



COMP 4161
NICTA Advanced Course

Advanced Topics in Software Verification

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DATATYPES IN ISAR

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Content



→ Intro & motivation, getting started

→ Foundations & Principles

- Lambda Calculus, natural deduction [2,3,4^a]
- Higher Order Logic [5,6^b,7]
- Term rewriting [8,9,10^c]

→ Proof & Specification Techniques

- Isar [11,12^d]
- Inductively defined sets, rule induction [13^e,15]
- Datatypes, recursion, induction [16,17^f,18,19]
- Calculational reasoning, mathematics style proofs [20]
- Hoare logic, proofs about programs [21^g,22,23]

^aa1 out; ^ba1 due; ^ca2 out; ^da2 due; ^esession break; ^fa3 out; ^ga3 due

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Datatype case distinction



```

proof (cases term)
  case Constructor1
  ⋮
next
  ⋮
next
  case (Constructork  $\vec{x}$ )
  ...  $\vec{x}$  ...
qed

```

case (Constructor_{*i*} \vec{x}) ≡
fix \vec{x} **assume** Constructor_{*i*} : "*term* = Constructor_{*i*} \vec{x} "

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Structural induction for type nat



```

show  $P\ n$ 
proof (induct  $n$ )
  case 0       $\equiv$  let ?case =  $P\ 0$ 
  ...
  show ?case
next
  case (Suc  $n$ )  $\equiv$  fix  $n$  assume Suc:  $P\ n$ 
  ...          let ?case =  $P\ (\text{Suc } n)$ 
  ...  $n$  ...
  show ?case
qed

```

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DEMO: DATATYPES IN ISAR



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Structural induction with \implies and \wedge



```

show " $\wedge x. A\ n \implies P\ n$ "
proof (induct  $n$ )
  case 0       $\equiv$  fix  $x$  assume 0: " $A\ 0$ "
  ...          let ?case = " $P\ 0$ "
  show ?case
next
  case (Suc  $n$ )  $\equiv$  fix  $n$  and  $x$ 
  ...          assume Suc: " $\wedge x. A\ n \implies P\ n$ "
  ...  $n$  ...    "A (Suc  $n$ )"
  ...          let ?case = " $P\ (\text{Suc } n)$ "
  show ?case
qed

```

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DEMO: REGULAR EXPRESSIONS



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We have seen today ...



- Datatypes in Isar
- Defining regular wexpressions as a data type
- Playing with recursion and induction

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