# COMP1917: Computing 1

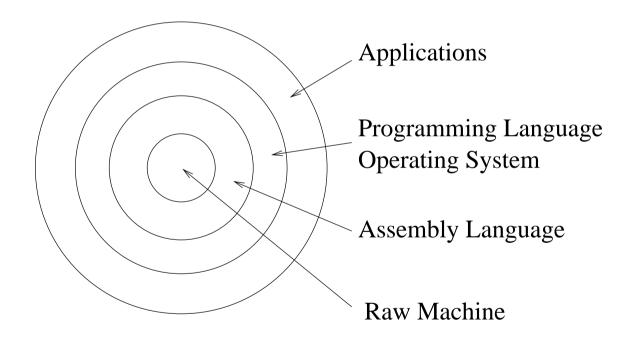
# 17. Memory and Stack Frames

### **Overview**

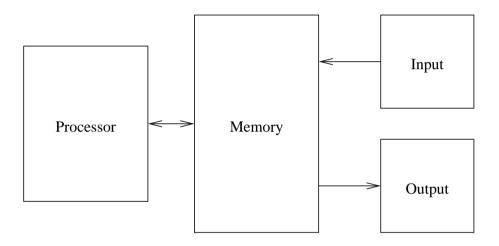
- Computer Systems
- Memory Map
- Static and Dynamic Variables
- Function Calls
- Stack Frames
- Stack Overflow

## **Computer Systems**

Modern computer systems are layered.



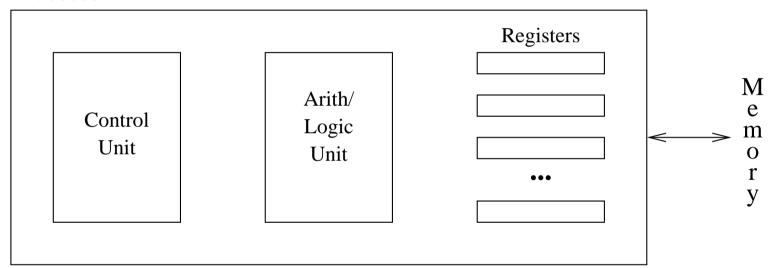
## **Computer Architecture**



- Processor: control, calculation
- Memory: data & program storage
- Input/output: interface to the world

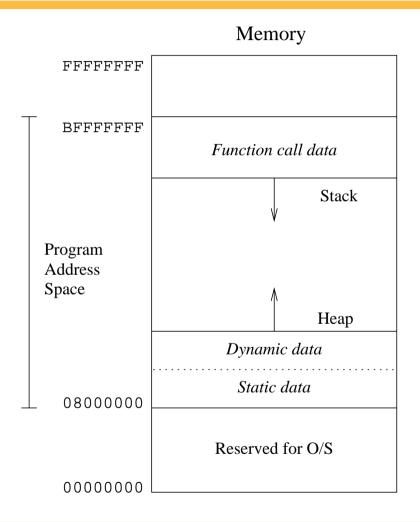
## **Central Processing Unit**

#### Processor



Registers are used as "working memory" to store intermediate values in a computation.

## **Memory Map**



### **Static Variables**

Recall: static variables keep their value from one function call to the next.

What is the output of this code?

```
void inc()
{
    static int k = 5;
    int l = 5;
    int l = 5;
    inc();
    k++;
    l++;
    printf("k = %d, l = %d\n", k, l );
}
```

## **Static and Dynamic Variables**

```
static int k:
  int 1;
  int *a =(int *)malloc( 10 * sizeof( int ));
 printf(" static variable k is stored at %8X\n", &k );
 printf("dynamic variable a is stored at %8X\n", a);
 printf(" local variable l is stored at %8X\n", &l );
Output:
          static variable k is stored at
                                           80496F0
         dynamic variable a is stored at 804A008
          local variable 1 is stored at BFD710D0
```

### **Function Calls**

If main calls f calls g calls h ...

Then h finishes, then g finishes, then f finishes and we're back in main.

Function call/return is last-in, first-out (LIFO) protocol.

 $\Rightarrow$  use a stack of return addresses.

Memory	Memory	Memory	Memory	Memory
f	f	f	f	f
	g	g	g	
		h		
inside f	inside g	inside h	inside g	inside f

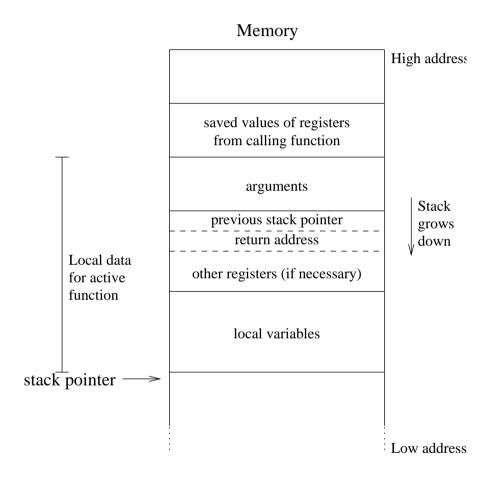
#### **Stack Frames**

- Inside a function, we need access to:
  - arguments
  - local variables
- When the function terminates, need to retrieve
  - return value
  - register values from previous function
  - previous stack pointer
  - return address

All of these are located on the stack. Thus, a small region on the stack is associated with the invocation of each function.

This is called a stack frame.

## gcc **Stack Frame**



### **Creating a Stack Frame**

#### On entry to a function:

- 1. compute size of stack frame and new stack pointer
- 2. allocate memory for frame
- 3. save registers
- 4. store arguments
- 5. save previous stack pointer and return address
- 6. change stack pointer
- 7. pass control to new function

## Removing a Stack Frame

#### On exit from a function:

- 1. save the return value
- 2. restore previous register values
- 3. pop stack frame by reverting to previous stack pointer
- 4. restore control to previous function by jumping to return address

#### factorial.c

```
int factorial( int n )
  printf("n at %X is equal to %d\n", &n, n);
   if( n <= 1 )
       return( 1 );
  else
       return( n * factorial( n-1 ));
int main( void )
   int fact; int n;
  printf("Enter number: ");
   scanf( "%d", &n );
  fact = factorial( n );
  printf("Factorial of %d is %d\n", n, fact );
```

### Output of factorial.c

```
Enter number: 4

n at BFD8E1B0 is equal to 4

n at BFD8E190 is equal to 3

n at BFD8E170 is equal to 2

n at BFD8E150 is equal to 1

Factorial of 4 is 24
```

### **Stack Overflow**

Enter number: 1000000

```
n at BFD8E1B0 is equal to 1000000
n at BFD8E190 is equal to 999999
n at BFD8E170 is equal to 999998
n at BFD8E150 is equal to 999997
n at BF46E700 is equal to 738089
n at BF46E6E0 is equal to 738088
n at BF46E6C0 is equal to 738087
Segmentation fault
```