COMP1521 25T2

Week 7 Lecture 2

File Systems

Adapted from Hammond Pearce, Andrew Taylor and John Shepherd's slides

Announcements

Test 5 and Test 6: due tomorrow

Assignment 1: Automarking starting today.

Assignment 2: Later this week. Announced via email and forum.

Today's Lecture

- Recap
 - File Operations
- Filesystems
 - File metadata
 - Permissions
 - "stat" system call
 - Hard Links and Symbolic Links
 - Working with directories



Recap Files

Question 1: What is better to use to read in a file? fgetc or fgets or fscanf?

Question 2: If I successfully open a file using FILE *f = fopen("data", "w");

- A. What will happen if the file already exists? What if it doesn't?
- B. What is the difference between mode "a" and "w"

Question 3: How many bytes would the following print to the file f:

```
A. fprintf(f, "%d", 255);
B. fputc(f, 255);
```

Recap Files

- These both return *s* on success, NULL on failure
 - o char* gets(char *s);
 - char* fgets(char *s, int size, FILE *stream);

- Revisit:
 - int sprintf(char *str, const char *format, ...)
 - int snprintf(char *str, size_t size, const char *format, ...)
- Do you need to remember *all* of these?

Filesystems

File Systems

File systems manage stored data (e.g. on disk, SSD)

File = named sequence of bytes, stored on device There is no meaning for bytes associated with a file (ASCII, jpg, MP4, ...)

Directory = named set of **files** or **Directories**

- File system maps name to location(s) on device
- File system maintains meta-data (e.g. permissions, time stamps)

Unix/Linux File System



Unix/Linux file system is tree-like

- symlinks actually make it into a graph
- when traversing you may find infinite loops

Unix-like File Names

Sequences of 1 or more bytes

- filenames can contain any byte except
- **0x00** bytes (ASCII '\0') used to terminate filenames
- **0x2F** bytes (ASCII '/') used to separate components of pathnames.
- maximum filename length, depends on file system, typically 255

Two filenames have a special meaning:

- . current directory
- .. parent directory

Some programs (shell, ls) treat filenames starting with . specially.

Paths and directories

Absolute pathnames start with a leading / and give full path from root e.g. - /usr/include/stdio.h

Relative pathnames do **not** start with a leading / e.g.

- ../../another/path/prog.c

- ./a.out
- main.c



Current Working Directory

Every process (running program) has a current working directory (CWD)

- this is an absolute pathname
- this is the directory from where the process was run from
- shell command pwd prints the CWD

Relative pathnames appended to CWD of process e.g.

- if CWD is /home/z5555555/lab07/
- and relative path is main.c
- absolute path would be /home/z5555555/lab07/main.c

Everything is a file

- Originally, files only managed data stored on a magnetic disk
- Unix philosophy: *Everything is a file*
- File systems used to access
 - Files
 - Directories
 - Devices
 - System information
 - Other filesystems

Unix-like File Metadata

Metadata for file system objects is stored in **inodes**, which hold

- location of file contents in file systems
- file type (regular file, directory, ...)
- file size in bytes
- file ownership
- file access permissions who can read, write, execute the file
- timestamps times of file was created, last accessed, last updated

File system implementations often add complexity to improve performance

• e.g. data of small files might be stored in the inode itself.

Inodes

Files system has large table of inodes containing metadata

- Inode-number is the inodes id
 - Unique for file system like zid within UNSW

Directories are effectively a collection of (name, inode-number) pairs

- ls -i prints inode-numbers

File Access: Behind the scenes

Access to files by name proceeds (roughly) as...

- open directory and scan for name
 - if not found, "No such file or directory"
- if found as (name, inumber), access inode table inodes[inumber]
- collect file metadata and...
 - check file access permissions given current user/group
 - if don't have required access, "Permission denied"
 - update access timestamp
- use data in inode (size, location of bytes) to access file contents

File Permissions

Every file and directory in linux has read, write and execute permissions (access rights) for each of the following user groups:

- user: the file's owner
- group: the members of the file's group -
- other: everyone else
- type **ls** -**l** on command line to see





File Permissions: read, write, execute

	Read	Write	Execute
File	View contents of file	Modify file	Run as executable
Directory	View names of file e.g. use ls	Create, delete, rename files	Can cd into it. Also needed to access (read, write, execute) items in directory

Modifying Permissions

You can think of permissions as a set of bits and then each 3 bits as an octal digit. e.g.

rwx r-x r-x 111 101 101 7 5 5

You can use the **chmod** command to set the permissions of a file or directory using the desired 3 digit octal code. e.g.

\$ chmod 700 f.txt

Hard Links and Symbolic Links

File system **links** allow multiple paths to access the same file

- Hard links

- multiple names referencing the same inode (file)
- the two entries must be on the same filesystem (inodes unique to filesystem)
- can not create a (extra) hard link to directories
- all hard links to a file have equal status
- file destroyed when last hard link removed

Assuming 'fileA' already exists, this creates a hard link named 'fileB' In fileA fileB

Hard Links and Symbolic Links

File system **links** allow multiple paths to access the same file

- Symbolic links (symlinks)
 - point to another path name
 - accessing the symlink (by default) accesses the file being pointed to
 - symbolic link can point to a directory
 - symbolic link can point to a pathname on another filesystems
 - symbolic links don't have permissions (not needed they are just a pointer)

Assuming 'fileA' already exists, this creates a symbolic link named 'fileB' **ln -s fileA fileB**

C library wrapper for stat system call

int stat(const char *pathname, struct stat *statbuf);

- returns metadata associated with pathname in statbuf
- metadata returned includes:
 - inode number
 - type (file, directory, symbolic link, device)
 - size of file in bytes (if it is a file)
 - permissions (read, write, execute)
 - times of last access/modification/status-change
- returns -1 and sets errno if metadata not accessible

C library wrapper for stat system call

int lstat(const char *pathname, struct stat *statbuf);

- same as stat() but doesn't follow symbolic links
 - in other words gives you metadata about the symbolic link and not the file it links to
 - important not to get stuck in infinite loops

int fstat(<u>int fd</u>, struct stat *statbuf);

- same as stat() but gets data via an open file descriptor

See man 2 stat man 3 stat man 7 inode

definition of struct stat

man 3 stat

struct stat {

dev_t	st_dev;
ino_t	st_ino;
mode_t	st_mode;
nlink_t	st_nlink
uid_t	st_uid;
gid_t	st_gid;
dev_t	st_rdev;
off_t	st_size;

- /* ID of device containing file */
 /* Inode number */
 /* File type and mode */
- /* Number of hard links */
- /* User ID of owner */
- /* Group ID of owner */
- /* Device ID (if special file) */
- /* Total size, in bytes */

···· };

st_mode field of struct stat

man 7 inode

st_mode is a bitwise-or of these values (& others):

S_IFLNK	0120000	symbolic link
S_IFREG	0100000	regular file
S_IFDIR	0040000	directory
S_IRUSR	0000400	owner has read permission
S_IWUSR	0000200	owner has write permission
S_IXUSR	0000100	owner has execute permission
S_IRGRP	0000040	group has read permission
S_IWGRP	0000020	group has write permission
S_IXGRP	0000010	group has execute permission
S_IROTH	0000004	others have read permission
S_IWOTH	0000002	others have write permission
S_IXOTH	0000001	others have execute permission

Code demos stat.c

stat.c

Another good sample program at bottom of man 2 stat

Making a directory

int mkdir(const char *pathname, mode_t mode);

returns 0 if successful, returns -1 and sets errno otherwise

- for example: mkdir("newDir", 0755)

if pathname is e.g. `a/b/c/d`

- all of the directories `a`, `b` and `c` must exist
- directory `c` must be writable to the caller
- directory `d` must not already exist

the new directory contains two initial entries

- `.` is a reference to itself
- `..` is a reference to its parent directory

Demo: mkdir.c

Opening and Reading directories

// open a directory stream for directory name
DIR *opendir(const char *name);

// return a pointer to next directory entry
struct dirent *readdir(DIR *dirp);

// close a directory stream
int closedir(DIR *dirp);

Found in man 3 Demo list_directory.c

Useful Linux (POSIX) functions

chmod(char *pathname, mode_t mode) // change permission of file/... unlink(char *pathname) // remove a file... **rename(char *oldpath, char *newpath)** // rename a file/directory **chdir(char *path)** // change current working directory getcwd(char *buf, size_t size) // get current working directory **link(char *oldpath, char *newpath)** // create hard link to a file symlink(char *target, char *linkpath) // create a symbolic link

Demo: chmod.c rm.c rename.c my_cd.c getcwd.c nest_directories.c many_links.c chain_links.c

Reach Out

Content Related Questions: Forum

Admin related Questions email: <u>cs1521@cse.unsw.edu.au</u>



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