

# COMP1521 26T1

## Week 4 Lecture 1

# Integers

# Announcements

- **Census Date** : Wednesday this week (12th March)
  - Last day to drop T1 courses without financial liability
- Weekly Test 3 Due: Thursday 9pm
- Assignment 1 Due: Week 5 Friday 18:00 (next week)
- See [Help Session Schedule](#)
- Lab marks usually come out about 2 weeks after deadline

# Today's Lecture

- Invalid C
- Integers



# Code Demos: Invalid C

invalid\_1.c

invalid\_4.c

See more interesting invalid C code examples at

[https://cgi.cse.unsw.edu.au/~cs1521/26T1/topic/mips\\_functions/code/invalid0.c](https://cgi.cse.unsw.edu.au/~cs1521/26T1/topic/mips_functions/code/invalid0.c)

[https://cgi.cse.unsw.edu.au/~cs1521/26T1/topic/mips\\_functions/code/invalid2.c](https://cgi.cse.unsw.edu.au/~cs1521/26T1/topic/mips_functions/code/invalid2.c)

[https://cgi.cse.unsw.edu.au/~cs1521/26T1/topic/mips\\_functions/code/invalid3.c](https://cgi.cse.unsw.edu.au/~cs1521/26T1/topic/mips_functions/code/invalid3.c)

# Integers

# Why Learn About Integers

- Fundamental topic in computing
  - Understand what you are seeing in mipsy web!
  - Understand limits of types and help you understand and debug code
- Prepare you for the next topic bitwise operators
- Understand the jokes in these slides

# There are 10 types of students

Those that understand binary,

And those that don't.

# Numbers

## 4705

It is equivalent to:  $4 \cdot 10^3 + 7 \cdot 10^2 + 0 \cdot 10^1 + 5 \cdot 10^0$   
 $= 4000 + 700 + 0 + 5$

# Numbers

## 4705

It is equivalent to:  $4 \cdot 10^3 + 7 \cdot 10^2 + 0 \cdot 10^1 + 5 \cdot 10^0$   
 $= 4000 + 700 + 0 + 5$

**If we assume it is base 10!**

# Base 10: Decimal

- In Base (or radix) 10 we have 10 digits e.g. 0..9
  - Then to get bigger numbers we start combining the digits e.g. 10
- Place Values

$10^3$	$10^2$	$10^1$	$10^0$
$1000_{10}$	$100_{10}$	$10_{10}$	$1_{10}$

- Example:

$$\begin{aligned} 4705_{10} &= 4 \cdot 10^3 + 7 \cdot 10^2 + 0 \cdot 10^1 + 5 \cdot 10^0 \\ &= 4000 + 700 + 0 + 5 \\ &= 4705_{10} \end{aligned}$$

# Base 10 was an arbitrary choice

- Possibly exists because we have 10 digits (fingers)
- Ancient Egyptians, Brahmi Numerals, Greek Numerals, Hebrew Numerals, Roman Numerals and Chinese Numerals:
  - All base 10!

# Code Demo: Digits.c

How can we get each digit from a number in decimal, one by one?

E.g. How can we get the 2 from 8192?

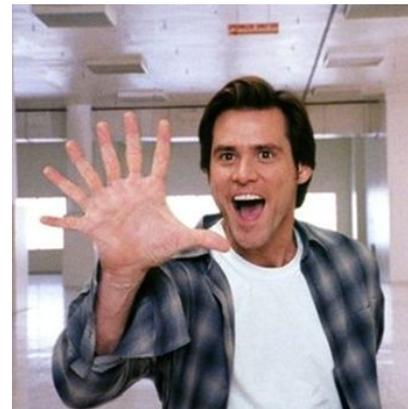
What about the 9?

# What about some other bases?

- Let's think about base 7 (not a very useful base)
- We have 7 digits 0..6
  - Then we start combining the digits
    - e.g.  $10_7$  represents  $7_{10}$
- Place Values

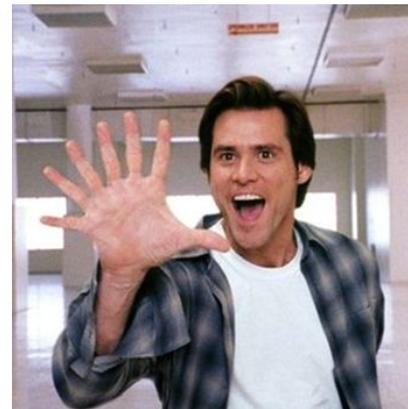
$7^3$	$7^2$	$7^1$	$7^0$
$343_{10}$	$49_{10}$	$7_{10}$	$1_{10}$

Here,  $1216_7 = ?$



# What about some other bases?

- Let's think about base 7 (not a very useful base)
- We have 7 digits 0..6
  - Then we start combining the digits  
e.g.  $10_7$  represents  $7_{10}$
- Place Values



$7^3$	$7^2$	$7^1$	$7^0$
$343_{10}$	$49_{10}$	$7_{10}$	$1_{10}$

$$\begin{aligned}\text{Here, } 1216_7 &= 1 * 7^3 + 2 * 7^2 + 1 * 7^1 + 6 * 7^0 \\ &= 1 * 343 + 2 * 49 + 1 * 7 + 6 * 1 \\ &= 454_{10}\end{aligned}$$

# Base 2: Computers like binary

- In Base (or radix) 2 we have 2 digits (bits) – 0 and 1
  - Easy to represent using “electricity” – **Off** and **On**
  - Then we start combining the digits e.g.  $10_2$  represents  $2_{10}$
- Place Values

$2^3$	$2^2$	$2^1$	$2^0$
$8_{10}$	$4_{10}$	$2_{10}$	$1_{10}$

$$1011_2 = ?_{10}$$

# Base 2: Computers like binary

- In Base (or radix) 2 we have 2 digits (bits) – 0 and 1
  - Easy to represent using “electricity” – **Off** and **On**
  - Then we start combining the digits e.g.  $10_2$  represents  $2_{10}$
- Place Values

$2^3$	$2^2$	$2^1$	$2^0$
$8_{10}$	$4_{10}$	$2_{10}$	$1_{10}$

$$\begin{aligned}1011_2 &= 1 * 2^3 + 0 * 2^2 + 1 * 2^1 + 1 * 2^0 \\ &= 1 \times 8 + 0 \times 4 + 1 \times 2 + 1 \times 1 \\ &= 11_{10}\end{aligned}$$

# More examples

**Question:** Convert  $1101_2$  to decimal

**Question:** Convert  $25_{10}$  to binary

# More examples

**Question:** Convert  $1101_2$  to decimal

**Answer:**  $1 * 2^3 + 1 * 2^2 + 0 * 2^1 + 1 * 2^0$   
 $= 1 \times 8 + 1 \times 4 + 0 \times 2 + 1 \times 1$   
 $= 13$

**Question:** Convert  $25_{10}$  to binary

# More examples

**Question:** Convert  $1101_2$  to decimal

**Answer:**  $1 * 2^3 + 1 * 2^2 + 0 * 2^1 + 1 * 2^0$   
 $= 1 * 8 + 1 * 4 + 0 * 2 + 1 * 1$   
 $= 13$

**Question:** Convert  $25_{10}$  to binary = 11001

- $25/2 = 12 \text{ R } 1$  (this is the least significant bit)
- $12/2 = 6 \text{ R } 0$
- $6/2 = 3 \text{ R } 0$
- $3/2 = 1 \text{ R } 1$
- $1/2 = 0 \text{ R } 1$  (this is the most significant bit)

# Binary numbers are hard to read!

- They get very long, very fast
- E.g.  $12345678_{10} = 101111000110000101001110_2$

# Binary numbers are hard to read!

- They get very long, very fast
- E.g.  $12345678_{10} = 101111000110000101001110_2$
- Solution: Write numbers in hexadecimal!
  - More compact than binary
    - 4 binary digits replaced by 1 hexadecimal digit
    - Bit patterns remain more obvious than in decimal

# Base 16: Hexadecimal

- In Base (or radix) 16 we have 16 digits (0..15)
  - 0 1 2 3 4 5 6 7 8 9 A B C D E F
  - Then we start combining the digits e.g. 10 represents  $16_{10}$
- Place Values

$16^3$	$16^2$	$16^1$	$16^0$
$4096_{10}$	$256_{10}$	$16_{10}$	$1_{10}$

- $3AF1_{16} = ?_{10}$

# Base 16: Hexadecimal

- In Base (or radix) 16 we have 16 digits
  - 0 1 2 3 4 5 6 7 8 9 A B C D E F
  - Then we start combining the digits e.g. 10 represents  $16_{10}$
- Place Values

$16^3$	$16^2$	$16^1$	$16^0$
$4096_{10}$	$256_{10}$	$16_{10}$	$1_{10}$

- $3AF1_{16} = 3 \times 16^3 + 10 \times 16^2 + 15 \times 16^1 + 1 \times 16^0$   
 $= 15089_{10}$

# More hexadecimal examples

**Question:** Convert  $1FF_{16}$  to decimal

**Question:** Convert  $13_{10}$  to hexadecimal

# More hexadecimal examples

**Question:** Convert  $1FF_{16}$  to decimal

**Answer:**  $1 \times 16^2 + 15 \times 16^1 + 15 \times 16^0 = 511_{10}$

**Question:** Convert  $13_{10}$  to hexadecimal

# More hexadecimal examples

**Question:** Convert  $1FF_{16}$  to decimal?

**Answer:**  $1 \times 16^2 + 15 \times 16^1 + 15 \times 16^0 = 511_{10}$

**Question:** Convert  $13_{10}$  to hexadecimal?

**Answer:**  $D_{16}$

# Binary -> Hexadecimal

- Binary gets very long very quick
  - e.g.  $12345678_{10} = 101111000110000101001110_2$
- Solution: Write numbers in hexadecimal!

- Place Values

$16^3$	$16^2$	$16^1$	$16^0$
$4096_{10}$	$256_{10}$	$16_{10}$	$1_{10}$

- Since **16 == 2<sup>4</sup>**
  - Each group of 4 bits to represents 1 hex digit

# Binary -> Hexadecimal

$$101111000110000101001110_2$$
$$= 1011 \ 1100 \ 0110 \ 0001 \ 0100 \ 1110_2$$

# Binary -> Hexadecimal

$$101111000110000101001110_2$$
$$= 1011 \ 1100 \ 0110 \ 0001 \ 0100 \ 1110_2$$

Each 4 bit group can be represented by **one** hexadecimal digit!

<b>Base 10</b>	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
<b>Base 16</b>	F	E	D	C	B	A	9	8	7	6	5	4	3	2	1	0
<b>Base 2</b>	1111	1110	1101	1100	1011	1010	1001	1000	0111	0110	0101	0100	0011	0010	0001	0000

# Binary -> Hexadecimal

$$\begin{aligned} & 101111000110000101001110_2 \\ = & \mathbf{1011} \mathbf{1100} \mathbf{0110} \mathbf{0001} \mathbf{0100} \mathbf{1110}_2 \\ = & \quad \mathbf{B} \quad \mathbf{C} \quad \mathbf{6} \quad \mathbf{1} \quad \mathbf{4} \quad \mathbf{E} \end{aligned}$$

Each 4 bit group can be represented by **one** hexadecimal digit!

<b>Base 10</b>	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
<b>Base 16</b>	F	E	D	C	B	A	9	8	7	6	5	4	3	2	1	0
<b>Base 2</b>	1111	1110	1101	1100	1011	1010	1001	1000	0111	0110	0101	0100	0011	0010	0001	0000

# More examples

Binary  $01101111_2 = ?_{16}$

Hexadecimal  $BAD2_{16} = ?_2$

# More examples

Binary  $01101111_2 = 6F_{16}$

Hexadecimal  $BAD2_{16} = ?_2$

# More examples

Binary  $01101111_2 = 6F_{16}$

Hexadecimal  $BAD2_{16} = 10111010\ 11010010_2$

# Base 8: Octal



- In Base (or radix) 8 we have 8 digits
  - 0 1 2 3 4 5 6 7
  - Then we start combining the digits e.g. 10 represents  $8_{10}$
- Similar advantages to hexadecimal
  - $8 = 2^3$  so group bits into 3:
  - Example:  $72_8 = 111\ 010_2 = 3A_{16} = 58_{10}$

<b>Base 10</b>	7	6	5	4	3	2	1	0
<b>Base 8</b>	7	6	5	4	3	2	1	0
<b>Base 2</b>	111	110	101	100	011	010	001	000

# Binary, Octal, Hexadecimal Summary

- In **binary**, (base 2), each digit represents **1** bit:
  - 01001000111110101011110010010111<sub>2</sub>
- In **octal**, (base 8), each digit represents **3** bits
  - 01 001 000 111 110 101 011 110 010 010 111<sub>2</sub>
  - 1 1 0 7 6 5 3 6 2 2 7<sub>8</sub>
- In **hexadecimal**, (base 16), each digit represents **4** bits:
  - 0100 1000 1111 1010 1011 1100 1001 0111<sub>2</sub>
  - 4 8 F A B C 9 7<sub>16</sub>

# Doing maths in binary

$$4 = 00000100 \quad +$$

$$1 = 00000001$$

-----

$$6 = 00000110 \quad +$$

$$7 = 00000111$$

-----

# Doing maths in binary

$$4 = 00000100 \quad +$$

$$1 = 00000001$$

-----

$$5 = 00000101$$

$$6 = 00000110 \quad +$$

$$7 = 00000111$$

-----

# Doing maths in binary

$$4 = 00000100 \quad +$$

$$1 = 00000001$$

-----

$$5 = 00000101$$

$$6 = 00000110 \quad +$$

$$7 = 00000111$$

-----

$$13 = 00001101$$

# Bits, Bytes and Bases in C

# Constants in C and MIPS assembly

- A number beginning with **0x** is hexadecimal
- A number beginning with **0** is octal
- A number beginning with **0b** is binary
- Otherwise, it is decimal
- Note that %d format specifier means decimal

```
printf("%d", 0x2A); // prints 42
```

```
printf("%d", 052); // prints 42
```

```
printf("%d", 0b101010); // prints 42
```

```
printf("%d", 42); // prints 42
```

# Easy Base Conversions in C

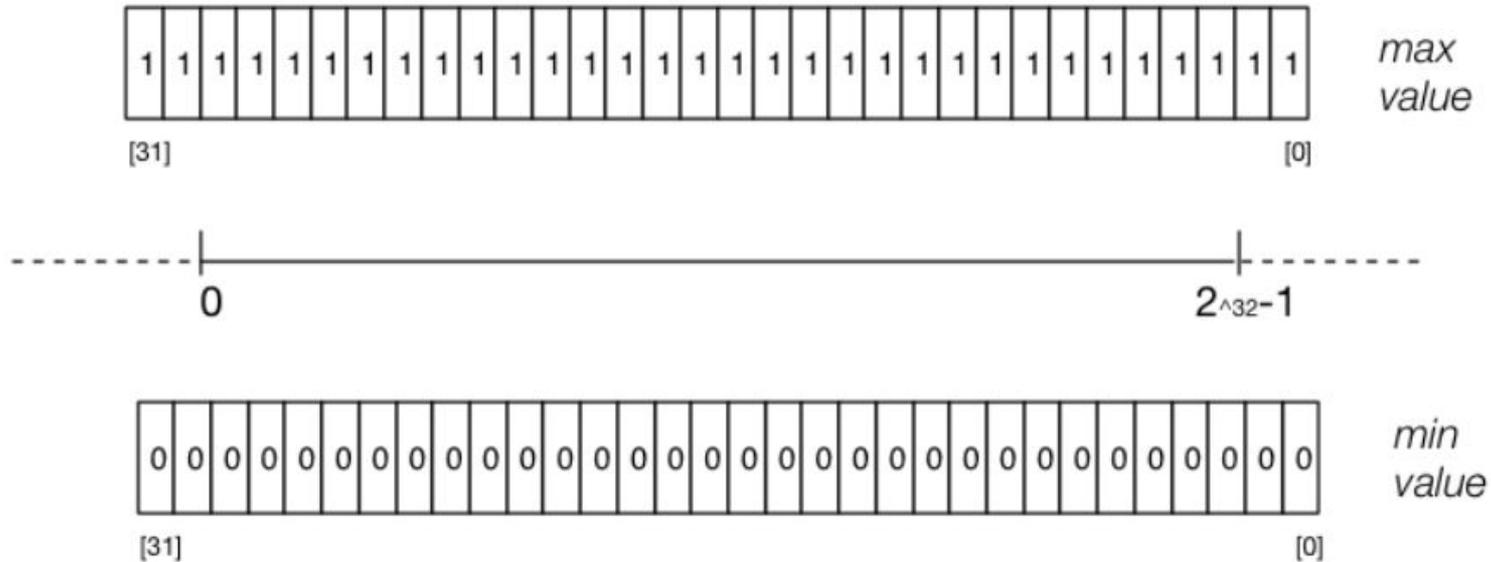
`integer_prefixes.c`

`integer_prefixes_args.c`

# Integer Representations

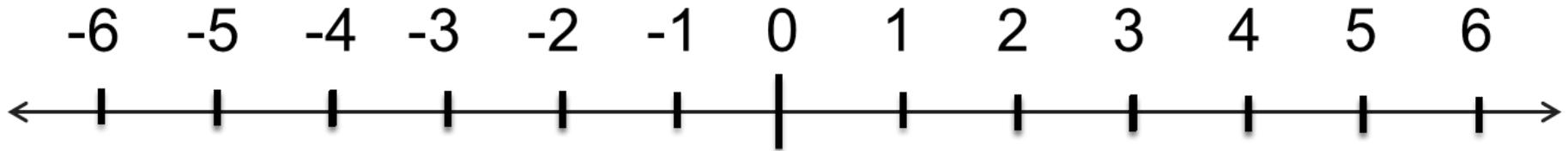
# Unsigned integers

- In C the `unsigned int` data type is 4 bytes on our system
  - means we can store values from the range  $0 \dots 2^{32}-1$

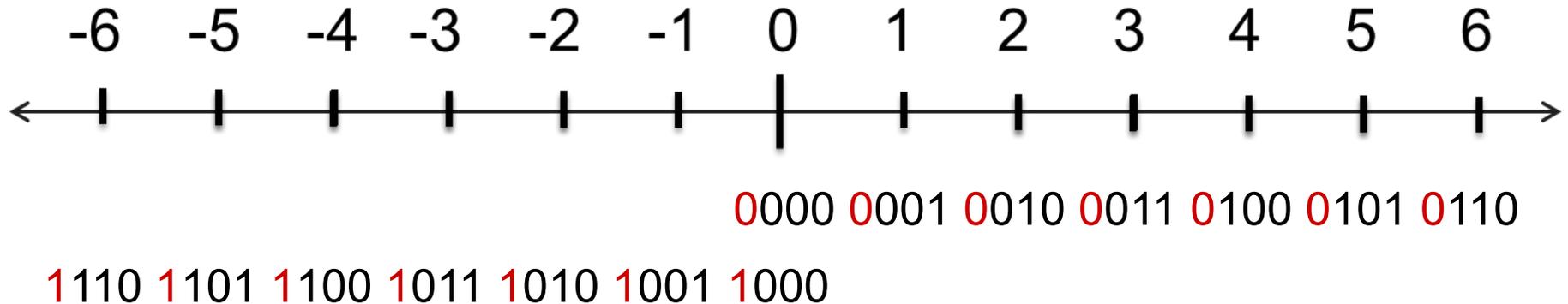


# How do we store signed integers?

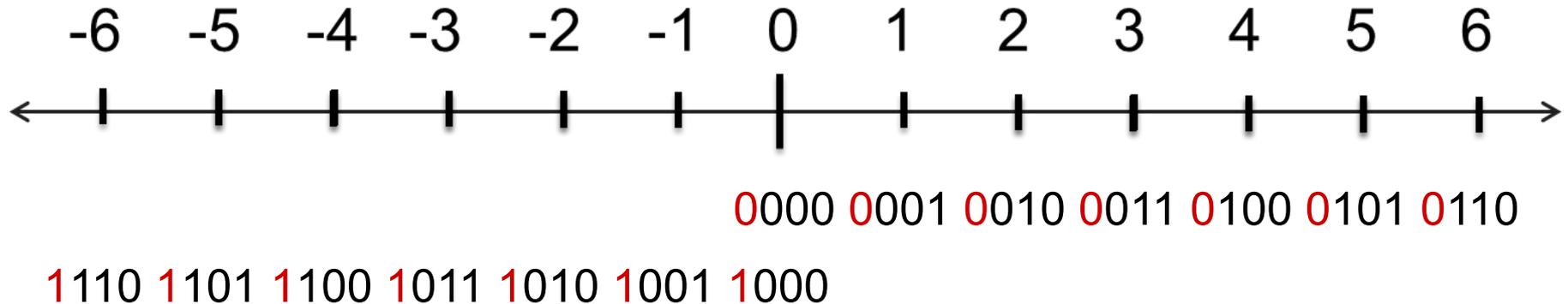
We COULD try to put a leading 0 for positive and a leading 1 for negative numbers



# How do we store signed integers?



# How do we store signed integers?



- Two representations for 0?
- Okay, but what algorithm for adding/subtracting numbers?

# What do signed binary numbers look like?

4 = 00000100 ) -1  
3 = 00000011 ) -1  
2 = 00000010 ) -1  
1 = 00000001 ) -1  
0 = 00000000 ) -1

# What do signed binary numbers look like?

4 = 00000100  
3 = 00000011  
2 = 00000010  
1 = 00000001  
0 = 00000000



Let's keep subtracting

# What do signed binary numbers look like?

4 = 00000100 ) -1  
3 = 00000011 ) -1  
2 = 00000010 ) -1  
1 = 00000001 ) -1  
0 = 00000000 ) -1  
-1 = 11111111 ) -1  
-2 = 11111110 ) -1  
-3 = 11111101 ) -1

# What do signed binary numbers look like?

4 = 00000100 } -1  
3 = 00000011 } -1  
2 = 00000010 } -1  
1 = 00000001 } -1  
0 = 00000000 } -1  
-1 = 11111111 } -1  
-2 = 11111110 } -1  
-3 = 11111101 } -1

This representation is called

## two's complement

- We don't have 2 zeros
- Subtraction works in the same way and so does addition. (Try it)

**Note: For a binary signed number**

Positive numbers start with **0**

Negative numbers start with **1**

# What do signed binary numbers look like?

- Modern computers use **two's complement** for integers
- Positive integers and zero represented as normal
- Negative integers represented in a way to make maths ✨easy✨ for the computer (not humans)
  - For an  $n$ -bit binary number, the number  $-b$  is  $2^n - b$
  - E.g. 8-bit number “-5” is represented as  $2^8 - 5 = 251 = 1111\ 1011_2$
  - Conversely,  $1111\ 1011_2$  represents “-5” because  $251 - 2^8 = -5$ .

# Two's Complement Tips and Tricks

- Another way of thinking about 2's complement
  - If you are trying to represent -5 in 8 bits

Take the +5 representation: 0000 0101

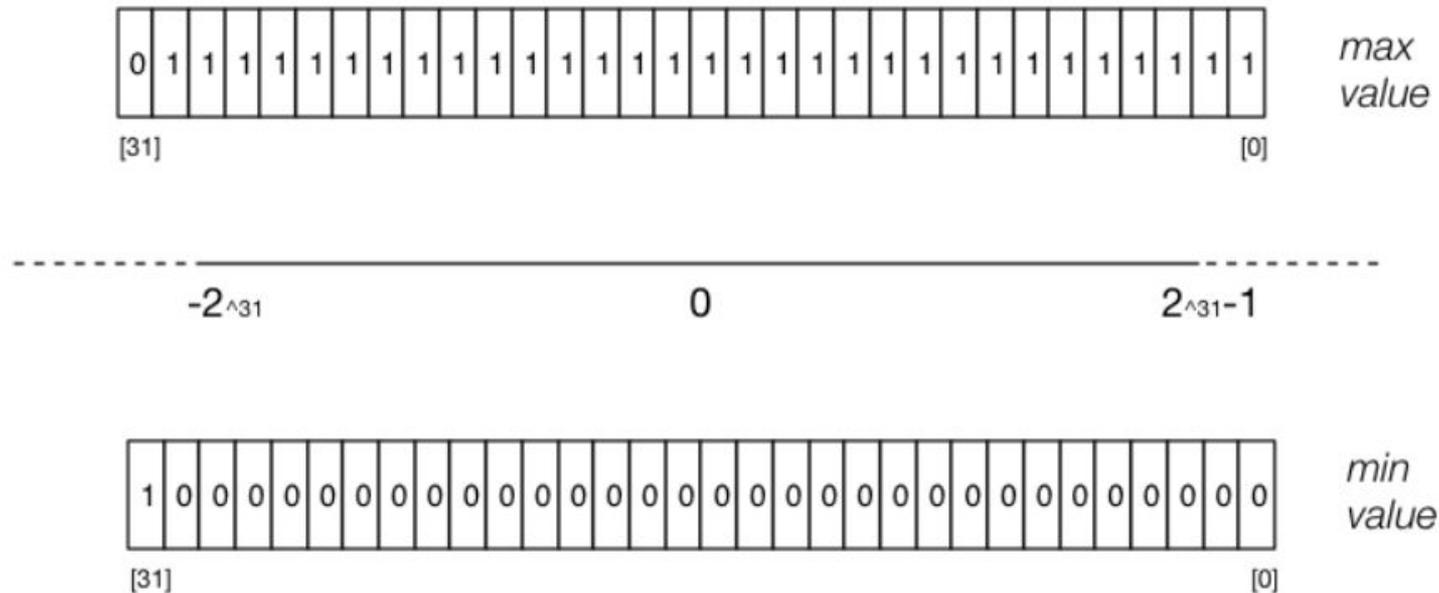
invert all the bits: 1111 1010

add 1: 1111 1011

- Repeat the process to go from -5 back to 5 again!

# Signed integers

- In C the `int` data type is 4 bytes on our system
  - we can store values from the range  $-2^{31}.. 2^{31} - 1$



# Example: 2's Complement Example

- Some simple code to examine 8-bit 2's complement numbers:

```
for (int i = -128; i < 128; i++) {  
    printf("%4d ", i);  
    print_bits(i, 8);  
    printf("\n");  
}
```

- `gcc 8_bit_twos_complement.c print_bits.c -o 8_bit_twos_complement`

# Example: Printing all 8-bit 2's complement

```
$ ./8_bit_twos_complement
-128 10000000
-127 10000001
-126 10000010
...
-3 11111101
-2 11111110
-1 11111111
0 00000000
1 00000001
2 00000010
3 00000011
...
125 01111101
126 01111110
127 01111111
```

# Exercise 8-bit numbers

What do the following bit patterns represent:

11111111

10000000

Assume they are unsigned 8 bit integers:

Now assume they are signed 8 bit integers:

# Exercise 8-bit numbers

What do the following bit patterns represent:

11111111

10000000

Assume they are unsigned 8 bit integers:

255

128

Now assume they are signed 8 bit integers:

# Exercise 8-bit numbers

What do the following bit patterns represent:

11111111

10000000

Assume they are unsigned 8 bit integers:

255

128

Now assume they are signed 8 bit integers: (Note  $2^8 == 256$ )

-1           (255 - 256)

-128        (128 - 256)

# Example: print\_bits\_of\_int.c

```
$ ./print_bits_of_int
Enter an int: 0
00000000000000000000000000000000
$ ./print_bits_of_int
Enter an int: 1
000000000000000000000000000000001
$ ./print_bits_of_int
Enter an int: -1
11111111111111111111111111111111
$ ./print_bits_of_int
Enter an int: 2147483647
01111111111111111111111111111111
$ ./print_bits_of_int
Enter an int: -2147483648
10000000000000000000000000000000
$
```

# Bits and Bytes on cse Servers

- On CSE servers, C types have these sizes
  - **char** = 1 byte = 8 bits
    - 42 is 00101010
  - **short** = 2 bytes = 16 bits,
    - 42 is 0000000000101010
  - **int** = 4 bytes = 32 bits,
    - 42 is 00000000000000000000000000000000101010
  - **double** = 8 bytes = 64 bits,
    - 42 = ?
- above are common sizes but not universal
- sizeof (int) might be 2 (bytes) on a small embedded CPU

# integer\_types.c - exploring integer types

	Type	Bytes	Bits
	char	1	8
	signed char	1	8
	unsigned char	1	8
	short	2	16
	unsigned short	2	16
	int	4	32
	unsigned int	4	32
	long	8	64
	unsigned long	8	64
	long long	8	64
	unsigned long long	8	64

# Exploring integer types and <limits.h>

Type	Min	Max
char	-128	127
signed char	-128	127
unsigned char	0	255
short	-32768	32767
unsigned short	0	65535
int	-2147483648	2147483647
unsigned int	0	4294967295
long	-9223372036854775808	9223372036854775807
unsigned long	0	18446744073709551615
long long	-9223372036854775808	9223372036854775807
unsigned long long	0	18446744073709551615

# stdint.h - guaranteed size integer types

- `#include <stdint.h>` to get below int types (and more) with known sizes
- We use these a lot in COMP1521!

```
        // range of values for type
        //          minimum          maximum
int8_t   i1; //          -128          127
uint8_t  i2; //           0          255
int16_t  i3; //        -32768        32767
uint16_t i4; //           0        65535
int32_t  i5; //    -2147483648    2147483647
uint32_t i6; //           0    4294967295
int64_t  i7; // -9223372036854775808  9223372036854775807
uint64_t i8; //           0 18446744073709551615
```

# Code Examples

overflow\_int.c

wrap\_around\_uint.c

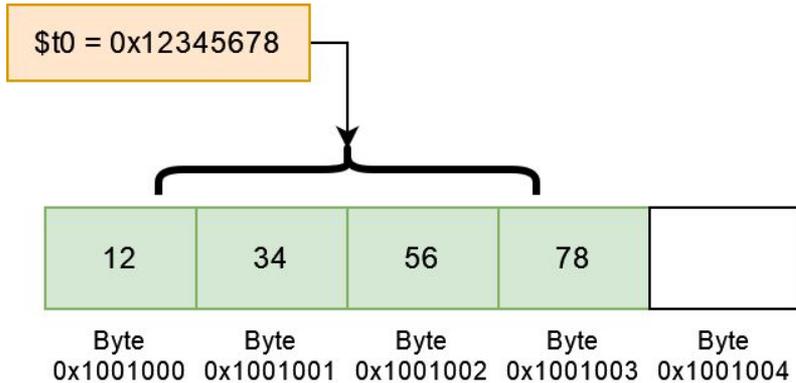
char\_bug.c

unsigned\_arithmetic.s

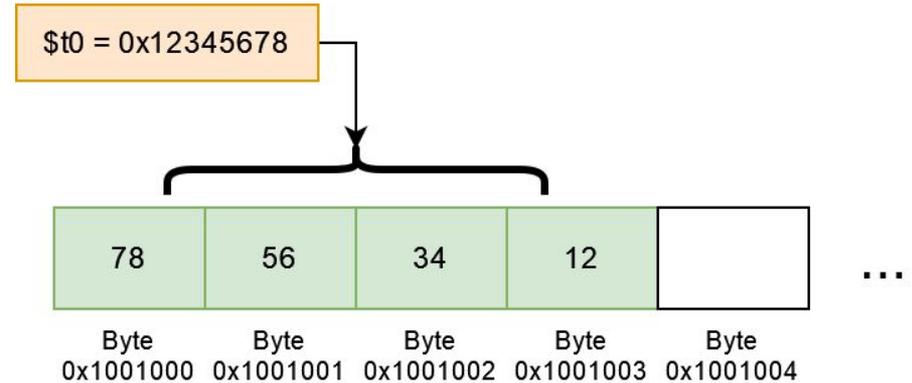
# New concept: Endian-ness

- “What order to put things in” is a hard question to answer
- Two schools of thought:
  - **Big-endian**: MSB at the “low address” - big bits “first!”
  - **Little-endian**: LSB at the “low address” - little bits “first!”

**BIG:**



**LITTLE:**



# Code example

- Mipsy-web is **little-endian**

```
.text
```

```
main:
```

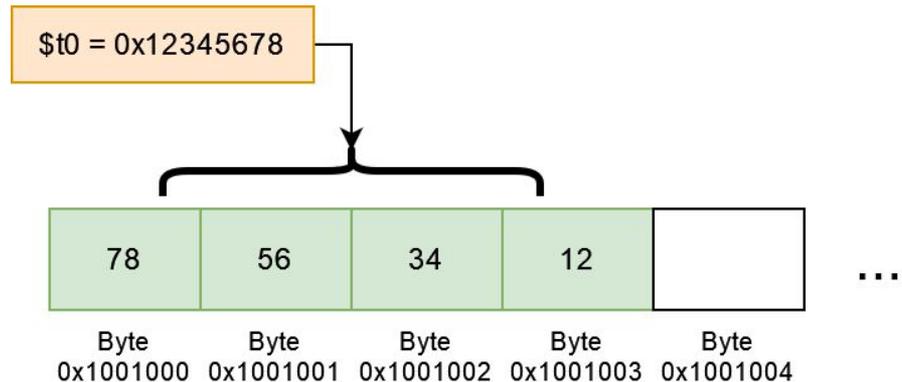
```
li $t0, 0x12345678
```

```
sw $t0, my_word
```

```
.data
```

```
my_word:
```

```
.space 4
```



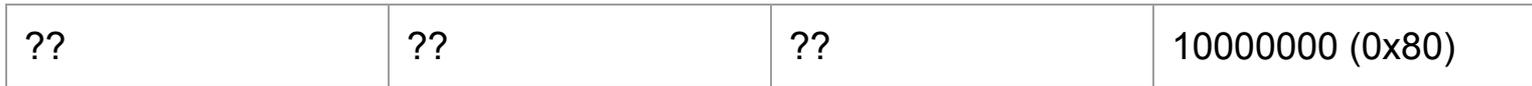
# Loading bytes, half-words

The results of these will depend on endianness:

- `lh/lb` assume the loaded byte/halfword is signed
- `lhu/lbu` for doing the same thing, but unsigned
  
- What might we need to do differently if we are loading a signed vs unsigned number?
- Why did we not need 2 versions for `lw`?

# Loading bytes, half-words

- Consider storing 0x80 in a 4 byte (32 bit register) register
- This only takes up 1 byte (8 bits)
- What do we store in the rest of the register?



# Loading bytes, half-words

- Consider storing 0x80 in a 4 byte (32 bit register) register
- This only takes up 1 byte (8 bits)
- What do we store in the rest of the register?

00000000 (0x00)	00000000 (0x00)	00000000 (0x00)	10000000 (0x80)
-----------------	-----------------	-----------------	-----------------

If we set them all the 0s, what happens if 0x80 is meant to be treated as a signed value?

# Loading bytes, half-words

- Consider storing 0x80 in a 4 byte (32 bit register) register
- This only takes up 1 byte (8 bits)
- What do we store in the rest of the register?

11111111 (0xFF)	11111111 (0xFF)	11111111 (0xFF)	10000000 (0x80)
-----------------	-----------------	-----------------	-----------------

If our value is meant to be signed we extend the sign bit. For a negative value this means extending with 1

# Loading Examples: lb

```
.text
```

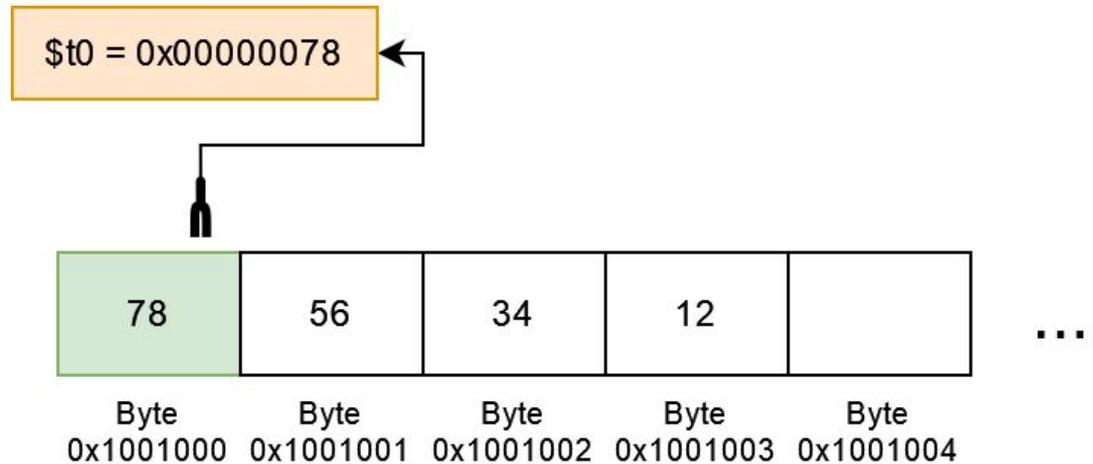
```
main:
```

```
    lb $t0, my_label
```

```
.data
```

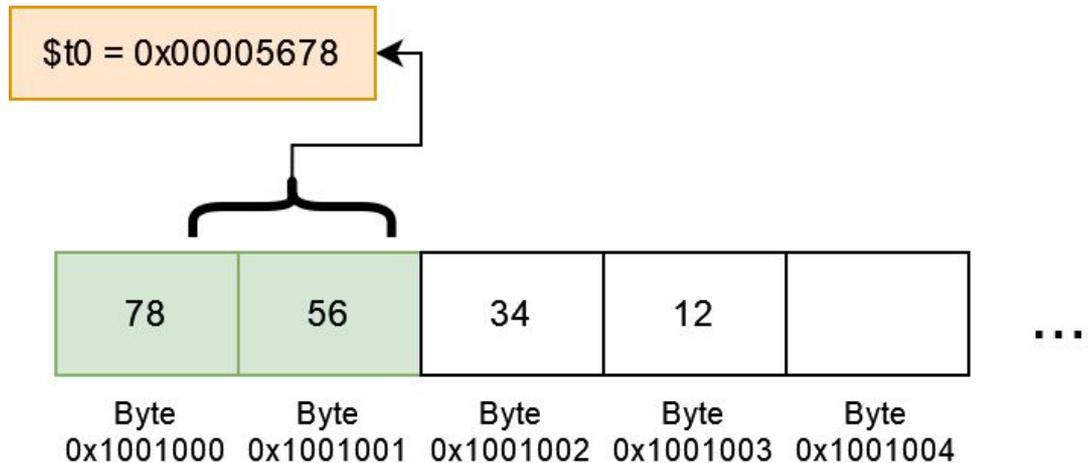
```
my_label:
```

```
    .word 0x12345678
```



# Loading Examples: lh

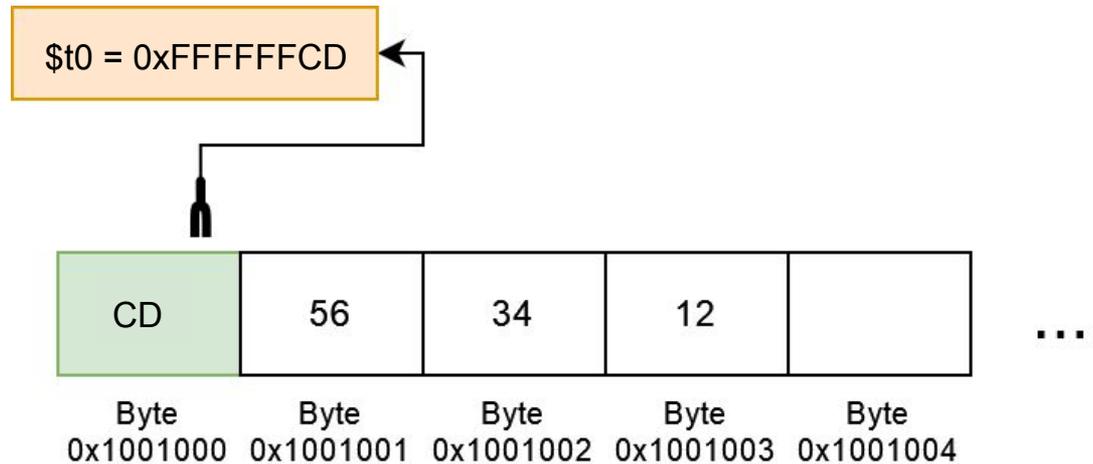
```
.text  
main:  
    lh $t0, my_label  
.data  
my_label:  
    .word 0x12345678
```



# Loading Examples Negative: lb

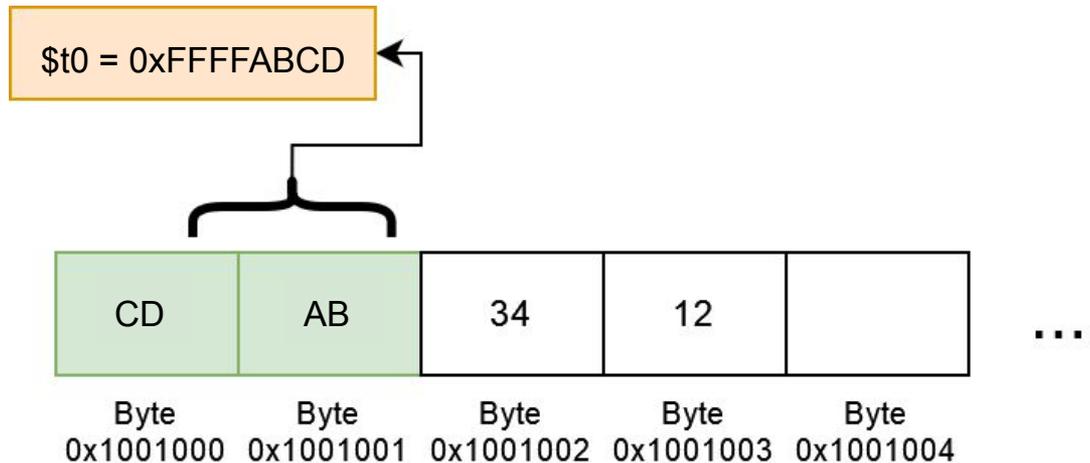
```
.text
main:
    lb $t0, my_label

.data
my_label:
    .word 0x1234ABCD
```



# Loading Examples Negative: lh

```
.text  
main:  
    lh $t0, my_label  
.data  
my_label:  
    .word 0x1234ABCD
```



# Loading Examples: lbu

```
.text
```

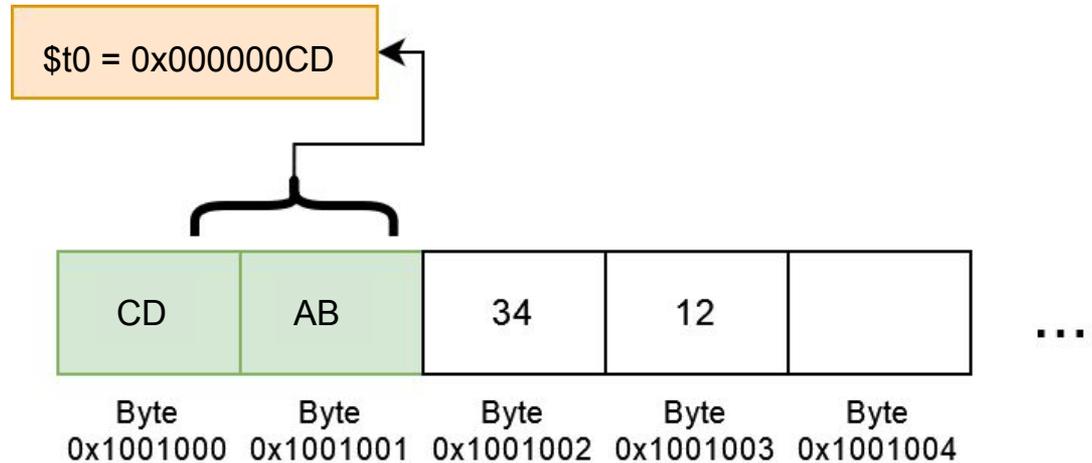
```
main:
```

```
    lbu $t0, my_label
```

```
.data
```

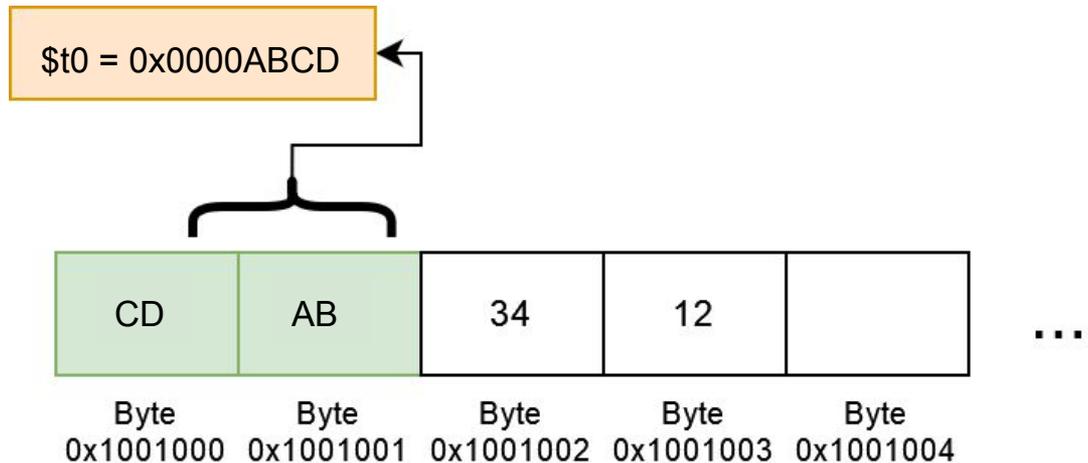
```
my_label:
```

```
    .word 0x1234ABCD
```



# Loading Examples Negative: lhu

```
.text  
main:  
    lhu $t0, my_label  
  
.data  
my_label:  
    .word 0x1234ABCD
```



# Endianness in C

endianness.c

# What did we learn today?

- Invalid C
- Integers
  - Endianness
  - Unsigned instructions in MIPS
- Next lecture:
  - Bitwise operators

# Feedback Please!

Your feedback is valuable!

If you have any feedback from today's lecture, please follow the link below or use the QR Code.

Please remember to keep your feedback constructive, so I can action it and improve your learning experience.



<https://forms.office.com/r/GXGX6vKDG6>

# Reach Out

Content Related Questions:

[Forum](#)

Admin related Questions email:

[cs1521@cse.unsw.edu.au](mailto:cs1521@cse.unsw.edu.au)



# Student Support | I Need Help With...

## My Feelings and Mental Health

Managing Low Mood, Unusual Feelings & Depression



**Mental Health Connect**

[student.unsw.edu.au/counselling](https://student.unsw.edu.au/counselling)  
Telehealth



**In Australia Call Afterhours  
UNSW Mental Health Support  
Line**

1300 787 026  
5pm-9am



**Mind HUB**

[student.unsw.edu.au/mind-hub](https://student.unsw.edu.au/mind-hub)  
Online Self-Help Resources



**Outside Australia  
Afterhours 24-hour  
Medibank Hotline**

+61 (2) 8905 0307

## Uni and Life Pressures

Stress, Financial, Visas, Accommodation & More



**Student Support  
Indigenous Student  
Support**

[student.unsw.edu.au/advisors](https://student.unsw.edu.au/advisors)

## Reporting Sexual Assault/Harassment



**Equity Diversity and Inclusion  
(EDI)**

[edi.unsw.edu.au/sexual-misconduct](https://edi.unsw.edu.au/sexual-misconduct)

## Educational Adjustments

To Manage my Studies and Disability / Health Condition



**Equitable Learning Service  
(ELS)**

[student.unsw.edu.au/els](https://student.unsw.edu.au/els)

## Academic and Study Skills



**Academic Language  
Skills**

[student.unsw.edu.au/skills](https://student.unsw.edu.au/skills)

## Special Consideration

Because Life Impacts our Studies and Exams



**Special Consideration**

[student.unsw.edu.au/special-consideration](https://student.unsw.edu.au/special-consideration)