

```
-- Model solution for Tut09
--
-- Copyright [2000..2004] Manuel M T Chakravarty
```

```
module Tut09
where
```

```
import Prelude hiding (abs, signum, (++))
```

```
abs           :: Int -> Int
```

```
abs k | k >= 0 = k
```

```
    | otherwise = -k
```

```
signum       :: Int -> Int
```

```
signum k | k > 0 = 1
```

```
    | k == 0 = 0
```

```
    | k < 0 = -1
```

```
{- -----
```

Property: $\text{abs } n * \text{signum } n = n$

PROOF:

Case $n > 0$:

```
= abs n * signum n
= {Due to  $n > 0$ , abs.1}
  n * signum n
= {Due to  $n > 0$ , signum.1}
  n * 1
= {neutral of *}
  n
```

Case $n = 0$:

```
= abs 0 * signum 0
= {abs.1}
  0 * signum 0
= {0 * a = 0}
  0
```

Case $n < 0$:

```
= abs n * signum n
= {Due to  $n < 0$ , abs.2}
  -n * signum n
= {Due to  $n < 0$ , signum.3}
  -n * -1
= {Arithmetic}
  n
```

QED

```
-}
```

```
(++)      :: [a] -> [a] -> [a]
[]        ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)
```

```
{- -----
```

Property: $P(xs) == (xs ++ ys) ++ zs = xs ++ (ys ++ zs)$

PROOF:

Induction over xs.

Base: $P([]) == ([] ++ ys) ++ zs = [] ++ (ys ++ zs)$

Left side:

```
([] ++ ys) ++ zs  
= {++.1}  
ys ++ zs
```

Right side:

```
[] ++ (ys ++ zs)  
= {++.1}  
ys ++ zs
```

Induction: $P(x:xs) == ((x:xs) ++ ys) ++ zs = (x:xs) ++ (ys ++ zs)$

Left side:

```
((x:xs) ++ ys) ++ zs  
= {++.2}  
(x:(xs ++ ys)) ++ zs  
= {++.2}  
x:((xs ++ ys) ++ zs)  
= {Induction hypothesis}  
x:(xs ++ (ys ++ zs))
```

Right side:

```
(x:xs) ++ (ys ++ zs)  
= {++.2}  
x:(xs ++ (ys ++ zs))
```

QED
-}