

COMP1521 25T2 — Floating-Point Numbers

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Floating Point Numbers

- C has three floating point types
 - **float** ... typically 32-bit (lower precision, narrower range)
 - **double** ... typically 64-bit (higher precision, wider range)
 - **long double** ... typically 128-bits (but maybe only 80 bits used)
- Floating point constants, e.g : **3.14159 1.0e-9** are **double**
- Reminder: division of 2 ints in C yields an int.
 - but division of double and int in C yields a double.

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Floating Point Number - Output

```
double d = 4/7.0;
// prints in decimal with (default) 6 decimal places
printf("%lf\n", d);          // prints 0.571429
// prints in scientific notation
printf("%le\n", d);          // prints 5.714286e-01
// picks best of decimal and scientific notation
printf("%lg\n", d);          // prints 0.571429
// prints in decimal with 9 decimal places
printf("%.9lf\n", d);        // prints 0.571428571
// prints in decimal with 1 decimal place and field width of 5
printf("%10.1lf\n", d);      // prints      0.6
```

source code for float_output.c

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The decimal fraction 0.75 means

- $7 \cdot 10^{-1} + 5 \cdot 10^{-2} = 0.7 + 0.05 = 0.75$
- or equivalently $75/10^2 = 75/100 = 0.75$

Similary 0.11_2 means

- $1 \cdot 2^{-1} + 1 \cdot 2^{-2} = 0.5 + 0.25 = 0.75$
- or equivalently $3/2^2 = 3/4 = 0.75$

Similarly $0.C_{16}$ means

- $12 \cdot 16^{-1} = 0.75$
- or equivalently $12/16^1 = 3/4 = 0.75$

Note: We call the \cdot a radix point rather than a decimal point when we are dealing with other bases.

Fractions in different Bases

The algorithm to convert a decimal fraction to another base is:

- take the fractional component and multiply by the base
- the whole number becomes the next digit to the right of the radix point in our fraction.
- repeat this process until the fractional part becomes exhausted or we have sufficient digits
- this process is not guaranteed to terminate.

Converting Decimal Fractions to Binary

For example if we want to convert 0.3125 to base 2

- $0.3125 \cdot 2 = 0.625$
- $0.625 \cdot 2 = 1.25$
- $0.25 \cdot 2 = 0.5$
- $0.5 \cdot 2 = 1.0$

Therefore $0.3125 = 0b0.0101$

Convert the following decimal values into binary

- 12.625
- 0.1

Floating Point Numbers

- can have fractional numbers in other bases, e.g.: $110.101_2 == 6.625_{10}$
- if we represent floating point numbers with a fixed small number of bits
 - there are only a finite number of bit patterns
 - can only represent a finite subset of reals
- almost all real values will have no exact representation
- value of arithmetic operations may be real with no exact representation
- we must use closest value which can be exactly represented
- this approximation introduces an error into our calculations
- often, does not matter
- sometimes ... can be disastrous

Fixed-Point Representation

- fixed-point is a simple trick to represent fractional numbers as integers
 - every value is multiplied by a particular constant, e.g. 1000 and stored as integer
 - so if constant is 1000, could represent 56.125 as an integer (56125)
 - but not 3.141592
- usable for some problems, but not ideal
- used on small embedded processors without silicon floating point
- major limitation is only small range of values can be represented
 - for example with 32 bits, and using 65536 (2^{16}) as constant
 - 16 bits used for integer part
 - 16 bits used for the fraction
 - minimum $2^{-16} \approx 0.000015$
 - maximum $2^{15} \approx 32768$

- you met scientific notation, e.g 6.0221515×10^{23} in physics or other science classes
- we can represent numbers on a computer in a similar way to scientific notation
- but using binary instead of base ten, e.g 10.6875

$$= 1010.1011 = 1.0101011 \times 2^{11_2} = (1 + 43/128) \times 2^3 = 1.3359375 \times 8 = 10.6875$$
- allows a much bigger range of values to be represented than fixed point
- using only 8 bits for the exponent, we can represent numbers from $10^{-38} \dots 10^{+38}$
- using only 11 bits for the exponent, we can represent numbers from $10^{-308} \dots 10^{+308}$
- leads to numbers close to zero having higher precision (more accurate) which is good

choosing which exponential representation

- exponent notation allows multiple representations for a single value
 - e.g $1.0101011 \times 2^{11_2} = 10.6875$ and $10.101011 \times 2^{10_2} = 10.6875$
- having multiple representations would make implementing arithmetic slower on CPU
- better to have only one representation (one bit pattern) representing a value
- decision - use representation with exactly one digit in front of decimal point
 - use $1.0101011 \times 2^{11_2}$ not $10.101011 \times 2^{10_2}$ or 1010.1011×2^{0_2}
 - this is called normalization
- weird hack: as we are using binary the first digit must be a one we don't need to store it
 - as we long we have a separate representation for zero

floating_types.c - print characteristics of floating point types

```
float f;
double d;
long double l;
printf("float      %2lu bytes  min=%-12g  max=%g\n", sizeof f, FLT_MIN, FLT_MAX);
printf("double     %2lu bytes  min=%-12g  max=%g\n", sizeof d, DBL_MIN, DBL_MAX);
printf("long double %2lu bytes  min=%-12Lg  max=%Lg\n", sizeof l, LDBL_MIN, LDBL_MAX);
```

source code for floating_types.c

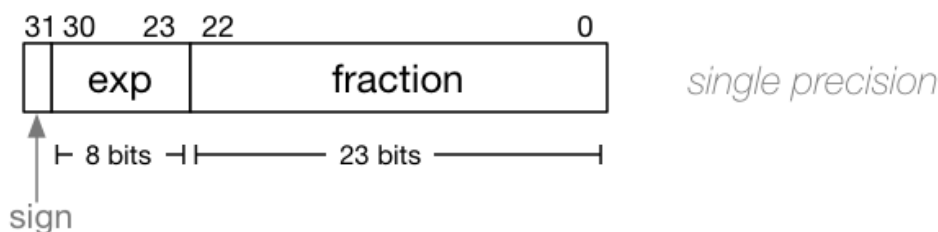
```
$ ./floating_types
float      4 bytes  min=1.17549e-38  max=3.40282e+38
double     8 bytes  min=2.22507e-308  max=1.79769e+308
long double 16 bytes min=3.3621e-4932  max=1.18973e+4932
```

- 1970s Intel building microprocessors (single-chip CPUs)
- 1976 Intel developing coprocessor (separate chip) for floating-point arithmetic
- Intel asked William Kahan, University of California to design format
- other manufacturers didn't want to be left out
- IEEE 754 standard working group formed
- Kahan and others produced well-designed robust specification
- accepted by manufacturers who begin using it for new architectures
- IEEE 754 standard released in 1985 (update to standard in 2008)
- today, almost all computers use IEEE 754

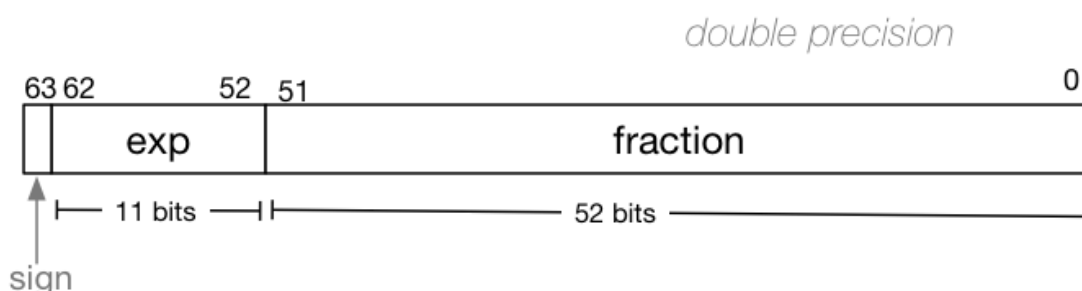
IEEE 754 standard

- C floats almost always IEEE 754 single precision (binary32)
- C double almost always IEEE 754 double precision (binary64)
- C long double might be IEEE 754 (binary128)
- IEEE 754 representation has 3 parts: *sign*, *fraction* and *exponent*
- numbers have form $sign \text{ fraction} * 2^{exponent}$, where *sign* is +/-
- *fraction* always has 1 digit before decimal point (*normalized*)
- *exponent* is stored as positive number by adding constant value (*bias*)

Internal structure of floating point values



single precision



double precision

Example of normalising the fraction part in binary:

- 1010.1011 is normalized as $1.0101011 * 2^{011}$
- $1010.1011 = 10 + 11/16 = 10.6875$
- $1.0101011 * 2^{011} = (1 + 43/128) * 2^3 = 1.3359375 * 8 = 10.6875$

The normalised fraction part always has 1 before the decimal point.

Example of determining the exponent in binary:

- if exponent is 8-bits, then the bias = $2^{8-1} - 1 = 127$
- valid bit patterns for exponent are **00000001 .. 11111110**
- these correspond to exponent values of -126 .. 127

Floating Point Numbers

Example (single-precision):

$$150.75 = 10010110.11$$

// normalise fraction, compute exponent

$$= 1.001011011 \times 2^7$$

// sign bit = 0

// exponent = 10000110

// fraction = 001011011000000000000000

$$= 01000011000101101100000000000000$$

Distribution of Floating Point Numbers

- floating point numbers not evenly distributed
- representations get further apart as values get bigger
 - this works well for most calculations
 - but can cause weird bugs
- double (IEEE 754 64 bit) has 52-bit fractions so:
 - between 2^n and 2^{n+1} there are 2^{52} doubles evenly spaced
 - e.g. in the interval 2^{-42} and 2^{-43} there are 2^{52} doubles
 - and in the interval between 1 and 2 there are 2^{52} doubles
 - and in the interval between 2^{42} and 2^{43} there are 2^{52}
 - near 0.001 - doubles are about 0.000000000000000002 apart
 - near 1000 - doubles are about 0.00000000000002 apart
 - near 1000000000000000 - doubles are about 0.25 apart
 - above 2^{53} - doubles are more than 1 apart

IEEE-754 Single Precision example: 0.15625

0.15625 is represented in IEEE-754 single-precision by these bits:

```
00111111000100000000000000000000
```

sign | exponent | fraction

```
0 | 01111100 | 0100000000000000000000
```

sign bit = 0

sign = +

```
raw exponent    = 01111100 binary
                 = 124 decimal
```

```
actual exponent = 124 - exponent_bias
                 = 124 - 127
                 = -3
```

```
number = +1.010000000000000000000000 binary * 2**-3
        = 1.25 decimal * 2**-3
        = 1.25 * 0.125
        = 0.15625
```

source code for explain_float_representation.c

IEEE-754 Single Precision example: -0.125

```
$ ./explain_float_representation -0.125
```

-0.125 is represented as a float (IEEE-754 single-precision) by these bits:

```
10111111000000000000000000000000
```

sign | exponent | fraction

```
1 | 01111100 | 0000000000000000000000
```

sign bit = 1

sign = -

```
raw exponent    = 01111100 binary
                 = 124 decimal
```

```
actual exponent = 124 - exponent_bias
                 = 124 - 127
                 = -3
```

```
number = -1.000000000000000000000000 binary * 2**-3
        = -1 decimal * 2**-3
        = -1 * 0.125
        = -0.125
```

IEEE-754 Single Precision example: 150.75

```
$ ./explain_float_representation 150.75
```

150.75 is represented in IEEE-754 single-precision by these bits:

```
01000011000101101100000000000000
```

sign | exponent | fraction

```
0 | 10000110 | 0010110110000000000000
```

sign bit = 0

sign = +

```
raw exponent    = 10000110 binary
                 = 134 decimal
```

```
actual exponent = 134 - exponent_bias
                 = 134 - 127
                 = 7
```

```
number = +1.001011011000000000000000 binary * 2**7
        = 1.17773 decimal * 2**7
        = 1.17773 * 128
        = 150.75
```

```
$ ./explain_float_representation -96.125
-96.125 is represented in IEEE-754 single-precision by these bits:
11000010111000000010000000000000
sign | exponent | fraction
  1  | 10000101 | 100000001000000000000000
sign bit = 1
sign = -
raw exponent    = 10000101 binary
                  = 133 decimal
actual exponent = 133 - exponent_bias
                  = 133 - 127
                  = 6
number = -1.100000001000000000000000 binary * 2**6
        = -1.50195 decimal * 2**6
        = -1.50195 * 64
        = -96.125
```

IEEE-754 Single Precision exploring bit patterns #1

```
$ ./explain_float_representation 00111101110011001100110011001101
sign bit = 0
sign = +
raw exponent    = 01111011 binary
                  = 123 decimal
actual exponent = 123 - exponent_bias
                  = 123 - 127
                  = -4
number = +1.10011001100110011001101 binary * 2**-4
        = 1.6 decimal * 2**-4
        = 1.6 * 0.0625
        = 0.1
```

infinity.c: exploring infinity

- IEEE 754 has a representation for +/- infinity
- propagates sensibly through calculations

```
double x = 1.0/0.0;
printf("%lf\n", x); //prints inf
printf("%lf\n", -x); //prints -inf
printf("%lf\n", x - 1); // prints inf
printf("%lf\n", 2 * atan(x)); // prints 3.141593
printf("%d\n", 42 < x); // prints 1 (true)
printf("%d\n", x == INFINITY); // prints 1 (true)
```

source code for infinity.c

- C (IEEE-754) has a representation for invalid results:
 - NaN (not a number)
- ensures errors propagates sensibly through calculations

```
double x = 0.0/0.0;
printf("%lf\n", x); //prints nan
printf("%lf\n", x - 1); // prints nan
printf("%d\n", x == x); // prints 0 (false)
printf("%d\n", isnan(x)); // prints 1 (true)
```

source code for nan.c

IEEE-754 Single Precision example: inf

```
$ ./explain_float_representation inf
inf is represented in IEEE-754 single-precision by these bits:
01111111110000000000000000000000
sign | exponent | fraction
  0  | 11111111 | 000000000000000000000000
sign bit = 0
sign = +
raw exponent    = 11111111 binary
                  = 255 decimal
number = +inf
```

IEEE-754 Single Precision exploring bit patterns #2

```
$ ./explain_float_representation 01111111110000000000000000000000
sign bit = 0
sign = +
raw exponent    = 11111111 binary
                  = 255 decimal
number = NaN
```

source code for explain_float_representation.c

```
double a, b;
a = 0.1;
b = 1 - (a + a + a + a + a + a + a + a + a + a);
if (b != 0) { // better would be fabs(b) > 0.000001
    printf("1 != 0.1+0.1+0.1+0.1+0.1+0.1+0.1+0.1+0.1+0.1\n");
}
printf("b = %g\n", b); // prints 1.11022e-16
```

source code for double_imprecision.c

- do not use == and != with floating point values
- instead check if values are close

Consequences of most reals not having exact representations

- The approximate representation of reals can produce unexpected errors.
- If we subtract or divide two values which are very close together a large relative error can result.
- This is called cancellation or catastrophic cancellation.
- For example, if x is close to 0, $\cos(x)$ is close to 1
 - calculating $1 - \cos(x)$ can produce a large error in a calculation
- we can avoid the error by replacing $1 - \cos(x)$ with $2 * \sin(x/2) * \sin(x/2)$

Consequences of most reals not having exact representations

```
double x;
printf("Enter x: ");
scanf("%lf", &x);
printf("(1 - cos(x)) / (x * x) = %lf\n", (1 - cos(x)) / (x * x));
printf("(2 * sin(x/2) * sin(x/2)) / (x * x) = %lf\n", (2 * sin(x/2) * sin(x/2)) / (x * x));
```

source code for cancelled.c

```
$ ./a.out
Enter x: 0.123
(1 - cos(x)) / (x * x) = 0.499370
(2 * sin(x/2) * sin(x/2)) / (x * x) = 0.499370
$ ./a.out
Enter x: 0.000000011
(1 - cos(x)) / (x * x) = 0.917540
(2 * sin(x/2) * sin(x/2)) / (x * x) = 0.500000
```

```

if (d == d) {
    printf("d == d is true\n");
} else {
    // will be executed if d is a NaN
    printf("d == d is not true\n");
}
if (d == d + 1) {
    // may be executed if d is large
    // because closest possible representation for d + 1
    // is also closest possible representation for d
    printf("d == d + 1 is true\n");
} else {
    printf("d == d + 1 is false\n");
}

```

source code for double_not_always.c

Another reason not to use == with floating point values

```

$ dcc double_not_always.c -o double_not_always
$ ./double_not_always 42.3
d = 42.3
d == d is true
d == d + 1 is false
$ ./double_not_always 4200000000000000000
d = 4.2e+18
d == d is true
d == d + 1 is true
$ ./double_not_always NaN
d = nan
d == d is not true
d == d + 1 is false

```

because closest possible representation for $d + 1$ is also closest possible representation for d

source code for double_not_always.c

Consequences of most reals not having exact representations

```

// loop looks to print 10 numbers but actually never terminates
double d = 9007199254740990;
while (d < 9007199254741000) {
    printf("%lf\n", d); // always prints 9007199254740992.000000
    // 9007199254740993 can not be represented as a double
    // closest double is 9007199254740992.0
    // so 9007199254740992.0 + 1 = 9007199254740992.0
    d = d + 1;
}

```

source code for double_disaster.c

- 9007199254740993 is $2^{53} + 1$
it is smallest integer which can not be represented exactly as a double
- The closest double to 9007199254740993 is 9007199254740992.0
- aside: 9007199254740993 can not be represented by a `int32_t`
it can be represented by `int64_t`

Convert the following floating point numbers to decimal.

Assume that they are in IEEE 754 single-precision format.

0 10000000 110000000000000000000000

1 01111110 100000000000000000000000