Scheduler Activations

With some slides modified from Raymond Namyst, U. Bordeaux
User-level Threads

User Mode

- Process A
  - Scheduler

- Process B
  - Scheduler

- Process C
  - Scheduler

Kernel Mode

Scheduler
User-level Threads

✓ Fast thread management (creation, deletion, switching, synchronisation…)

✗ Blocking blocks all threads in a process
  – Syscalls
  – Page faults

✗ No thread-level parallelism on multiprocessor
Kernel-Level Threads

User Mode

Kernel Mode

Scheduler

Process A

Process B

Process C
Kernel-level Threads

❌ Slow thread management (creation, deletion, switching, synchronisation…)
  • System calls

✅ Blocking blocks only the appropriate thread in a process

✅ Thread-level parallelism on multiprocessor
Hybrid Multithreading

User Mode

Scheduler

Process A

Scheduler

Process B

Scheduler

Process C

Kernel Mode

Scheduler
Hybrid Multithreading

- Can get real thread parallelism on multiprocessor
- Blocking still a problem!!!
Scheduler Activations

- First proposed by [Anderson et al. 91]
- Idea: Both schedulers co-operate
  - User scheduler uses system calls
  - Kernel scheduler uses upcalls!
- Two important concepts
  - Upcalls
    - Notify the user-level of kernel scheduling events
  - Activations
    - A new structure to support upcalls and execution
      - approximately a kernel thread
    - As many running activations as (allocated) processors
    - Kernel controls activation creation and destruction
Scheduler Activations

- Instead of
  User Space
  Kernel Space
  Hardware

- Rather use the following scheme:
  User Space
  Kernel Space
  Hardware
Upcalls to User-level scheduler

- **New**
  - Allocated a new virtual CPU
  - Can schedule a user-level thread
- **Preempted**
  - DEALLOCATED a virtual CPU
  - Can schedule one less thread
- **Blocked**
  - Notifies thread has blocked
  - Can schedule another user-level thread
- **Unblocked**
  - Notifies a thread has become runnable
  - Must decided to continue current or unblocked thread
Working principle

- Blocking syscall scenario on 2 processors
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Scheduler Activations

- Thread management at user-level
  - Fast
- Real thread parallelism via activations
  - Number of activations (virtual CPU) can equal CPUs
- Blocking (syscall or page fault) creates new activation
  - User-level scheduler can pick new runnable thread.
- Fewer stacks in kernel
  - Blocked activations + number of virtual CPUs
Adoption

- Adopters
  - BSD “Kernel Scheduled Entities”
  - K42
  - Digital UNIX
  - Solaris
  - Mach
- Linux -> kernel threads