COMP1521 21T3 — Files

https://www.cse.unsw.edu.au/~cs1521/21T3/
Operating system - What Does it Do.

- Operating system sits between the user and the hardware
- Operating system effectively provides a virtual machine to each user
- This virtual machine is much simpler and more convenient than real machine
- The virtual machine interface can be consistent across different hardware.
  - program can portably access hardware across different hardware configurations
  - linux available for almost all suitable hardware
- can coordinate/share access to resources between users
- can provide privileges/security
needs hardware to provide a **privileged** mode which:
- allows access to all hardware/memory
- Operating System (kernel) runs in **privileged** mode
- allows transfer to running code a **non-privileged** mode

needs hardware to provide a **non-privileged** mode which:
- prevents access to hardware
- limits access to memory
- provides mechanism to make requests to operating system

operating system requests are called **system calls**
- system calls transfers execution back to kernel code in **privileged** mode
System Call - What is It

- System call allows programs to request hardware operations.
- System call transfers execution to OS code in **privileged** mode:
  - Includes arguments specifying details of request being made.
  - OS checks if operation is valid and permitted.
  - OS carries out the operation.
  - Transfers execution back to user code in **non-privileged** mode.
- Different operating systems have different system calls:
  - E.g., Linux provides completely different system calls to Windows.
- Linux provides 400+ system calls.
- Operations likely to be provided by system calls:
  - Read/write bytes to a file.
  - Request more memory.
  - Create a process (run a program).
  - Terminate a process.
  - Send or receive information via a network.
System Call in SPIM

- SPIM provides a virtual machine which can execute MIPS programs
- SPIM also provides a tiny operating system
- Small number of SPIM system calls for I/O and memory allocation
- Access is via the `syscall` instruction
- MIPS programs running on real hardware also use `syscall`
  - On Linux `syscall`, will pass execution to Linux kernel
- SPIM system calls are designed for students writing tiny programs
  - E.g. SPIM system call 1 - print an integer
- System calls on real operating systems more general
  - E.g. system call might be write n bytes
- In real operating system library systems calls more general
  - Library functions like `printf` provide convenient operations
// hello world implemented with a direct syscall
#include <unistd.h>

int main(void) {
    char bytes[16] = "Hello, Andrew!\n";
    // argument 1 to syscall is system call number, 1 == write
    // remaining arguments are specific to each system call
    // write system call takes 3 arguments:
    // 1) file descriptor, 1 == stdout
    // 2) memory address of first byte to write
    // 3) number of bytes to write
    syscall(1, 1, bytes, 15); // prints Hello, Andrew! on stdout
    return 0;
}
Using read & write system calls to copy stdin to stdout

```c
// copy stdin to stdout with read & write syscalls
while (1) {
    char bytes[4096];
    // system call number 0 == read
    // read system call takes 3 arguments:
    // 1) file descriptor, 1 == stdin
    // 2) memory address to put bytes read
    // 3) maximum number of bytes read
    // returns number of bytes actually read
    long bytes_read = syscall(0, 0, bytes, 4096);
    if (bytes_read <= 0) {
        break;
    }
    syscall(1, 1, bytes, bytes_read); // prints bytes to stdout
}
```

source code for cat_syscalls.c

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What Really are Files and Directories?

- **file systems** manage persistent stored data e.g. on magnetic disk or SSD

- On Unix-like systems:
  - a **file** is sequence (array) of zero or more bytes.
  - no meaning for bytes associated with file
    - file metadata doesn’t record that it is e.g. ASCII, MP4, JPG, ...
    - Unix-like files are just bytes
  - a **directory** is an object containing zero or more files or directories.

- file systems maintain metadata for files & directories, e.g. permissions

- system calls provide operations to manipulate files.

- libc provides a low-level API to manipulate files.

- stdio.h provides more portable, higher-level API to manipulate files.
Unix-like Files & Directories

- Unix-like filenames are sequences of 1 or more bytes.
  - filenames can contain any byte except \x00 and \x2F
  - \x00 bytes (ASCII ‘\0’) used to terminate filenames
  - \x2F bytes (ASCII ‘/’) used to separate components of pathnames.
  - maximum filename length, depends on file system, typically 255

- Two filenames can not be used - they have a special meaning:
  - current directory
  - parent directory

- Some programs (shell, ls) treat filenames starting with specially.

- Unix-like directories are sets of files or directories
Unix/Linux Pathnames

- Files & directories accessed via pathnames, e.g: `/home/z5555555/lab07/main.c`
- **absolute** pathnames start with a leading `/` and give full path from root
  - e.g. `/usr/include/stdio.h`, `/cs1521/public_html/`
- every process (running program) has an associated **absolute** pathname called the **current working directory** (CWD)
- shell command `pwd` prints CWD
- **relative** pathname do not start with a leading `/` e.g. `../../another/path/prog.c`, `./a.out`, `main.c`
- **relative** pathnames appended to CWD of process using them
- Assume process CWD is `/home/z5555555/lab07/`
  - `main.c` translated to absolute path `/home/z5555555/lab07/main.c`
  - `../a.out` translated to absolute path `/home/z5555555/lab07/..../a.out`
  - which is equivalent to absolute path `/home/z5555555/a.out`
Everything is a File

- Originally files only managed data stored on a magnetic disk.
- Unix philosophy is: *Everything is a File*.
- File system can be used to access:
  - files
  - directories (folders)
  - storage devices (disks, SSD, ...)
  - peripherals (keyboard, mouse, USB, ...)
  - system information
  - inter-process communication
  - ...
Unix/Linux file system is tree-like

- We think of file-system as a **tree**
- But beware if you follow symbolic links it is a **graph**.
  - and you may infinitely loop attempting to traverse a file system
Metadata for file system objects is stored in **inodes**, which hold:

- location of file contents in file systems
- file type (regular file, directory, ...)
- file size in byte
- file ownership
- file access permissions - who can read, write, execute the file
- timestamps - time of creation/access/update

Note: file systems add much complexity to improve performance

- e.g. very small files might be stored in an inode itself
unix-like file systems effectively have an array of inodes

- every inode has a **inode-number** (or **i-number**) - its index in this array
- directories are effectively a list of (name, inode-number) pairs
- inode-number uniquely identify files within a filesystem
  - just a zid uniquely identifies a student within UNSW

- **ls -i** prints **inode-number**, e.g.:

  $ ls -i file.c
  109988273 file.c
  109988273 file.c
  $
Access to files by name proceeds (roughly) as...

- open directory and scan for *name*
- if not found, “No such file or directory”
- if found as \((name, inumber)\), access inode table \(\text{inodes}[inumber]\)
- collect file metadata and...
  - check file access permissions given current user/group
    - if don’t have required access, “Permission denied”
  - collect information about file’s location and size
  - update access timestamp
- use data in inode to access file contents
Hard Links & Symbolic Links

File system *links* allow multiple paths to access the same file

**Hard links**
- multiple names referencing the same file (inode)
- the two entries must be on the same filesystem
- all hard links to a file have equal status
- file destroyed when last hard link removed
- can not create a (extra) hard link to directories

**Symbolic links (symlinks)**
- point to another path name
- accessing the symlink (by default) accesses the file being pointed to
- symbolic link can point to a directory
- symbolic link can point to a pathname on another filesystems
- symbolic links don’t have permissions (just a pointer)
$ echo 'Hello Andrew' > hello
$ ln hello hola  # create hard link
$ ln -s hello selamat  # create symbolic link
$ ls -l hello hola selamat
-rw-r--r-- 2 andrewt 13 Oct 23 16:18 hello
-rw-r--r-- 2 andrewt 13 Oct 23 16:18 hola
lrwxrwxrwx 1 andrewt 5 Oct 23 16:20 selamat -> hello
$ cat hello
Hello Andrew
$ cat hola
Hello Andrew
$ cat selamat
Hello Andrew
System Calls to Manipulate files

Unix presents a uniform interface to file system objects

- system calls manipulate objects as a *stream of bytes*
- accessed via a *file descriptor*
  - file descriptors are small integers
  - index to a per-process operating system table (array)

Some important system calls:

- **open()** — open a file system object, returning a file descriptor
- **close()** — stop using a file descriptor
- **read()** — read some bytes into a buffer from a file descriptor
- **write()** — write some bytes from a buffer to a file descriptor
- **lseek()** — move to a specified offset within a file
- **stat()** — get file system object metadata
Using system call directly to create a file

// cp <file1> <file2> with syscalls, no error handling!
// system call number 2 == open, takes 3 arguments:
// 1) address of zero-terminated string containing file pathname
// 2) bitmap indicating whether to write, read, ... file
// 0x41 == write to file creating if necessary
// 3) permissions if file will be newly created
// 0644 == readable to everyone, writeable by owner

long read_file_descriptor = syscall(2, argv[1], 0, 0);
long write_file_descriptor = syscall(2, argv[2], 0x41, 0644);

while (1) {
    char bytes[4096];
    long bytes_read = syscall(0, read_file_descriptor, bytes, 4096);
    if (bytes_read <= 0) {
        break;
    }
    syscall(1, write_file_descriptor, bytes, bytes_read);
}

source code for cp_syscalls.c

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On Unix-like systems there are C library functions corresponding to each system call,
  - e.g. open, read, write, close
  - the `syscall` function is not used in normal coding

These functions are not portable - absent from many platforms/implementations

POSIX standardizes some of these functions
  - some non-Unix systems provide implementations of these functions

better to use functions from standard C library, available everywhere
  - e.g fopen, fgets, fputc from `stdio.h`
  - on unix-like systems these will call open, read, write,

but sometimes need to use lower level functions
Unix-like (POSIX) systems add some extra file-system-related C types in these include files:

```c
#include <sys/types.h>
#include <sys/stat.h>
```

- **off_t** — offsets within files
  - typically `int64_t` - signed to allow backward references
- **size_t** — number of bytes in some object
  - typically `uint64_t` - unsigned since objects can’t have negative size
- **ssize_t** — sizes of read/written bytes
  - like `*size_t`, but signed to allow for error values
- **struct stat** — file system object metadata
  - stores information about file, not its contents
  - requires other types: `ino_t`, `dev_t`, `time_t`, `uid_t`, ...
int open(char *pathname, int flags)

- open file at **pathname**, according to **flags**
- **flags** is a bit-mask defined in `<fcntl.h>`
  - O_RDONLY — open for reading
  - O_WRONLY — open for writing
  - O_APPEND — append on each write
  - O_RDWR — open object for reading and writing
  - O_CREAT — create file if doesn’t exist
  - O_TRUNC — truncate to size 0
- flags can be combined e.g. (O_WRONLY|O_CREAT)
- if successful, return file descriptor (small non-negative int)
- if unsuccessful, return -1 and set errno

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C library wrapper for close system call

```c
int close(int fd)
```

- release open file descriptor `fd`
- if successful, return 0
- if unsuccessful, return −1 and set `errno`
  - could be unsuccessful if `fd` is not an open file descriptor
    - e.g. if `fd` has already been closed

An aside: removing a file e.g. via `rm`

- removes the file's entry from a directory
- but the inode and data persist until
  - all references to the inode from other directories are removed
  - all processes accessing the file `close()` their file descriptor
- after this, the inode and the space used for file contents is recycled
C library wrapper for read system call

\[ \text{ssize_t read(int fd, void *buf, size_t count)} \]

- read (up to) \textbf{count} bytes from \textbf{fd} into \textbf{buf}
  - \textbf{buf} should point to array of at least \textbf{count} bytes
  - read does (can) not check \textbf{buf} points to enough space

- if successful, number of bytes actually read is returned
- 0 returned, if no more bytes to read
- -1 returned if error and \texttt{errno} set to reason
- next call to \textbf{read} will return next bytes from file
- repeated calls to \textbf{reads} will yield entire contents of file
  - associated with a file descriptor is “current position” in file
  - can also modify this position with \texttt{lseek}
C library wrapper for write system call

```c
ssize_t write(int fd, const void *buf, size_t count)
```

- attempt to write \textbf{count} bytes from \textit{buf} into
  - stream identified by file descriptor \textit{fd}
- if successful, number of bytes actually written is returned
- if unsuccessful, return \text{−1} and set \textbf{errno}
- does (can) not check \textit{buf} points to \textbf{count} bytes of data
- next call to \textbf{write} will follow bytes already written
- file often created by repeated calls to \textbf{write}
  - associated with a file descriptor is “current position” in file
  - can also modify this position with \textbf{lseek}
// hello world implemented with libc
#include <unistd.h>
int main(void) {
    char bytes[16] = "Hello, Andrew!\n";
    // write takes 3 arguments:
    // 1) file descriptor, 1 == stdout
    // 2) memory address of first byte to write
    // 3) number of bytes to write
    write(1, bytes, 15); // prints Hello, Andrew! on stdout
    return 0;
}
Using read & write to copy stdin to stdout

```c
while (1) {
    char bytes[4096];
    // system call number 0 == read
    // read system call takes 3 arguments:
    // 1) file descriptor, 1 == stdin
    // 2) memory address to put bytes read
    // 3) maximum number of bytes read
    // returns number of bytes actually read
    ssize_t bytes_read = read(0, bytes, 4096);
    if (bytes_read <= 0) {
        break;
    }
    write(1, bytes, bytes_read); // prints bytes to stdout
}
```

source code for cat_libc.c

https://www.cse.unsw.edu.au/~cs1521/21T3/COMP1521.21T3—Files
Using open to copy a file

```c
// open takes 3 arguments:
// 1) address of zero-terminated string containing pathname of file to open
// 2) bitmap indicating whether to write, read, ... file
// 3) permissions if file will be newly created
// 0644 == readable to everyone, writeable by owner
int read_file_descriptor = open(argv[1], O_RDONLY);
int write_file_descriptor = open(argv[2], O_WRONLY | O_CREAT, 0644);
while (1) {
    char bytes[4096];
    ssize_t bytes_read = read(read_file_descriptor, bytes, 4096);
    if (bytes_read <= 0) {
        break;
    }
    write(write_file_descriptor, bytes, bytes_read);
}
```

Source code for cp_libc.c

https://www.cse.unsw.edu.au/~cs1521/21T3/COMP1521_21T3---Files
C library wrapper for lseek system call

```c
off_t lseek(int fd, off_t offset, int whence)
```

- change the ‘current position’ in the file of `fd`
- `offset` is in units of bytes, and can be negative
- `whence` can be one of ...
  - SEEK_SET — set file position to Offset from start of file
  - SEEK_CUR — set file position to Offset from current position
  - SEEK_END — set file position to Offset from end of file

- seeking beyond end of file leaves a gap which reads as 0’s
- seeking back beyond start of file sets position to start of file

Example: `lseek(fd, 0, SEEK_END);` (move to end of file)
Using lseek to read the last byte then the first byte of a file

```c
int read_file_descriptor = open(argv[1], O_RDONLY);
char bytes[1];

// move to a position 1 byte from end of file
// then read 1 byte
lseek(read_file_descriptor, -1, SEEK_END);
read(read_file_descriptor, bytes, 1);
printf("last byte of the file is 0x%02x\n", bytes[0]);

// move to a position 0 bytes from start of file
// then read 1 byte
lseek(read_file_descriptor, 0, SEEK_SET);
read(read_file_descriptor, bytes, 1);
printf("first byte of the file is 0x%02x\n", bytes[0]);
```

source code for lseek.c
Using lseek to read bytes in the middle of a file

```c
printf("first byte of the file is 0x%02x\n", bytes[0]);
// move to a position 41 bytes from start of file
// then read 1 byte
lseek(read_file_descriptor, 41, SEEK_SET);
read(read_file_descriptor, bytes, 1);
printf("42nd byte of the file is 0x%02x\n", bytes[0]);
// move to a position 58 bytes from current position
// then read 1 byte
lseek(read_file_descriptor, 58, SEEK_CUR);
read(read_file_descriptor, bytes, 1);
printf("100th byte of the file is 0x%02x\n", bytes[0]);
```

source code for lseek.c

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- `stdio.h` is part of standard C library
- available in every C implementation that can do I/O
- `stdio.h` functions are portable, convenient & efficient
- use them by default for file operations
- on Unix-like systems they will call open/read/write/...
  - but with buffering for efficiency
FILE *fopen(const char *pathname, const char *mode)

- **stdio.h** equivalent to `open`
- **mode** is string of 1 or more characters including:
  - `r` open text file for reading.
  - `w` open text file for writing truncated to 0 zero length if it exists created if does not exist
  - `a` open text file for writing writes append to it if it exists created if does not exist
- `fopen` returns a **FILE** * pointer
- **FILE** is an opaque struct - we can not access fields

int fclose(FILE *stream)

- **stdio.h** equivalent to `close`
stdio.h - read and writing

```c
int fgetc(FILE *stream)  // read a byte
int fputc(int c, FILE *stream)  // write a byte

char *fputs(char *s, FILE *stream)  // write a string

char *fgets(char *s, int size, FILE *stream)  // read a line

// formatted input
int fscanf(FILE *stream, const char *format, ...)  // formatted output
int fprintf(FILE *stream, const char *format, ...)

// read array of bytes (fgetc + loop often better)
size_t fread(void *ptr, size_t size, size_t nmemb, FILE *stream);

// write array of bytes (fputc + loop often better)
size_t fwrite(const void *ptr, size_t size, size_t nmemb, FILE *stream)
```
stdio.h - using fputc to output bytes

char bytes[] = "Hello, stdio!\n"; // 15 bytes
// write 14 bytes so we don't write (terminating) 0 byte
for (int i = 0; i < (sizeof bytes) - 1; i++) {
    fputc(bytes[i], stdout);
}
// or as we know bytes is 0-terminated
for (int i = 0; bytes[i] != '\0'; i++) {
    fputc(bytes[i], stdout);
}
// or if you prefer pointers
for (char *p = &bytes[0]; *p != '\0'; p++) {
    fputc(*p, stdout);
}

source code for hello_stdio.c

https://www.cse.unsw.edu.au/~cs1521/21T3/COMP1521%2021T3%20—%20Files
**stdio.h** - using `fputs`, `fwrite` & `fprintf` to output bytes

```c
char bytes[] = "Hello, stdio!\n"; // 15 bytes

// fputs relies on bytes being 0-terminated
fputs(bytes, stdout);

// write 14 1 byte items
fwrite(bytes, 1, (sizeof bytes) - 1, stdout);

// %s relies on bytes being 0-terminated
fprintf(stdout, "%s", bytes);
```

source code for `hello_stdio.c`
stdio.h - using fgetc to copy stdin to stdout

// c can not be char (common bug)
// fgetc returns 0..255 and EOF (usually -1)
int c;
// return bytes from the stream (stdin) one at a time
while ((c = fgetc(stdin)) != EOF) {
    fputc(c, stdout); // write the byte to standard output
}

source code for cat_fgetc.c
stdio.h - using fgets to copy stdin to stdout

// return bytes from the stream (stdin) line at a time
// BUFSIZ is defined in stdio.h - its an efficient value to use
// but any value would work
char line[BUFSIZ];
while (fgets(line, sizeof line, stdin) != NULL) {
  fputs(line, stdout);
}

// NOTE: fgets returns a null-terminated string
// in other words a 0 byte marks the end of the bytes read
// fgets can not be used to read bytes which are 0
// fputs takes a null-terminated string
// so fputs can not be used to write bytes which are 0
// hence you can't use fget/fputs for binary data e.g. jpgs

source code for cat_fgets.c

https://www.cse.unsw.edu.au/~cs1521/21T3/COMP1521-21T3---Files
while (1) {
    char bytes[4096];
    ssize_t bytes_read = fread(bytes, 1, sizeof bytes, stdin);
    if (bytes_read <= 0) {
        break;
    }
    fwrite(bytes, 1, bytes_read, stdout);
}
// create file "hello.txt" containing 1 line: Hello, Andrew
#include <stdio.h>
#include <stdlib.h>

int main(int argc, char *argv[]) {
    FILE *output_stream = fopen("hello.txt", "w");
    if (output_stream == NULL) {
        perror("hello.txt");
        return 1;
    }
    fprintf(output_stream, "Hello, Andrew!\n");
    // fclose will flush data to file
    // best to close file ASAP
    // but doesn't matter as file autoamtically closed on exit
    fclose(output_stream);
    return 0;
}

source code for create_file_fopen.c
FILE *input_stream = fopen(argv[1], "rb");
if (input_stream == NULL) {
    perror(argv[1]);  // prints why the open failed
    return 1;
}

FILE *output_stream = fopen(argv[2], "wb");
if (output_stream == NULL) {
    perror(argv[2]);
    return 1;
}

int c;  // not char!
while ((c = fgetc(input_stream)) != EOF) {
    fputc(c, output_stream);
}

fclose(input_stream);  // optional as close occurs
fclose(output_stream);  // automatically on exit

source code for cp_fgetc.c
```c
FILE *input_stream = fopen(argv[1], "rb");
FILE *output_stream = fopen(argv[2], "wb");

// this will be slightly faster than an a fgetc/fputc loop
while (1) {
    char bytes[BUFSIZ];
    size_t bytes_read = fread(bytes, 1, sizeof bytes, input_stream);
    if (bytes_read <= 0) {
        break;
    }
    fwrite(bytes, 1, bytes_read, output_stream);
}
fclose(input_stream); // optional as close occurs
fclose(output_stream); // automatically on exit
```
```c
int fseek(FILE *stream, long offset, int whence);
```

- **fseek** is stdio equivalent to lseek
- Like lseek, **offset** can be positive or negative
- Like lseek, **whence** can be SEEK_SET, SEEK_CUR or SEEK_END making **offset** relative to file start, current position or file end

```c
int fflush(FILE *stream);
```

- Flush any buffered data on output stream
Using fseek to read the last byte then the first byte of a file

FILE *input_stream = fopen(argv[1], "rb");
// move to a position 1 byte from end of file
// then read 1 byte
fseek(input_stream, -1, SEEK_END);
printf(”last byte of the file is 0x%02x
”, fgetc(input_stream));
// move to a position 0 bytes from start of file
// then read 1 byte
fseek(input_stream, 0, SEEK_SET);
printf(”first byte of the file is 0x%02x
”, fgetc(input_stream));

• NOTE: important error checking is missing above
Using fseek to read bytes in the middle of a file

// move to a position 41 bytes from start of file
// then read 1 byte
fseek(input_stream, 41, SEEK_SET);
printf("42nd byte of the file is 0x%02x\n", fgetc(input_stream));
// move to a position 58 bytes from current position
// then read 1 byte
fseek(input_stream, 58, SEEK_CUR);
printf("100th byte of the file is 0x%02x\n", fgetc(input_stream));

NOTE: important error checking is missing above
Using fseek to change a random file bit

FILE *f = fopen(argv[1], "r+");  // open for reading and writing
fseek(f, 0, SEEK_END);  // move to end of file
long n_bytes = ftell(f);  // get number of bytes in file
srandom(time(NULL));  // initialize random number generator with current time

long target_byte = random() % n_bytes;  // pick a random byte
fseek(f, target_byte, SEEK_SET);  // move to byte
int byte = fgetc(f);  // read byte
int bit = random() % 8;  // pick a random bit
int new_byte = byte ^ (1 << bit);  // flip the bit
fseek(f, -1, SEEK_CUR);  // move back to same position
fputc(new_byte, f);  // write the byte
fclose(f);

* source code for fuzz.c

- random changes to search for errors/vulnerabilities called fuzzing
// Create a 16 terabyte sparse file
// error checking omitted for clarity
#include <stdio.h>

int main(void) {

    FILE *f = fopen("sparse_file.txt", "w");
    fprintf(f, "Hello, Andrew!\n");
    fseek(f, 16L * 1000 * 1000 * 1000 * 1000, SEEK_CUR);
    fprintf(f, "Goodbye, Andrew!\n");
    fclose(f);

    return 0;
}

source code for create_gigantic_file.c

- almost all the 16Tb are zeros which the file system doesn't actually store
as we often read/write to stdin/stdout. `stdio.h` provides convenience functions, we can use:

```c
int getchar()  // fgetc(stdin)
int putchar(int c)  // fputc(c, stdin)
int puts(char *s)  // fputs(s, stdout)
int scanf(char *format, ...)  // fscanf(stdin, format, ...)
int printf(char *format, ...)  // fprintf(stdout, format, ...)
char *gets(char *s);  // NEVER USE
```
stdio.h provides useful functions which operate on strings

\[
\text{int snprintf(char *str, size_t size, const char *format, ...);}\\
\]
- like printf, but output goes to char array \text{str}
- handy for creating strings passed to other functions
- do not use unsafe related function: ‘sprintf

\[
\text{int sscanf(const char *str, const char *format, ...);}\\
\]
- like scanf, but input comes from char array \text{str}

\[
\text{int sprintf(char *str, const char *format, ...);} // DO NOT USE\\
\]
- like \text{snprintf} but dangerous because can overflow \text{str}
C library wrapper for stat system call

```c
int stat(const char *pathname, struct stat *statbuf)
```

- Returns metadata associated with `pathname` in `statbuf`
- Metadata returned includes:
  - Inode number
  - Type (file, directory, symbolic link, device)
  - Size of file in bytes (if it is a file)
  - Permissions (read, write, execute)
  - Times of last access/modification/status-change
- Returns -1 and sets `errno` if metadata not accessible

```c
int fstat(int fd, struct stat *statbuf)
```

- Same as `stat()` but gets data via an open file descriptor

```c
int lstat(const char *pathname, struct stat *statbuf)
```

- Same as `stat()` but doesn’t follow symbolic links
```c
struct stat {
    dev_t st_dev;    /* ID of device containing file */
    ino_t st_ino;    /* Inode number */
    mode_t st_mode;  /* File type and mode */
    nlink_t st_nlink; /* Number of hard links */
    uid_t st_uid;    /* User ID of owner */
    gid_t st_gid;    /* Group ID of owner */
    dev_t st_rdev;   /* Device ID (if special file) */
    off_t st_size;   /* Total size, in bytes */
    blksize_t st_blksize; /* Block size for filesystem I/O */
    blkcnt_t st_blocks; /* Number of 512B blocks allocated */
    struct timespec st_atim; /* Time of last access */
    struct timespec st_mtim; /* Time of last modification */
    struct timespec st_ctim; /* Time of last status change */
};
```
The **st_mode** field of the struct stat is a bitwise-or of these values (and others):

- **S_IFLNK 0120000** symbolic link
- **S_IFREG 0100000** regular file
- **S_IFBLK 0060000** block device
- **S_IFDIR 0040000** directory
- **S_IFCHR 0020000** character device
- **S_IFIFO 0010000** FIFO
- **S_IRUSR 0000400** owner has read permission
- **S_IWUSR 0000200** owner has write permission
- **S_IXUSR 0000100** owner has execute permission
- **S_IRGRP 0000040** group has read permission
- **S_IWGRP 0000020** group has write permission
- **S_IXGRP 0000010** group has execute permission
- **S_IROTH 0000004** others have read permission
- **S_IWOTH 0000002** others have write permission
- **S_IXOTH 0000001** others have execute permission
Using stat

```c
struct stat s;
if (stat(pathname, &s) != 0) {
    perror(pathname);
    exit(1);
}
printf("ino = %10ld # Inode number\n", s.st_ino);
printf("mode = %10o # File mode \n", s.st_mode);
printf("nlink =%10ld # Link count \n", (long)s.st_nlink);
printf("uid = %10u # Owner uid\n", s.st_uid);
printf("gid = %10u # Group gid\n", s.st_gid);
printf("size = %10ld # File size (bytes)\n", (long)s.st_size);
printf("mtime =%10ld # Modification time (seconds since 1/1/70)\n", (long)s.st_mtime);
```

source code for stat.c

[https://www.cse.unsw.edu.au/~cs1521/21T3/](https://www.cse.unsw.edu.au/~cs1521/21T3/)
mkdir

**int mkdir(const char *pathname, mode_t mode)**

- create a new directory called **pathname** with permissions **mode**
- If **pathname** is e.g. `a/b/c/d`
  - all of the directories `a`, `b` and `c` must exist
  - directory `c` must be writeable to the caller
  - directory `d` must not already exist
- The new directory contains two initial entries
  - `. ` is a reference to itself
  - `.. ` is a reference to its parent directory
- Returns 0 if successful, returns -1 and sets **errno** otherwise

For example:

```c
mkdir("newDir", 0755);
```
Example of using mkdir to create directories

```c
#include <stdio.h>
#include <sys/stat.h>

// create the directories specified as command-line arguments
int main(int argc, char *argv[]) {
    for (int arg = 1; arg < argc; arg++) {
        if (mkdir(argv[arg], 0755) != 0) {
            perror(argv[arg]); // prints why the mkdir failed
            return 1;
        }
    }
    return 0;
}
```

source code for mkdir.c

https://www.cse.unsw.edu.au/~cs1521/21T3/
Other useful Linux (POSIX) functions

```c
chmod(char *pathname, mode_t mode) // change permission of file/

unlink(char *pathname) // remove a file/directory/

rename(char *oldpath, char *newpath) // rename a file/directory

chdir(char *path) // change current working directory

getcwd(char *buf, size_t size) // get current working directory

link(char *oldpath, char *newpath) // create hard link to a file

symlink(char *target, char *linkpath) // create a symbolic link
```
file permissions

- file permissions are separated into three types:
  - **read** - permission to get bytes of file
  - **write** - permission to change bytes of file
  - **execute** - permission to execute file

- read/write/execute often represented as bits of an octal digit

- file permissions are specified for 3 groups of users:
  - **owner** - permissions for the file owner
  - **group** - permissions for users in the group of the file
  - **other** - permissions for any other user
changing file permissions

```
// first argument is mode in octal
mode_t mode = strtol(argv[1], &end, 8);

// check first argument was a valid octal number
if (argv[1][0] == '\0' || end[0] != '\0') {
    fprintf(stderr, "%s: invalid mode: %s\n", argv[0], argv[1]);
    return 1;
}

for (int arg = 2; arg < argc; arg++) {
    if (chmod(argv[arg], mode) != 0) {
        perror(argv[arg]); // prints why the chmod failed
        return 1;
    }
}
```

source code for chmod.c

https://www.cse.unsw.edu.au/~cs1521/21T3/COMP1521_21T3—Files
removing files

```c
int main(int argc, char *argv[]) {
    for (int arg = 1; arg < argc; arg++) {
        if (unlink(argv[arg]) != 0) {
            perror(argv[arg]); // prints why the unlink failed
            return 1;
        }
    }
    return 0;
}

source code for rm.c

$ dcc rm.c
$ ./a.out rm.c
$ ls -l rm.c
ls: cannot access 'rm.c': No such file or directory
```
int main(int argc, char *argv[]) {
    if (argc != 3) {
        fprintf(stderr, "Usage: %s <old-filename> <new-filename>\n", argv[0]);
        return 1;
    }
    char *old_filename = argv[1];
    char *new_filename = argv[2];
    if (rename(old_filename, new_filename) != 0) {
        fprintf(stderr, "%s rename %s %s:\n", argv[0], old_filename, new_filename);
        perror(""');
        return 1;
    }
    return 0;
}
// use repeated chdir("..") to climb to root of the file system
char pathname[PATH_MAX];

while (1) {
    if (getcwd(pathname, sizeof(pathname)) == NULL) {
        perror("getcwd");
        return 1;
    }

    printf("getcwd() returned %s\n", pathname);
    if (strcmp(pathname, "/") == 0) {
        return 0;
    }
    if (chdir("..") != 0) {
        perror("chdir");
        return 1;
    }
}

source code for getcwd.c

https://www.cse.unsw.edu.au/~cs1521/21T3/
making a 1000-deep directory

```c
for (int i = 0; i < 1000; i++) {
    char dirname[256];
    snprintf(dirname, sizeof dirname, "d%d", i);
    if (mkdir(dirname, 0755) != 0) {
        perror(dirname);
        return 1;
    }
    if (chdir(dirname) != 0) {
        perror(dirname);
        return 1;
    }
}
char pathname[1000000];
if (getcwd(pathname, sizeof pathname) == NULL) {
    perror("getcwd");
    return 1;
}
printf("Current directory now: %s\n", pathname);
```

source code for nest_directories.c

https://www.cse.unsw.edu.au/~cs1521/21T3/COMP1521.21T3—Files
int main(int argc, char *argv[]) {
    char pathname[256] = "hello.txt";
    // create a target file
    FILE *f1;
    if ((f1 = fopen(pathname, "w")) == NULL) {
        perror(pathname);
        return 1;
    }
    fprintf(f1, "Hello Andrew!\n");
    fclose(f1);
}

source code for many_links.c
creating 1000 hard links to a file (checking the file)

for (int i = 0; i < 1000; i++) {
    printf("Verifying \'%s\' contains: ", pathname);
    FILE *f2;
    if ((f2 = fopen(pathname, "r")) == NULL) {
        perror(pathname);
        return 1;
    }
    int c;
    while ((c = fgetc(f2)) != EOF) {
        fputc(c, stdout);
    }
    fclose(f2);
}

source code for many_links.c
```c
char new_pathname[256];
snprintf(new_pathname, sizeof new_pathname, "hello_%d.txt", i);
printf("Creating a link %s -> %s\n", new_pathname, pathname);
if (link(pathname, new_pathname) != 0) {
    perror(pathname);
    return 1;
}
return 0;
```

source code for many_links.c


```c
#include <sys/types.h>
#include <dirent.h>

// open a directory stream for directory name
DIR *opendir(const char *name);

// return a pointer to next directory entry
struct dirent *readdir(DIR *dirp);

// close a directory stream
int closedir(DIR *dirp);
```
for (int arg = 1; arg < argc; arg++) {
    DIR *dirp = opendir(argv[arg]);
    if (dirp == NULL) {
        perror(argv[arg]);  // prints why the open failed
        return 1;
    }
    struct dirent *de;
    while ((de = readdir(dirp)) != NULL) {
        printf("%ld %s\n", de->d_ino, de->d_name);
    }
    closedir(dirp);
}
int array[10] = { 10, 11, 12, 13, 14, 15, 16, 17, 18, 19 };  
FILE *f = fopen("array.save", "w");  
if (f == NULL) {  
  perror("array.save");  
  return 1;  
}  
// assuming int are 4 bytes, this will  
// write 40 bytes of array to "array.save"  
if (fwrite(array, 1, sizeof array, f) != sizeof array) {  
  perror("array.save");  
  return 1;  
}  
fclose(f);
int array[10];
FILE *f = fopen("array.save", "r");
if (f == NULL) {
    perror("array.save");
    return 1;
}
// read array: NOT-PORTABLE: depends on size of int and byte-order
if (fread(array, 1, sizeof array, f) != sizeof array) {
    perror("array.save");
    return 1;
}
fclose(f);
for (int i = 0; i < 10; i++) {
    printf("%d ", array[i]);
}
printf("\n");

source code for read_array.c

https://www.cse.unsw.edu.au/~cs1521/21T3/
int array[10] = { 10, 11, 12, 13, 14, 15, 16, 17, 18, 19 }; 
int *p = &array[5]; 
FILE *f = fopen("array.save", "w"); 

if (fwrite(array, 1, sizeof array, f) != sizeof array) {
    perror("array.save");
    return 1;
}

if (fwrite(&p, 1, sizeof p, f) != sizeof p) {
    perror("array.save");
    return 1;
}

fclose(f);
int array[10];
int *p;
FILE *f = fopen("array.save", "r");

if (fread(array, 1, sizeof array, f) != sizeof array) {
    perror("array.save");
    return 1;
}

// BROKEN - address of array has almost certainly changed
// BROKEN - so address p needs to point has changed
if (fread(&p, 1, sizeof p, f) != sizeof p) {
    perror("array.save");
    return 1;
}

fclose(f);
```c
int read_file_descriptor = open(argv[1], O_RDONLY);
int write_file_descriptor = open(argv[2], O_WRONLY | O_CREAT, 0644);
// copy bytes 1 at a time
while (1) {
    char bytes[1];
    ssize_t bytes_read = read(read_file_descriptor, bytes, 1);
    if (bytes_read <= 0) {
        break;
    }
    write(write_file_descriptor, bytes, 1);
}
```

- similar to earlier example [source code for cp_libc.c] but one byte at time
$ clang -O3 cp_libc_one_byte.c -o cp_libc_one_byte
$ dd bs=1M count=10 </dev/urandom >random_file
10485760 bytes (10 MB, 10 MiB) copied, 0.183075 s, 57.3 MB/s
$ time ./cp_libc_one_byte random_file random_file_copy
  real 0m5.262s
  user 0m0.432s
  sys  0m4.826s

much slower than previous version which copies 4096 bytes at a time

$ clang -O3 cp_libc.c -o cp_libc
$ time ./cp_libc random_file random_file_copy
  real 0m0.008s
  user 0m0.001s
  sys  0m0.007s

main reason - system calls are expensive
FILE *input_stream = fopen(argv[1], "rb");
if (input_stream == NULL) {
    perror(argv[1]); // prints why the open failed
    return 1;
}
FILE *output_stream = fopen(argv[2], "wb");
if (output_stream == NULL) {
    perror(argv[2]);
    return 1;
}
int c; // not char!
while ((c = fgetc(input_stream)) != EOF) {
    fputc(c, output_stream);
}
fclose(input_stream); // optional as close occurs
fclose(output_stream); // automatically on exit

source code for cp_fgetc.c

https://www.cse.unsw.edu.au/~cs1521/21T3/
I/O Performance & Buffering - stdio Copying 1 Byte Per Time

```
$ clang -O3 cp_fgetc.c -o cp_fgetc
$ time ./cp_fgetc random_file random_file_copy
real  0m0.059s
user  0m0.042s
sys   0m0.009s
```

- at the user level copies 1 byte at time using `fgetc/fputc`
- much faster than coping 1 byte at time using `read/write`
- little slower than coping 4096 bytes at time using `read/write`
- how?
I/O Performance & Buffering - stdio buffering

- Assume stdio buffering size (BUFSIZ) is 4096 (typical)
- Stdio buffers 1 byte fgetc/fputc into 4096 bytes read/write
- First fgetc reads 4096 bytes into an array (input buffer)
  - Next 4095 fgetc calls get byte from array
- First 4095 fputc put bytes into another array (output buffer)
  - Next 4095 fgetc get byte from array
- Output buffer* emptied by exit or main returning
- Data in output buffer
- Program can force empty of output buffer with fflush call
// re-implementation of stdio functions fopen, fgetc, fputc, fclose
// with no buffering and *zero* error handling for clarity
#include <unistd.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <stdint.h>
#include <stdlib.h>
#include <assert.h>
#include <stdio.h>
#define MY_EOF -1
// struct to hold data for a stream
typedef struct my_file {
    int fd;
} my_file_t;

source code for cp_unbuffered.c
my_file_t *my_fopen(char *file, char *mode) {
    int fd = -1;
    if (mode[0] == 'r') {
        fd = open(file, O_RDONLY);
    } else if (mode[0] == 'w') {
        fd = open(file, O_WRONLY | O_CREAT, 0666);
    } else if (mode[0] == 'a') {
        fd = open(file, O_WRONLY | O_APPEND);
    }
    if (fd == -1) {
        return NULL;
    }
    my_file_t *f = malloc(sizeof *f);
    f->fd = fd;
    return f;
}
int my_fgetc(my_file_t *f) {
    uint8_t byte;
    int bytes_read = read(f->fd, &byte, 1);
    if (bytes_read == 1) {
        return byte;
    } else {
        return MY_EOF;
    }
}

source code for cp_unbuffered.c
int my_fputc(int c, my_file_t *f) {
    uint8_t byte = c;
    if (write(f->fd, &byte, 1) == 1) {
        return byte;
    } else {
        return MY_EOF;
    }
}

source code for cp_unbuffered.c
int my_fclose(my_file_t *f) {
    int result = close(f->fd);
    free(f);
    return result;
}
reimplementing stdio.h - buffering (advanced topic)

- reimplementing stdio with input buffering
  
  source code for cp_input_buffered.c

- and output buffering
  
  source code for cp_output_buffered.c
Operating systems provide a *file system*

- as an abstraction over physical storage devices (e.g. disks)
- providing named access to chunks of related data (files)
- providing access (sequential/random) to the contents of files
- allowing files to be arranged in a hierarchy of directories
- providing control over access to files and directories
- managing other metadata associated with files (size, location, ...)

Operating systems also manage other resources

- memory, processes, processor time, i/o devices, networking, ...