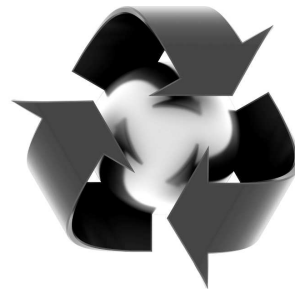


COMP 4161

NICTA Advanced Course

Advanced Topics in Software Verification

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Content

Rough timeline

- Intro & motivation, getting started [1]

- Foundations & Principles
 - Lambda Calculus, natural deduction [2,3,4^a]
 - Higher Order Logic [5,6^b,7]
 - Term rewriting [8,9,10^c]

- Proof & Specification Techniques
 - Isar [11,12^d]
 - Inductively defined sets, rule induction [13^e,15]
 - Datatypes, recursion, induction [16,17^f,18,19]
 - Calculational reasoning, mathematics style proofs [20]
 - Hoare logic, proofs about programs [21^g,22,23]

^a a1 out; ^b a1 due; ^c a2 out; ^d a2 due; ^e session break; ^f a3 out; ^g a3 due

DATATYPES IN ISAR

Datatype case distinction

```
proof (cases term)  
  case Constructor1  
  ⋮  
next  
⋮  
next  
  case (Constructork  $\vec{x}$ )  
  ...  $\vec{x}$  ...  
qed
```

case (Constructor_{*i*} \vec{x}) ≡
fix \vec{x} **assume** Constructor_{*i*} : "*term* = Constructor_{*i*} \vec{x} "

Structural induction for type nat

show $P\ n$

proof (induct n)

case 0 \equiv **let** $?case = P\ 0$

...

show $?case$

next

case (Suc n) \equiv **fix** n **assume** Suc: $P\ n$

...

let $?case = P\ (\text{Suc } n)$

... n ...

show $?case$

qed

Structural induction with \implies and \wedge

show " $\wedge x. A\ n \implies P\ n$ "

proof (induct n)

case 0

\equiv **fix** x **assume** 0: " $A\ 0$ "

...

let $?case = "P\ 0"$

show $?case$

next

case (Suc n)

\equiv **fix** n and x

...

assume Suc: " $\wedge x. A\ n \implies P\ n$ "

... n ...

" $A\ (\text{Suc } n)$ "

...

let $?case = "P\ (\text{Suc } n)"$

show $?case$

qed

DEMO: DATATYPES IN ISAR

DEMO: REGULAR EXPRESSIONS

We have seen today ...

- Datatypes in Isar
- Defining regular wxpressions as a data type
- Playing with recursion and induction