COMP1521 25T1

Week 1 Lecture 2

MIPS: Basics and Control

Adapted from slides by Abiram Nadarajah, Hammond Pearce, Andrew Taylor and John Shepherd's slides

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C Revision and recursion Lab Week 2

- Monday 11-1pm lute
- Tuesday 7-9pm online on BlackBoard Collaborate
- Bookings required: <u>Book ticket</u>
- Code: (you may be asked to enter a code when booking your ticket) COMP1521

Content:

- C revision and recursion lab style based questions
- Can also get help with regular lab week 1 if struggling

Today's Lecture

- Recap Lecture 1
- System Calls
- Style
- Simplified C
 - \circ and goto
- MIPS Control
 - if statements
 - boolean expressions
 - while loops/for loops



DISCLAIMER:

Code written in lectures may not necessarily have the best style!

- Live lecture code is meant to be quick and dirty, to demonstrate a concept
- Will quickly overview good style soon, but refer to your tutor, tut solutions, lab solutions and assignment resources.

Recap of Lecture 1 Intro to MIPS

- Exploring different types of storage/memory
- RAM contains everything a program needs in a given moment
- Instructions!
- Assembly language (MIPS)!
- Registers!
- We did not get up to System calls!

Recap exercise

- Open Mipsy
- Store the value 1 in \$t0
- Store the value 4 in \$t1
- Sum these and put the result in **\$t2**

Assembly Syntax overview

- Assembly instructions, each on their own line
- Generally a 1:1 mapping from assembly instructions to binary instructions
- However, assemblers also provide pseudo-instructions for convenience
- pseudo-instructions turn into 1-3 real CPU instructions
 - Example:
 - li \$t0, 5 gets mapped to real equivalent CPU instruction
 - addi \$t0, \$zero, 5
 - You will see many more as you write more code in MIPS.

More about registers

Registers have symbolic names and also numeric names \$t0 is also known as \$8

There are many registers we won't learn about or use till week 3.

Number	Names	Conventional Usage	
0	zero	Constant 0	
1	at	Reserved for assembler	
2,3	v0,v1	Expression evaluation and results of a function	
47	a0a3	Arguments 1-4	
816	t0t7	Temporary (not preserved across function calls)	
1623	s0s7	Saved temporary (preserved across function calls)	
24,25	t8,t9	Temporary (not preserved across function calls)	
26,27	k0,k1	Reserved for Kernel use	
28	gp	Global Pointer	
29	sp	Stack Pointer	
30	fp	Frame Pointer	
31	ra	Return Address (used by function call instructions)	

What do MIPS instructions look like?

- 32 bits long
- Specify:
 - An operation
 - (The thing to do)
 - 0 or more operands
 - (The thing to do it over)
- For example:

OPCODE	R1	R2	R3	R4	OPCODE	R-type
6 bits	►5 bits	►5 bits	►5 bits	5 bits	-6 bits-	

I-type

001000010000100100000000000001100 addi \$t1, \$t0, 12

Aside: Hexadecimal

0x in C and mipsy means hexadecimal.

Hexadecimal uses 16 digits. It uses 0-9 then A-F

We will learn more about this later in the course.

Decimal	Hexadecimal	Decimal	Hexadecimal
0	0	10	A
1	1	11	В
2	2	12	С
3	3	13	D
4	4	14	E
5	5	15	F
6	6	16	10
7	7	17	11
8	8	18	12
9	9	19	13

Aside: Hexadecimal

- We often use Hexadecimal to represent addresses and other binary data like instructions.
 - Easier for humans to read than binary as it is compact
 - Maps more easily to binary than decimal
- 8 hex digits can represent 32 bits
 - For example the instruction addi \$t4, \$zero, 7 maps to binary instruction

0010000000010110000000000000111 Which can be represented in hexadecimal by 0x200b0007

MIPS and mipsy documentation

Literally your best friend (it'll even be there for you in the exam $\overline{\mathfrak{S}}$)

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MIPS Instruction Set

An overview of the instruction set of the MIPS32 architecture as implemented by the mipsy and SPIM emulators. Adapted from reference documents from the University of Stuttgart and Drexel University, from material in the appendix of Patterson and Hennessey's *Computer Organization and Design*, and from the MIPS32 (r5.04) Instruction Set reference.

- Registers
- Memory
- Syntax
- Instructions
 - CPU Arithmetic Instructions
 - CPU Logical Instructions
 - CPU Shift Instructions

But how can we do input and output?

System calls

- None of the instructions we have access to can interact with the outside world (eg. printing, scanning)
- Instead, we request the operating system to perform these tasks for us this process is called a **system call**
- The operating system can access privileged instructions on the CPU (eg. communicating to other devices)
- *mipsy* simulates a very basic operating system
- Will explore real system calls in the second half of the course

Common mipsy syscalls

Service	\$v0	Arguments	Returns	
<pre>printf("%d")</pre>	1	int in \$a0		
fputs	4	string in \$a0		
<pre>scanf("%d")</pre>	5	none int in \$v0		
fgets	8	line in \$a0, length in \$a1		
exit(0)	10	none		
<pre>printf("%c")</pre>	11	char in \$a0		
<pre>scanf("%c")</pre>	12	none char in \$v0		

Service	\$v0	Arguments	Returns
printf("%f")	2	float in \$f12	
<pre>printf("%lf")</pre>	3	double in \$f12	
<pre>scanf("%f")</pre>	6	none	float in \$f0
<pre>scanf("%lf")</pre>	7	none	double in \$f0
sbrk(nbytes)	9	nbytes in \$a0	address in \$v0
<pre>open(filename, flags, mode)</pre>	13	filename in \$a0, flags in \$a1, mode \$a2	fd in \$v0
<pre>read(fd, buffer, length)</pre>	14	fd in \$a0, buffer in \$a1, length in \$a2	number of bytes read in \$∨0
write(fd, buffer, length)	15	fd in \$a0, buffer in \$a1, length in \$a2	number of written in \$v0
close(fd)	16	fd in \$a0	
exit(status)	17	status in \$a0	

Probably only used for challenge exercises in COMP1521

Let's try to print out the number 42

The system call workflow

- We specify which system call we want in v0
 - eg. print_int is syscall 1:
 - o li \$v0, 1
- We specify arguments (if any)

o li \$a0, 42

- We transfer execution to the operating system
 - The OS will fulfill our request if it looks sane
 - syscall
- Some syscalls may return a value check syscall table
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Let's try to print out the number 42 and then 99 on the next line

Let's add 2 numbers and print out the result

How do we print strings?

Printing Strings and the Data segment

- We need to define our string in the data section
- Then pass the address of our string to our system call in \$a0
- We need to use the .data directive so we can create global data in our program
- We need to use the .asciiz directive so we can define a string and give the string a label!
- We need to use the **la** to load the address of the string!!

Hello COMP1521 revisited

.text

main:

li \$v0, 4
la \$a0, hello_msg
syscall

syscall 4: print_string
#
printf("Hello COMP1521!!\n");

```
li $v0,0
jr $ra # return 0;
```

.data

hello_msg:

.asciiz "Hello COMP1521!!\n"

Example: Integer Average

```
// Translate into MIPS
int main(void) {
    int a, b;
    printf("Enter a number: ");
    scanf("%d", &a);
    printf("Enter another number: ");
    scanf("%d", &b);
   printf("The average is dn", (a + b)/2);
    return 0;
```

Simplified C

- Translating C code directly to MIPS is not fun
- Simplify your C code and then translate it to "simplified C":
 - Simplified C is generally written so that each line of C code maps to one MIPS instruction
 - Compile your simplified C and make sure it still works as expected
 - Translate each line of simplified C to MIPS

Putting data in registers

- li (load immediate) is loading a fixed value into a register
 li \$t0, 7
- la (load address) is for loading a fixed address into a register
 - remember, labels really just represent addresses!
 - o la \$t0, my_label
- move is for copying value from a register into another register
 - o move \$t0, \$t1

Assembly Language Syntax Recap

- Labels
 - Appended with :
 - They represent memory addresses
- Comments
 - Start with #
- Directives
 - Symbols beginning with . eg .asciiz .text .data
- Constant definitions
 - Like #define in C e.g.
 - MAX_NUMBERS = 10

After adding two numbers successfully in Assembly Language

We deserve a quick break now!

MIPS Control

So far

- Our programs have implemented fixed, predictable behaviour.
 - Execute linearly we always go down to the next instruction
- However, what if we want to implement **logic** in our code?
 - if statements conditional code execution
 - **for/while** loops repeat some instructions?

if/else and **loops don't exist** in MIPS - we have to use branching to implement these ourselves

Branch Instructions

- We have many conditional branch instructions of the form:
 - "if condition is true, jump to instruction at a given label" e.g.
 - o ble \$t0, \$t1, label1 # if (\$t0 <= \$t1)</pre>
 - o bgt \$t0, 5, label1 # if (\$t0 > 5)
- We have an unconditional branch instruction too
 b label1
- How do we implement this in our simplified C code?

Branch/jump instructions

b label	pc += I«2	pseudo-instruction
$\mathbf{beq} \ r_s$, r_t , label	if (r_s == r_t) pc += I«2	000100ssssstttttIIIIIIIIIIIIIII
bne r_s , r_t , label	if (r_s != r_t) pc += I«2	000101ssssstttttIIIIIIIIIIIIIII
ble r_s , r_t , label	if ($r_s \ll r_t$) pc += I«2	pseudo-instruction
${\tt bgt}r_s$, r_t , label	if (r_s > r_t) pc += I«2	pseudo-instruction
${\tt blt}r_s{\tt,r}_t{\tt,label}$	if (r_s < r_t) pc += I«2	pseudo-instruction
bge r_s , r_t , label	if ($r_s >= r_t$) pc += I«2	pseudo-instruction
blez r_s , label	if (r_s <= 0) pc += I«2	000110sssss00000IIIIIIIIIIIIIII
${f bgtz}r_s$, ${f label}$	if (r_s > 0) pc += I«2	000111sssss00000IIIIIIIIIIIIIII
${f bltz} r_s$, ${f label}$	if (r_s < 0) pc += I«2	000001sssss00000IIIIIIIIIIIIIII
${f bgez}r_s$, ${f label}$	if (r_s >= 0) pc += I«2	000001sssss00001IIIIIIIIIIIIIII
bnez r_s , label	if (r_s != 0) pc += I«2	pseudo-instruction
$\mathbf{beqz} \; r_s$, label	if (r_s == 0) pc += I«2	pseudo-instruction

- Allows you to transfer the flow of execution to a different instruction *conditionally*
 - except **b**, which is unconditional
- Can replace r_t with a constant in mipsy

COMP1511 staff hid this simple trick!

In C, **goto** allows jumping to any arbitrary label within a program.

This means we can effectively jump around within a program however we wish.

What will this code do?

```
int main(void) {
    goto sleep;
    printf("Please pay close attention\n");
sleep:
    printf("You are getting sleepy\n");
    goto sleep;
    printf("Please wake up now!");
    return 0;
```

With great power comes great responsibility

Edgar Dijkstra: Go To Statement Considered Harmful

Go To Statement Considered Harmful

Key Words and Phrases: go to statement, jump instruction, branch instruction, conditional clause, alternative clause, repetitive clause, program intelligibility, program sequencing *CR* Categories: 4.22, 5.23, 5.24

EDITOR:

For a number of years I have been familiar with the observation (hat the quality of programmers is a decreasing function of the density of go to statements in the programs they produce. More recently I discovered why the use of the go to statement has such disastrous effects, and I became convinced that the go to statement should be abolished from all "higher level" programming languages (i.e. everything except, perhaps, plain machine code). At that time I did not attach too much importance to this discovery; I now submit my considerations for publication because in very recent discussions in which the subject turned up, I have been urged to do so.

My first remark is that, although the programmer's activity ends when he has constructed a correct program, the process dynamic progress is only characterized when we also give to which call of the procedure we refer. With the inclusion of procedures we can characterize the progress of the process via a sequence of textual indices, the length of this sequence being equal to the dynamic depth of procedure calling.

Let us now consider repetition clauses (like, while B repeat A or repeat A until B). Logically speaking, such clauses are now superfluous, because we can express repetition with the aid of recursive procedures. For reasons of realism I don't wish to exclude them: on the one hand, repetition clauses can be implemented quite comfortably with present day finite equipment; on the other hand, the reasoning pattern known as "induction" makes us well equipped to retain our intellectual grasp on the processes generated by repetition clauses. With the inclusion of the repetition clauses textual indices are no longer sufficient to describe the dynamic progress of the process. With each entry into a repetition clause, however, we can associate a so-called "dynamic index," inexorably counting the ordinal number of thes procedure calls) may be applied nestedly, we find that now the

Go To Considered Harmful (1968)

Don't (ab)use goto

Don't use it in your actual C programs.

- goto makes programs more difficult to read
- **goto** makes it hard for compilers to optimise code, resulting in slower programs
- In general, do not use **goto** without good reason!
 - Typically only kernel/embedded programmers use goto
- We will use it in this course ONLY for writing simplified C to translate into MIPS.

Simplifying if-else statements

```
int main(void) {
    int n;
    printf("Enter a number: ");
    scanf("%d", &n);

    if (n % 2 == 0) {
        printf("even\n");
    }
    return 0;
}
```

```
int main(void) {
    int n;
    printf("Enter a number: ");
    scanf("%d", &n);
    int tmp = n % 2;
    if (tmp != 0) goto if_even_end;
        printf("even\n");
if_even_end:
    return 0;
}
```

Now we can write it in MIPS.

Exercise: add an else statement for odd numbers

Style

- Have equivalent C code as *inline comments*
- Huge recommendation: indent with 8-wide tabs
- We generally don't indent to show structure
 - i.e no indenting within loops or if statements, etc.
- Instead:
 - don't indent labels
 - indent instructions by one step
- For this course: focus on readable code, not reducing number of registers used or lines of code

More complex conditionals:

Split combined "or" conditions

```
if (milk_age > 48 ||
	milk_level < 10) {
	printf("Replace milk\n");
} else {
	printf("Milk okay!\n");
}
printf("Done!\n");
```

More complex conditionals:

Split combined "or" conditions

```
if (milk_age > 48 ||
	milk_level < 10) {
	printf("Replace milk\n");
} else {
	printf("Milk okay!\n");
} mi
printf("Done!\n");
```

```
if (milk_age > 48) goto milk_replace;
if (milk_level < 10) goto milk_replace;</pre>
```

```
printf("Milk okay!\n");
goto milk_replace__end;
```

```
milk_replace:
    printf("Replace milk\n");
```

```
milk_replace__end:
    printf("Done!");
```

More complex conditionals: &&

Invert the condition to use || (De Morgan's Law)

```
if (x >= 0 && x <= 100) {
    // in bounds
} else {
    // out of bounds
}
return 0;</pre>
```

More complex conditionals: &&

Invert the condition to use || (De Morgan's Law)

```
if (x >= 0 && x <= 100) {
    // in bounds
} else {
    // out of bounds
}
return 0;
    return 0;
    if (x < 0 || x > 100) {
         // out of bounds
         // out of bounds
         // in bounds
         return 0;
         return 0;
         return 0;
         // in bounds
         // in bounds
```

More complex conditionals:

Split into separate conditionals:

```
if (x < 0 || x > 100) {
    // out of bounds
} else {
    // in bounds
}
return 0;
```

More complex conditionals:

Split into separate conditionals:

```
if (x < 0) goto x_out_of_bounds;
                               if (x > 100) goto x_out_of_bounds;
if (x < 0 | | x > 100) {
   // out of bounds
                                   in bounds
} else {
    // in bounds
                               goto epiloque;
return 0;
                               x out of bounds:
                                    // out of bounds
                               epilogue:
                                    return 0;
```

```
if (y < 10 || z > 50) {
    // condition met
} else {
    // condition not met
}
return 1;
```



```
if (y < 10 || z > 50) {
    // condition met
} else {
    // condition not met
}
return 1;
```

if (y < 10) goto condition met; if (z > 50) goto condition met; goto condition not met; condition met: // condition met goto epilogue; condition not met: // condition not met epilogue: return 1;

```
if (y < 10 || z > 50) {
    // condition met
} else {
    // condition not met
}
return 1;
```

if (y < 10) goto condition_met; if (z > 50) goto condition_met; // condition not met goto epilogue; condition_met: // condition met

epilogue:

return 1;

Your turn to try

```
if (y < 10 || (z > 50 && w < 5)) {
    // condition met
} else {
    // condition not met
}
return 1;</pre>
```

```
if (y < 10 || (z > 50 && w < 5)) {
    // condition met
} else {
    // condition not met
}
return 1;</pre>
```

if (y < 10) goto condition_met; if (z <= 50) goto condition_not_met; if (w >= 5) goto condition_not_met; condition_met: // condition met goto epilogue; condition_not_met: // condition not met epilogue: return 1;

Simplifying loop structures

- **for** loops should be broken down to **while** loops
- while loops should be broken down into if/goto

General structure:

- loop init
- loop condition (do we need to *exit* the loop?)
- loop body
- loop step
- loop end

Use labels to show structure!

Simplifying for loops: Counting

Counting

```
int i;
                                loop_i_to_10__init:
                                    i = 0:
                                loop_i_to_10__cond:
                                    if (i >= 10) goto loop_i_to_10__end;
int i = 0;
                                loop_i_to_10__body:
while (i < 10) {
    printf("%d\n", i);</pre>
                                    printf("%d", i);
                                    putchar('\n');
    i++;
                                loop_i_to_10__step:
                                    i++:
                                    goto loop_i_to_10__cond;
                                loop_i_to_10__end:
                                    // ...
```

Exercise: Sum 100 squares

Convert to MIPS

```
int sum = 0;
for (int i = 1; i <= 100; i++) {
    sum += i * i;
}
```

Sidenote: C break/continue

break can be used in a loop to completely exit the loop. The loop condition here makes this look like an infinite loop:

```
while (1) {
    int c = getchar();
    if (c == EOF) break;
}
```

but **break** means it's possible for the loop to be exited.

In simplified C/MIPS, a **break** is really just equivalent to going to the loop's end label.

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Sidenote: C break/continue

continue can be used to proceed to the next iteration of a for loop.

This would be a (terrible) way to print even numbers:

```
for (int i = 0; i < 10; i++) {
    if (i % 2 != 0) continue;
    printf("%d\n", i);
}</pre>
```

In simplified C/MIPS, a **continue** is really just equivalent to going to the loop's step label.

Beware: Writing this as a while loop in C needs care not to miss the i++

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What did we learn today?

- MIPS
 - recap of basics from lecture 1
 - system calls
 - printing out and reading in integers, and chars
 - printing out strings
 - simplified C
 - \circ control
 - goto statements
 - if statements,
 - boolean expressions
 - loops

Feedback Please!

Your feedback is valuable!

If you have any feedback from today's lecture, please follow the link below or use the QR Code.

Please remember to keep your feedback constructive, so I can action it and improve your learning experience.

https://forms.office.com/r/EYPYy0KG5E

Reach Out

Content Related Questions: Forum

Admin related Questions email: <u>cs1521@cse.unsw.edu.au</u>

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