COMP 4161

NICTA Advanced Course

Advanced Topics in Software Verification

Gerwin Klein, June Andronick, Toby Murray, Rafal Kolanski

Slide 1

Content

- ➜ Intro & motivation, getting started [1]
- ➜ Foundations & Principles
  - Lambda Calculus, natural deduction [1,2]
  - Higher Order Logic [3]
  - Term rewriting [4[a]]
- ➜ Proof & Specification Techniques
  - Inductively defined sets, rule induction [5]
  - Datatypes, recursion, induction [6, 7]
  - Code generation, type classes [7]
  - Hoare logic, proofs about programs, refinement [8[b],9[c],10]
  - Isar, locales [11[d],12]

[a] a1 due; [b] a2 due; [c] session break; [d] a3 due

Slide 2

Overview

Automatic Proof and Disproof

- ➜ Sledgehammer: automatic proofs
- ➜ Quickcheck: counter example by testing
- ➜ Nipick: counter example by SAT

Based on slides by Jasmin Blanchette, Lukas Bulwahn, and Tobias Nipkow (TUM).

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Automation

Dramatic improvements in fully automated proofs in the last 2 decades.

- ➜ First-order logic (ATP): Otter, Vampire, E, SPASS
- ➜ Propositional logic (SAT): MiniSAT, Chaff, RSat
- ➜ SAT modulo theory (SMT): CVC3, Yices, Z3

The key:

Efficient reasoning engines, and restricted logics.

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Automation in Isabelle

1980s rule applications, write ML code

1990s simplifier, automatic provers (blast, auto), arithmetic

2000s embrace external tools, but don't trust them (ATP/SMT/SAT)

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Sledgehammer

Slegehammer:

- ➜ Connects Isabelle with ATPs and SMT solvers:
  E, SPASS, Vampire, CVC3, Yices, Z3
- ➜ Simple invocation:
  - ➜ Users don't need to select or know facts
  - ➜ or ensure the problem is first-order
  - ➜ or know anything about the automated prover
- ➜ Exploits local parallelism and remote servers

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DEMO: SLEDGEHAMMER

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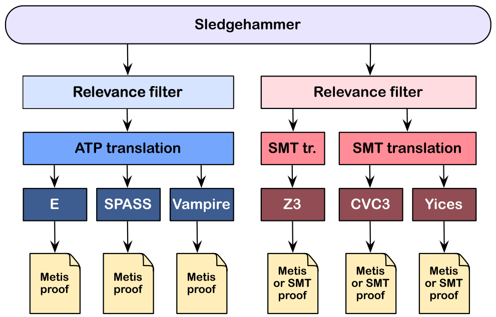
Sledgehammer Architecture



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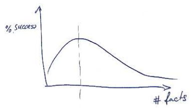
Fact Selection

Provers perform poorly if given 1000s of facts.

- ➜ Best number of facts depends on the prover
- ➜ Need to take care which facts we give them
- ➜ Idea: order facts by relevance, give top n to prover $(n = 250, 1000, \ldots)$
- ➜ Meng & Paulson method: lightweight, symbol-based filter
- ➜ Machine learning method:
  look at previous proofs to get a probability of relevance



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From HOL to FOL

Source: higher-order, polymorphism, type classes
Target: first-order, untyped or simply-typed

- ➜ **First-order:**
  - ➜ SK combinators, λ-lifting
  - ➜ Explicit function application operator
- ➜ **Encode types:**
  - ➜ Monomorphise (generate multiple instances), or
  - ➜ Encode polymorphism on term level

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Reconstruction

We don't want to trust the external provers.
Need to check/reconstruct proof.

- ➜ Re-find using Metis
  Usually fast and reliable (sometimes too slow)
- ➜ Rerun external prover for trusted replay
  Used for SMT. Re-runs prover each time!
- ➜ Recheck stored explicit external representation of proof
  Used for SMT, no need to re-run. Fragile.
- ➜ Recast into structured Isar proof
  Fast, experimental.

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Judgement Day

Evaluating Sledgehammer:

- ➜ 1240 goals out of 7 existing theories.
- ➜ How many can sledgehammer solve?
- ➜ **2010:** E, SPASS, Vampire (for 5-120s). 46%
  $ESV \times 5s \approx V \times 120s$
- ➜ **2011:** Add E-SInE, CVC2, Yices, Z3 (30s).
  $Z3 > V$
- ➜ **2012:** Better integration with SPASS. 64%
  SPASS best (small margin)
- ➜ **2013:** Machine learning for fact selection. 69%
  Improves a few percent across provers.

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Evaluation

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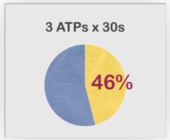
2010



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Evaluation

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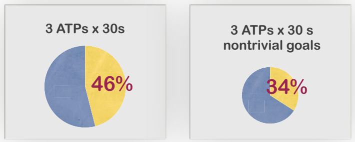
2010



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Evaluation

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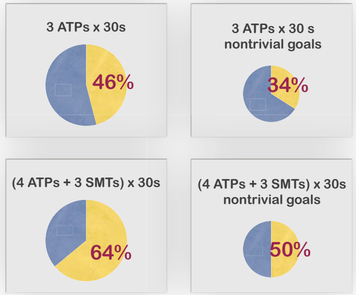
2010

2012



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Sledgehammer rules!

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Example application:

- ➜ Large Isabelle/HOL repository of algebras for modelling imperative programs (Kleene Algebra, Hoare logic, . . ., ≈ 1000 lemmas)
- ➜ Intricate refinement and termination theorems
- ➜ Sledgehammer and Z3 automate algebraic proofs at textbook level.

"The integration of ATP, SMT, and Nitpick is for our purposes very very helpful." – G. Struth

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DISPROOF

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Theorem proving and testing

Testing can show only the presence of errors, but not their absence. (Dijkstra)

Testing cannot prove theorems, but it can refute conjectures!

Sad facts of life:

- ➜ Most lemma statements are wrong the first time.
- ➜ Theorem proving is expensive as a debugging technique.

Find counter examples automatically!

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Quickcheck

Lightweight validation by testing.

- ➜ Motivated by Haskell's QuickCheck
- ➜ Uses Isabelle's code generator
- ➜ Fast
- ➜ Runs in background, proves you wrong as you type.
  (current version: PG only, next release also in jEdit)

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Quickcheck

Covers a number of testing approaches:

- ➜ Random and exhausting testing.
- ➜ Smart test data generators.
- ➜ Narrowing-based (symbolic) testing.

Creates test data generators automatically.

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DEMO: QUICKCHECK

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Test generators for datatypes

Fast iteration in continuation-passing-style

datatype α list = Nil | Cons α (α list)

Test function:

$\text{test}_{\alpha\ list}$ P = P Nil *andalso* test_{α} (λx. $\text{test}_{\alpha\ list}$ (λxs. P (Cons x xs)))

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Test generators for predicates

distinct xs $\Longrightarrow$ distinct (remove1 x xs)

Problem:
Exhaustive testing creates many useless test cases.

Solution:
Use definitions in precondition for smarter generator.
Only generate cases where *distinct xs* is true.

$\text{test-distinct}_{\alpha\ list}$ P = P Nil *andalso*
test_{α} (λx. $\text{test-distinct}_{\alpha\ list}$ (if x $\notin$ xs then (λxs. P (Cons x xs)) else True))

Use data flow analysis to figure out which variables
must be computed and which generated.

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Narrowing

Symbolic execution with demand-driven refinement
➜ Test cases can contain variables
➜ If execution cannot proceed: instantiate with further symbolic terms

Pays off if large search spaces can be discarded:

distinct (Cons 1 (Cons 1 x))

False for any *x*, no further instantiations for x necessary.

Implementation:
Lazy execution with outer refinement loop.
Many re-computations, but fast.

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Quickcheck Limitations

Only executable specifications!

➜ No equality on functions with infinite domain
➜ No axiomatic specifications

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NITPICK

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Nitpick

Finite model finder

➜ Based on SAT via Kodkod (backend of Alloy prover)
➜ Soundly approximates infinite types

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Nitpick Successes

➜ Algebraic methods
➜ C++ memory model
➜ Found soundness bugs in TPS and LEO-II

Fan mail:

"Last night I got stuck on a goal I was sure was a theorem. After 5–10 minutes I gave Nitpick a try, and within a few secs it had found a splendid counterexample—despite the mess of locales and type classes in the context!"

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DEMO: NITPICK

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We have seen today ...

- ➜ Proof: Sledgehammer
- ➜ Counter examples: Quickcheck
- ➜ Counter examples: Nitpick

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