

```
-- Model solution for Lab11
--
-- Copyright [2000..2001] Gabriele Keller
```

```
module Lab11
```

```
where
```

```
foo :: [Int] -> Int
foo [] = []
foo (x:xs) = (bar xs) + (foo xs)
```

```
{- DERIVATION OF
```

```
T_foo (n), (n: length of input list)
```

```
T_foo (0) = 1
T_foo (n) = T_bar (n-1) + T_foo (n-1)
```

```
1. assume
for T_bar (n) = 2
```

```
T_foo (n) = 2 * n + 1
```

It has linear work complexity, since the timing function is $O(n)$.

Proof:

For $c = 3$ and all $n > n_0 = 2$ we have

$T_foo (n) = 2 * n + 1 < 3 * n$, therefore T_foo is in $O(n)$

```
1. assume
for T_bar (n) = 3 * n + 1
```

```
T_foo (n) = 3 * (n-1) + 1 + T_foo (n-1)
           = 3n - 2 + T_foo (n-1)
```

```
T_foo (n) = 1 + sum (i=1) (n) (3n - 2)
           = 1 - 2n + 3 * sum(i=1) (n) n
           = 1 - 2n + 3 * (n * (n+1) / 0.5)
           = 1.5 n^2 - 0.5n + 1
```

It has linear work complexity, since the timing function is $O(n^2)$,
According to the observations discussed in the lecture, we can
omit all the constant factors and all components of the polynome but the
one with the highest exponent (2 in this case)

```
-}
```

```
Ord a => [a] -> [a]
```

```
bubbleSort :: Ord a => [a] -> [a]
```

```
bubbleSort xs
```

```
  | xs == xs' = xs
```

```
  | otherwise = bubbleSort xs'
```

```
where
```

```
  xs' = bubble xs
```

```
  bubble (x:y:xs)
```

```
    | x < y = x : (bubble (y:xs))
```

```
    | otherwise = y : (bubble (x:xs))
```

```
  bubble xs = xs
```

{- DERIVATION

First, we derive the timing function for T_{bubble} depending on the length n of the input list

$T_{\text{bubble}}(0) = 1$

$T_{\text{bubble}}(1) = 1$

$T_{\text{bubble}}(n) = 2 + T_{\text{bubble}}(n-1)$

*$\Rightarrow T_{\text{bubble}}(n) = 2 * n + 1$*

$T_{\text{bubblesort}}$ does not follow the same pattern as previous derivations, since `bubbleSort` is called recursively on a list of the same length.

*$T_{\text{bubbleSort}}(n) = 1 + T_{\text{bubble}}(n)$, if $xs == xs'$
 $1 + T_{\text{bubble}}(n) + T_{\text{bubblesort}}(n)$, otherwise*

However, if we look more closely at `bubble` we can see that

1) `bubble` applied to a sorted list (and only to the sorted list) returns its input list

2) every element in a list which is to the right (i.e., further back in the list) than its proper position (i.e., the position it will have in the sorted list) will be moved at least by one index to the left with each application of `bubble`

This means that after at most n calls of `bubble`, `bubble xs` has to be sorted.

Therefore, we have in the best case

*$T_{\text{bubbleSort}}(n) = 1 + T_{\text{bubble}}(n) = 2 * n + 2$*

in the worst case

$T_{\text{bubbleSort}}(n) = 1 + T_{\text{bubble}}(n) = 2n^2 + n + 1$

-}