COMP1917: Computing 1

17. Memory and Stack Frames

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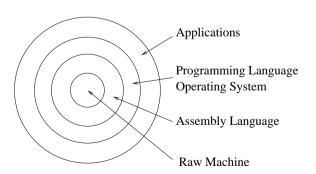
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Computer Systems

Modern computer systems are layered.



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Overview

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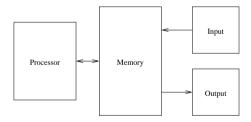
- Computer Systems
- Memory Map
- Static and Dynamic Variables
- Function Calls
- Stack Frames
- Stack Overflow

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Computer Architecture



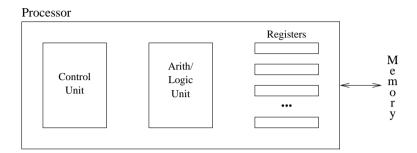
- Processor: control, calculation
- Memory: data & program storage
- Input/output: interface to the world

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Central Processing Unit



Registers are used as "working memory" to store intermediate values in a computation.

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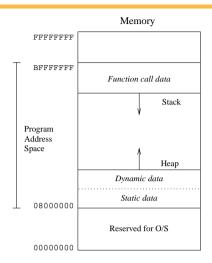
Static Variables

Recall: static variables keep their value from one function call to the next.

What is the output of this code?

```
void inc()
{
    static int k = 5;
    int l = 5;
    inc();
    k++;
    l++;
    printf("k = %d, l = %d\n", k, l );
}
```

Memory Map



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Static and Dynamic Variables

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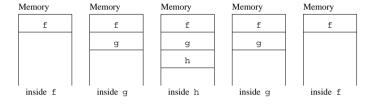
Function Calls

If main calls f calls g calls h ...

Then h finishes, then g finishes, then f finishes and we're back in main.

Function call/return is last-in, first-out (LIFO) protocol.

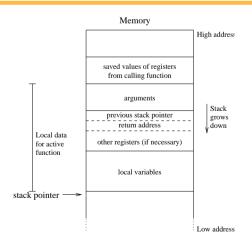
 \Rightarrow use a stack of return addresses.



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gcc Stack Frame



Stack Frames

- Inside a function, we need access to:
 - arguments
 - local variables
- When the function terminates, need to retrieve
 - return value
 - ▶ register values from previous function
 - previous stack pointer
 - return address

All of these are located on the stack. Thus, a small region on the stack is associated with the invocation of each function.

This is called a stack frame.

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Creating a Stack Frame

On entry to a function:

- 1. compute size of stack frame and new stack pointer
- 2. allocate memory for frame
- 3. save registers
- 4. store arguments
- 5. save previous stack pointer and return address
- 6. change stack pointer
- 7. pass control to new function

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Removing a Stack Frame

On exit from a function:

- 1. save the return value
- 2. restore previous register values
- 3. pop stack frame by reverting to previous stack pointer
- 4. restore control to previous function by jumping to return address

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Output of factorial.c

```
Enter number: 4
n at BFD8E1B0 is equal to 4
n at BFD8E190 is equal to 3
n at BFD8E170 is equal to 2
n at BFD8E150 is equal to 1
Factorial of 4 is 24
```

factorial.c

```
int factorial( int n )
   printf("n at %X is equal to %d\n", &n, n );
   if( n <= 1 )
       return(1);
   else
       return( n * factorial( n-1 ));
int main( void )
   int fact; int n;
   printf("Enter number: ");
   scanf( "%d", &n );
   fact = factorial( n );
   printf("Factorial of %d is %d\n", n, fact );
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```

Stack Overflow

```
Enter number: 1000000
n at BFD8E1B0 is equal to 1000000
n at BFD8E190 is equal to 999999
n at BFD8E170 is equal to 999998
n at BFD8E150 is equal to 999997
. . .
n at BF46E700 is equal to 738089
n at BF46E6E0 is equal to 738088
n at BF46E6C0 is equal to 738087
Segmentation fault
```