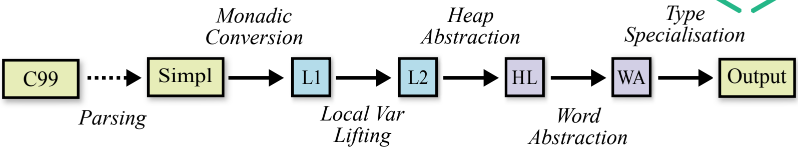

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On-the-fly proof:
Simpl refines **L1** refines **L2** refines **HL** refines **WA** refines **Output**

Example: C99

We will use the following example program to illustrate each of the phases.

```
unsigned some_func(unsigned *a, unsigned *b, unsigned c) {
  unsigned *p = NULL;

  if (c > 10u){
    p = a;
  } else {
    p = b;
  }

  return *p;
}
```

Example: Simpl

```
some_func_body ≡
TRY
  ´p :== ptr_coerce (Ptr (scast 0));;
  IF 0xA < ´c THEN
    ´p :== ´a
  ELSE
    ´p :== ´b
  FI;;
  Guard C_Guard {|c_guard ´p|}
   (creturn global_exn_var_'_update ret__unsigned_'_update
     (λs. h_val (hrs_mem (t_hrs_' (globals s))) (p_' s)));;
  Guard DontReach {} SKIP
CATCH SKIP END
```


Example: L1 (monadic shallow embedding)

```
l1_some_func  ≡  L1_seq (L1_init ret__unsigned_'_update)
 (L1_seq (L1_modify (p_'_update (λ_. ptr_coerce (Ptr (scast 0)))))
   (L1_seq (L1_condition (λs. 0xA < c_' s)
                         (L1_modify (λs. s(|p_' := a_' s|)))
                         (L1_modify (λs. s(|p_' := b_' s|))))
     (L1_seq (L1_guard (λs. c_guard (p_' s)))
       (L1_seq (L1_modify (λs. s(|ret__unsigned_' :=
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```

State type is the same as Simpl, namely a record with fields:

- ➜ **globals**: heap and type information
- ➜ **a_'**, **b_'**, **c_'**, **p_'** (parameters and local variables)
- ➜ **ret__unsigned_'**, **global_exn_var_'** (return value, exception type)

Example: L2 (local variables lifted)

```
l2_some_func a b c ≡
L2_seq (L2_condition (λs. 0xA < c)
                     (L2_gets (λs. a) [''p''])
                     (L2_gets (λs. b) [''p'']))
  (λp. L2_seq (L2_guard (λs. c_guard p))
     (λ_. L2_gets (λs. h_val (hrs_mem (t_hrs_' s)) p) [''ret'']))
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```

State is a record with just the **globals** field

- ➜ function now takes its parameters as arguments
- ➜ local variable **p** now passed via λ-binding
- ➜ **L2_gets** annotated with local variable names
- ➜ This ensures preservation by later AutoCorres phases

Example: HL (heap abstracted into typed heaps)

```
hl_some_func a b c ≡
L2_seq (L2_condition (λs. 0xA < c)
                      (L2_gets (λs. a) [''p''])
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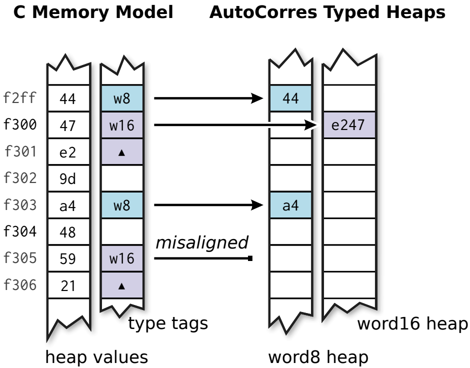
State is a record with a set of **is_valid_** and **heap_** fields:

- ➜ These store **pointer validity** and **heap contents** respectively, per type
- ➜ above example has only 32-bit word pointers

Heap Abstraction



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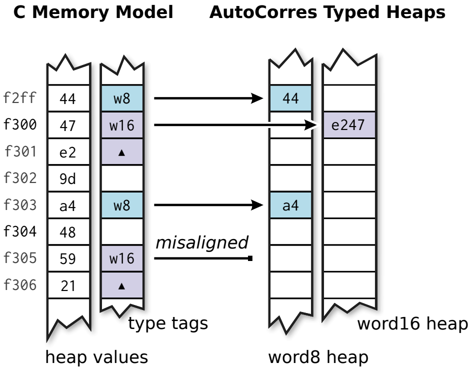

C Memory Model: by Harvey Tuch

- ➜ **Heap** is a mapping from 32-bit addresses to bytes: 32 word $\Rightarrow$ 8 word
- ➜ **Heap Type Description** stores type information for each heap location

Example: WA (words abstracted to ints and nats)

```
wa_some_func a b c ≡
L2_seq (L2_condition (λs. 10 < c)
                     (L2_gets (λs. a) [''p''])
                     (L2_gets (λs. b) [''p'']))
 (λr. L2_seq (L2_guard (λs. is_valid_w32 s r))
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Word abstraction: C **int** → Isabelle int, C **unsigned** → Isabelle nat

- ➜ Guards inserted to ensure absence of unsigned underflow and overflow
- ➜ Signed under/overflow already has guards, because is undefined behaviour

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In the example, the **unsigned** argument **c** is now of type **nat**

- ➜ The function also returns a nat result
- ➜ The heap is not abstracted, hence the call to **unat**

Example: Output (type strengthening and polish)

```
some_func' a b c ≡
DO p ← oreturn (if 10 < c then a else b);
   oguard (λs. is_valid_w32 s p);
   ogets (λs. unat (heap_w32 s p))
OD
```


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Type Strengthening:

- ➜ Tries to convert output to a more restricted monad
- ➜ The above is in the **option** monad because it doesn't modify the state, but might fail
- ➜ The **type** of the option monad implies it cannot modify state

Example: Output (type strengthening and polish)

```
some_func' a b c ≡
DO p ← oreturn (if 10 < c then a else b);
   oguard (λs. is_valid_w32 s p);
   ogets (λs. unat (heap_w32 s p))
OD
```

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Polish:

- Simplify output as much as possible
- The **condition** has been rewritten to a **return** because the condition **10** $<$ **c** doesn't depend on the state

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Can be controlled by the **ts_force** option of AutoCorres

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Hoare Rules for Exceptions

New kind of Hoare triples to model normal and exceptional cases:

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Weakest Precondition Rules:

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(other rules analogous)

Today we have seen

- The automated proof method **wp**
- The C Parser and translating C into Simpl
- AutoCorres and translating Simpl into monadic form
- The option and exception monads