

Random Stuff

- No tutorials or comp3891/comp9283 lecture this week
- Release the warm-up exercise tomorrow
- New weeks tutorial questions – probably tomorrow also



Introduction to Operating Systems

Chapter 1 – 1.3

Chapter 1.5 – 1.9

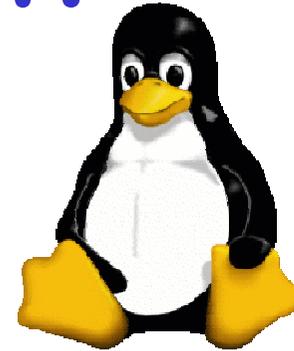


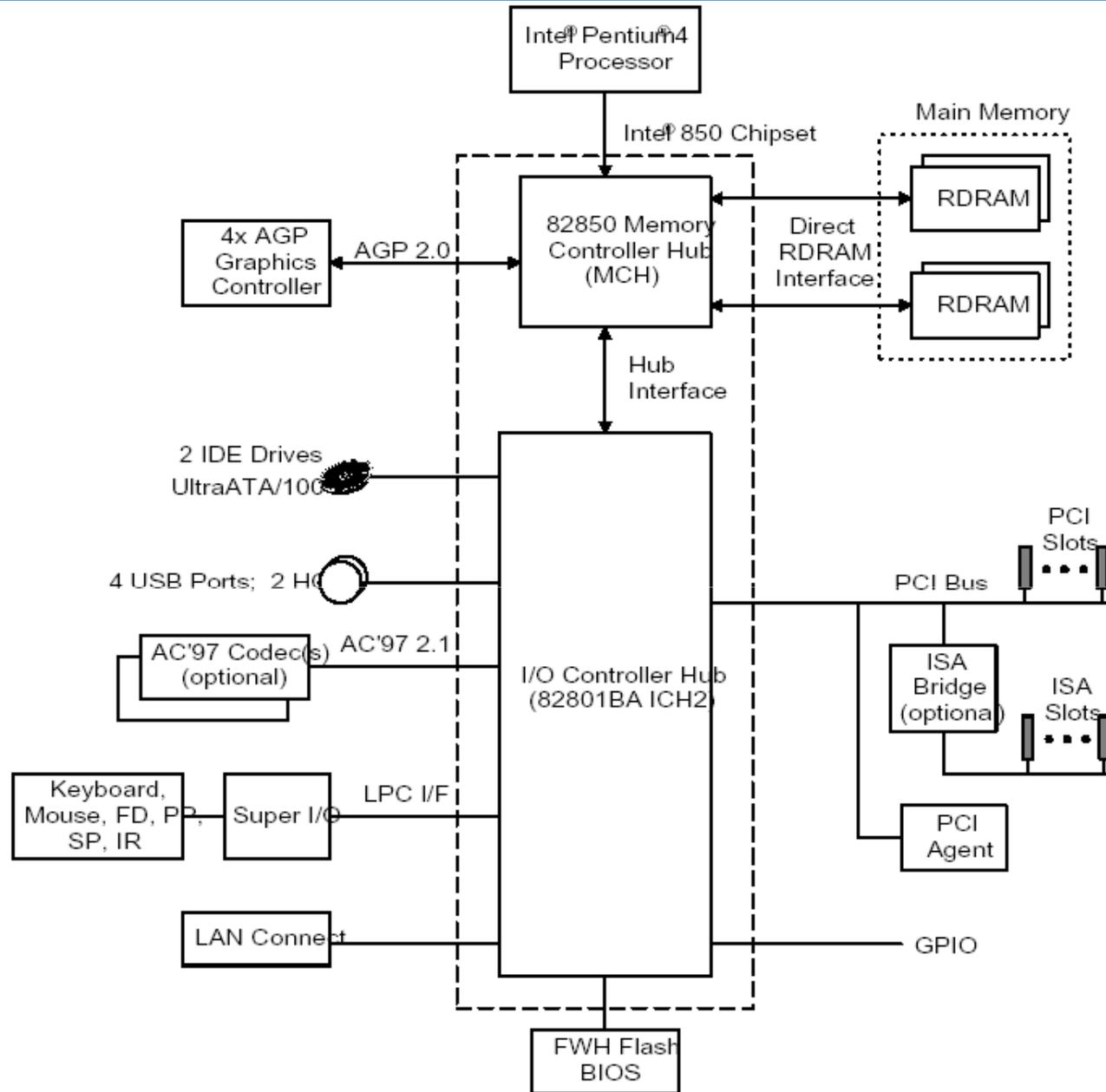
Learning Outcomes

- High-level understand what is an operating system and the role it plays
- A high-level understanding of the structure of operating systems, applications, and the relationship between them.
- Some knowledge of the services provided by operating systems.
- Exposure to some details of major OS concepts.



What is an Operating System?





Viewing the Operating System as an Abstract Machine

- Extends the basic hardware with added functionality
- Provides high-level abstractions
 - More programmer friendly
 - Common core for all applications
- It hides the details of the hardware
 - Makes application code portable



Disk



Memory



CPU

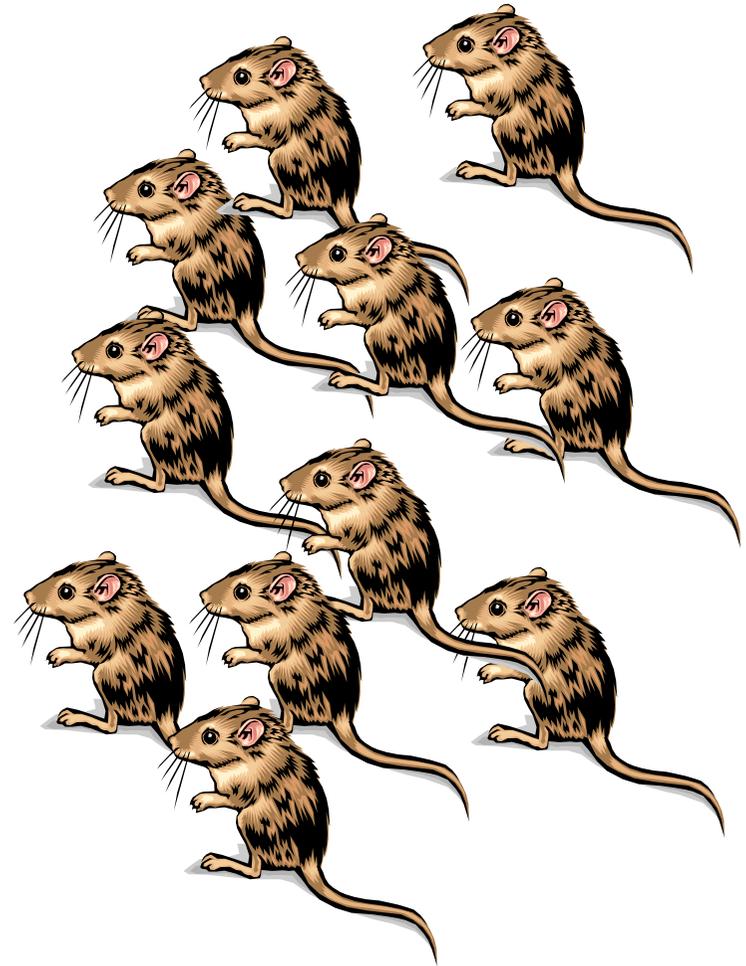


Network

Bandwidth



Users



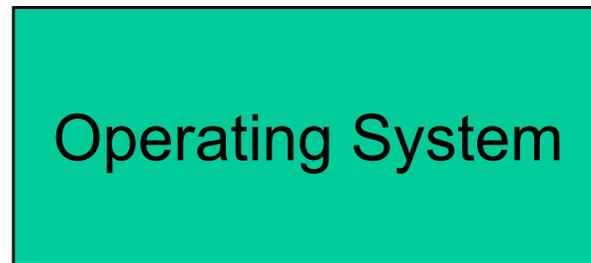
Viewing the Operating System as a Resource Manager

- Responsible for allocating resources to users and processes
- Must ensure
 - No Starvation
 - Progress
 - Allocation is according to some desired policy
 - First-come, first-served; Fair share; Weighted fair share; limits (quotas), etc...
 - Overall, that the system is efficiently used



Traditional View: the Operating System as the Privileged Component

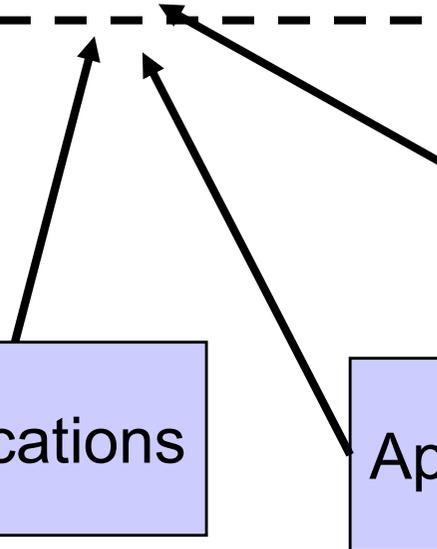
Privileged Mode



Requests
(System Calls)



User Mode



Kernel

- Portion of the operating system that is running in *privileged mode*
- Usually resident in main memory
- Contains fundamental functionality
 - Whatever is required to implement other services
 - Whatever is required to provide security
- Contains most-frequently used functions
- Also called the nucleus or supervisor



The Operating System is Privileged

- Applications should not be able to interfere or bypass the operating system
 - OS can enforce the “extended machine”
 - OS can enforce its resource allocation policies
 - Prevent applications from interfering with each other
- Note: Some Embedded OSs have no privileged component, e.g. PalmOS
 - Can implement OS functionality, but cannot enforce it.
- Note: Some operating systems implement significant OS functionality in user-mode, e.g. User-mode Linux

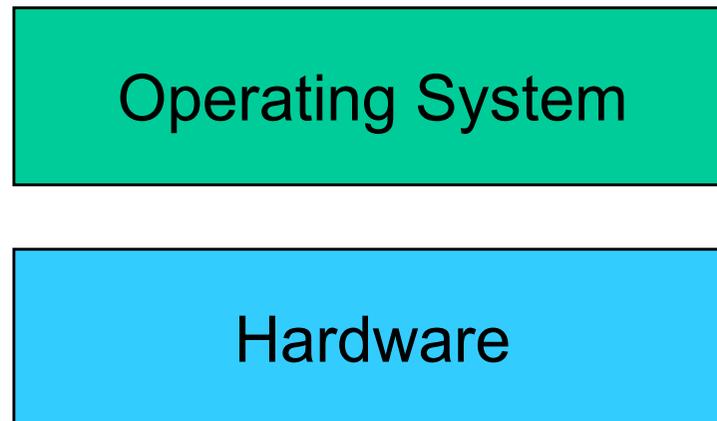


Structure of a Computer System

User Mode



Kernel Mode

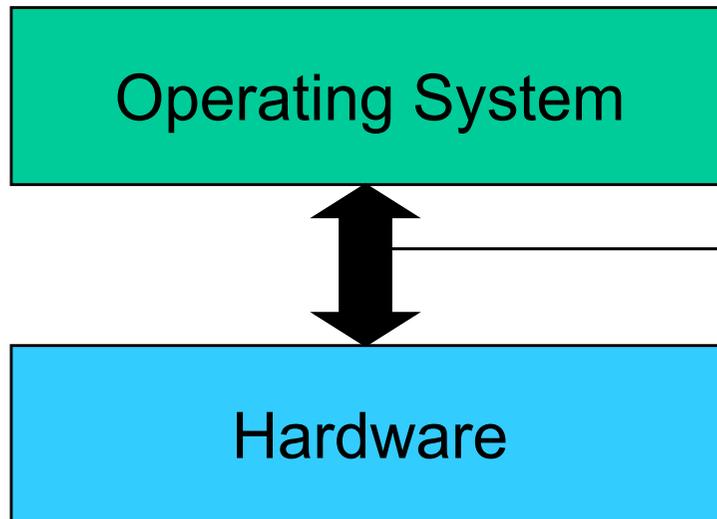


Structure of a Computer System

User Mode



Kernel Mode



Interacts via load and store instructions to CPU and device registers, and interrupts



Structure of a Computer System

User Mode

Application

Interaction via
function calls to
library procedures

System Libraries

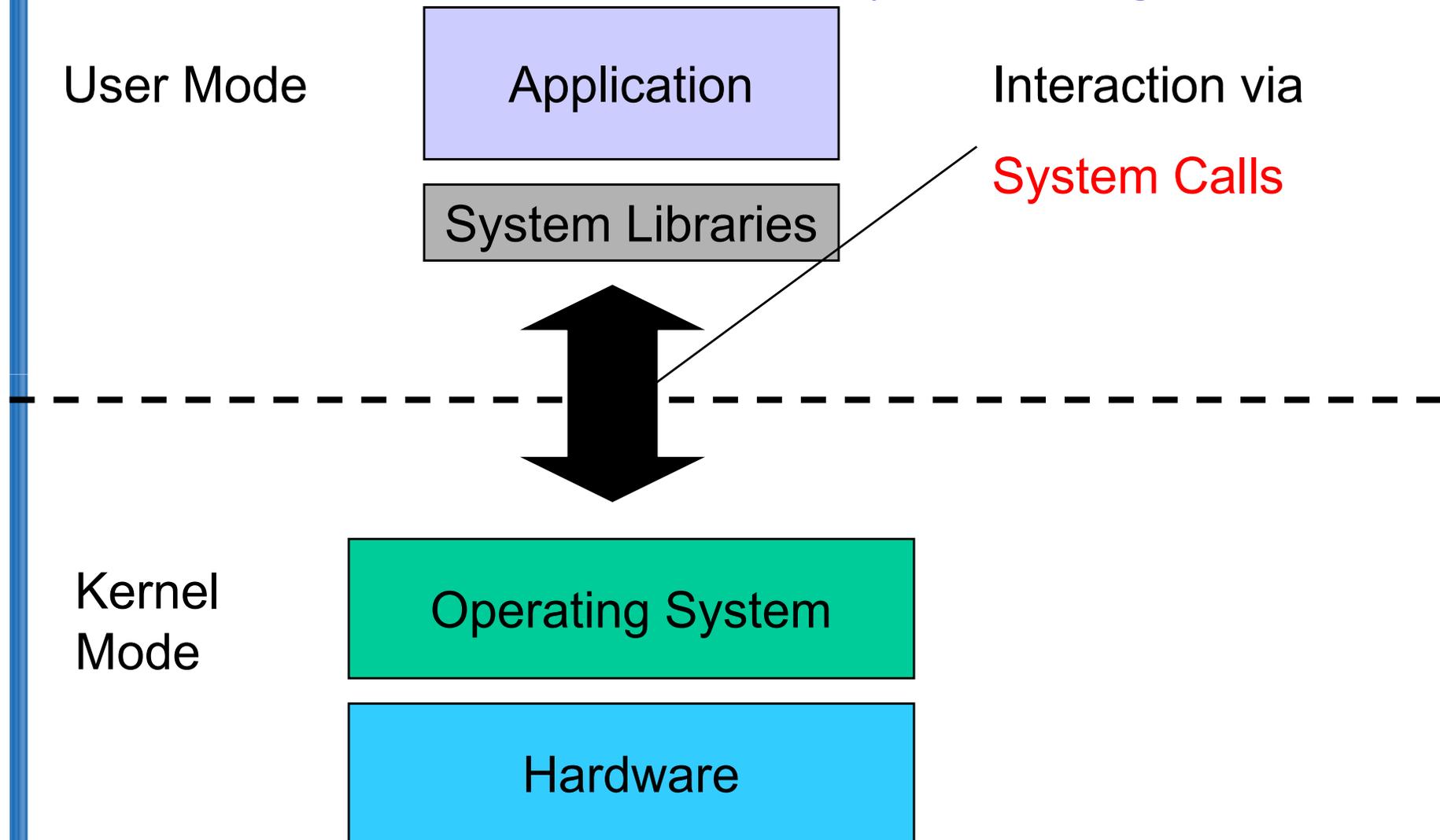
Kernel
Mode

Operating System

Hardware



Structure of a Computer System



A note on System Libraries

- System libraries are just that, libraries of support functions (procedures, subroutines)
 - Only a subset of library functions are actually systems calls
 - strcmp(), memcpy(), are pure library functions
 - open(), close(), read(), write() are system calls
 - System call functions are in the library for convenience



Operating System

- Convenience **Objectives**
 - Make the computer more convenient to use
- Abstraction
 - Hardware-independent programming model
- Efficiency
 - Allows the computer system to be used in an efficient manner
- Ability to evolve
 - Permit effective development, testing, and introduction of new system functions without interfering with existing services
- Protection



Services Provided by the Operating System

- Program development
 - Editors, compilers, debuggers
 - Not so much these days
- Program execution
 - Load a program and its data
- Access to I/O devices
- Controlled access to files
 - Access protection
- System access
 - User authentication



Services Provided by the Operating System

- Error detection and response
 - internal and external hardware errors
 - memory error
 - device failure
 - software errors
 - arithmetic overflow
 - access forbidden memory locations
 - operating system cannot grant request of application



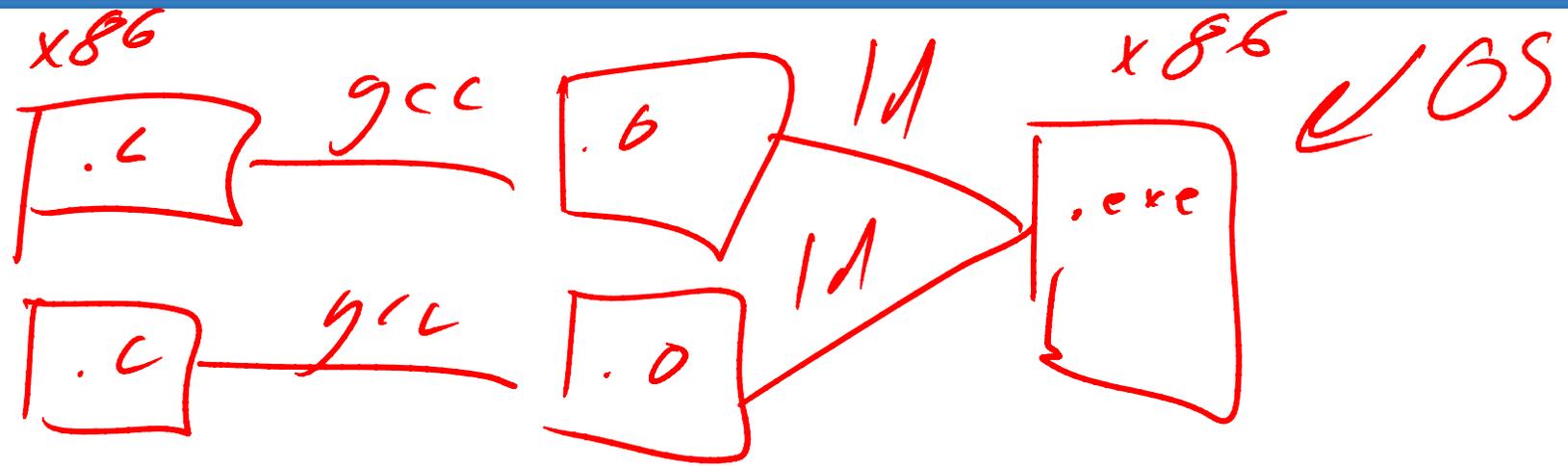
Services Provided by the Operating System

- Accounting
 - collect statistics
 - monitor performance
 - used to anticipate future enhancements
 - used for billing users

Operating System Software

- Fundamentally, OS functions the same way as ordinary computer software
 - It is a program that is executed (just like apps)
 - It has more privileges
- Operating system relinquishes control of the processor to execute other programs
 - Reestablishes control after
 - System calls
 - Interrupts (especially timer interrupts)





CROSS COMPILER

cs161-gcc.



Major OS Concepts

- Processes
- Concurrency and deadlocks
- Memory management
- Files
- Information Security and Protection
- Scheduling and resource management

Processes

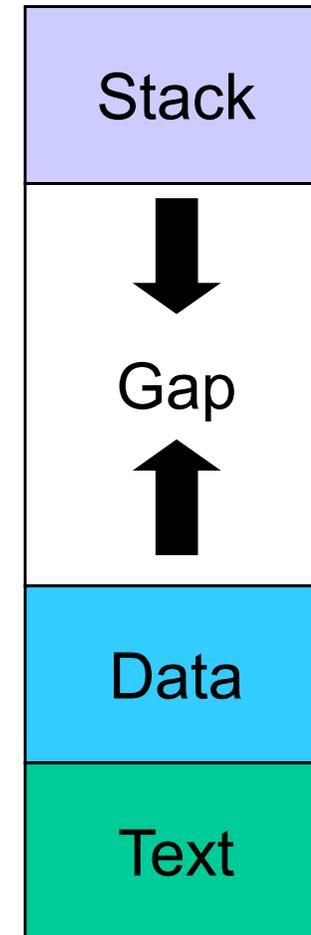
- A program in execution
- An instance of a program running on a computer
- The entity that can be assigned to and executed on a processor
- A unit of resource ownership
- A unit of activity characterized by a single sequential thread of execution, a current state, and an associated set of system resources
 - Nowadays the execution abstraction is separated out:
Thread
 - Single process can contain many threads



Process

- Consist of three segments
 - Text
 - contains the code (instructions)
 - Data
 - Global variables
 - Stack
 - Activation records of procedure
 - Local variables
- Note:
 - data can dynamically grow up
 - The stack can dynamically grow down

Memory

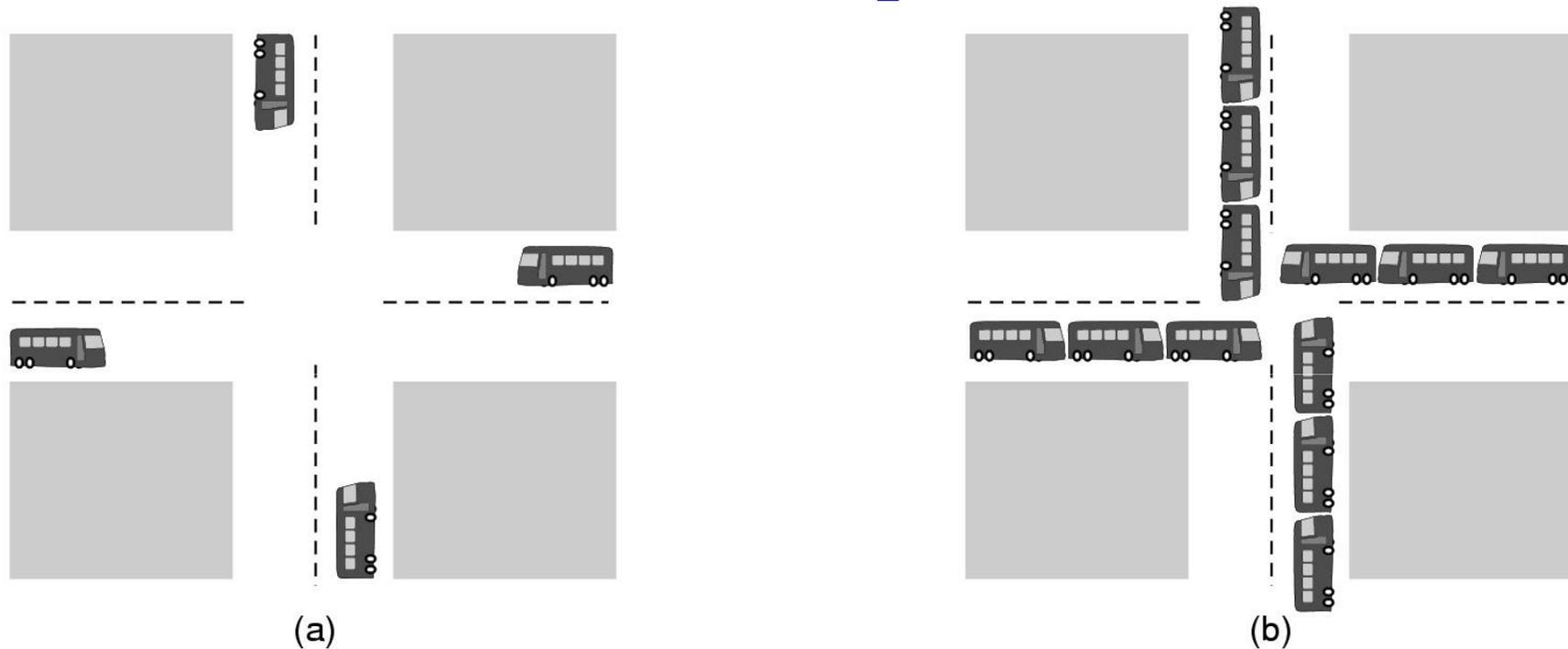


Process

- Consists of three components
 - An executable program
 - text
 - Associated data needed by the program
 - Data and stack
 - Execution context of the program
 - All information the operating system needs to manage the process
 - Registers, program counter, stack pointer, etc...
 - A multithread program has a stack and execution context for each thread



Multiple processes creates concurrency issues



(a) A potential deadlock. **(b)** an actual deadlock.

Memory Management

- The view from thirty thousand feet
 - Process isolation
 - Prevent processes from accessing each others data
 - Automatic allocation and management
 - Don't want users to deal with physical memory directly
 - Protection and access control
 - Still want controlled sharing
 - Long-term storage
 - OS services
 - Virtual memory
 - File system



Virtual Memory

- Allows programmers to address memory from a logical point of view
 - Gives apps the illusion of having RAM to themselves
 - Logical addresses are independent of other processes
 - Provides isolation of processes from each other
- Can overlap execution of one process while swapping in/out others.



Virtual Memory Addressing

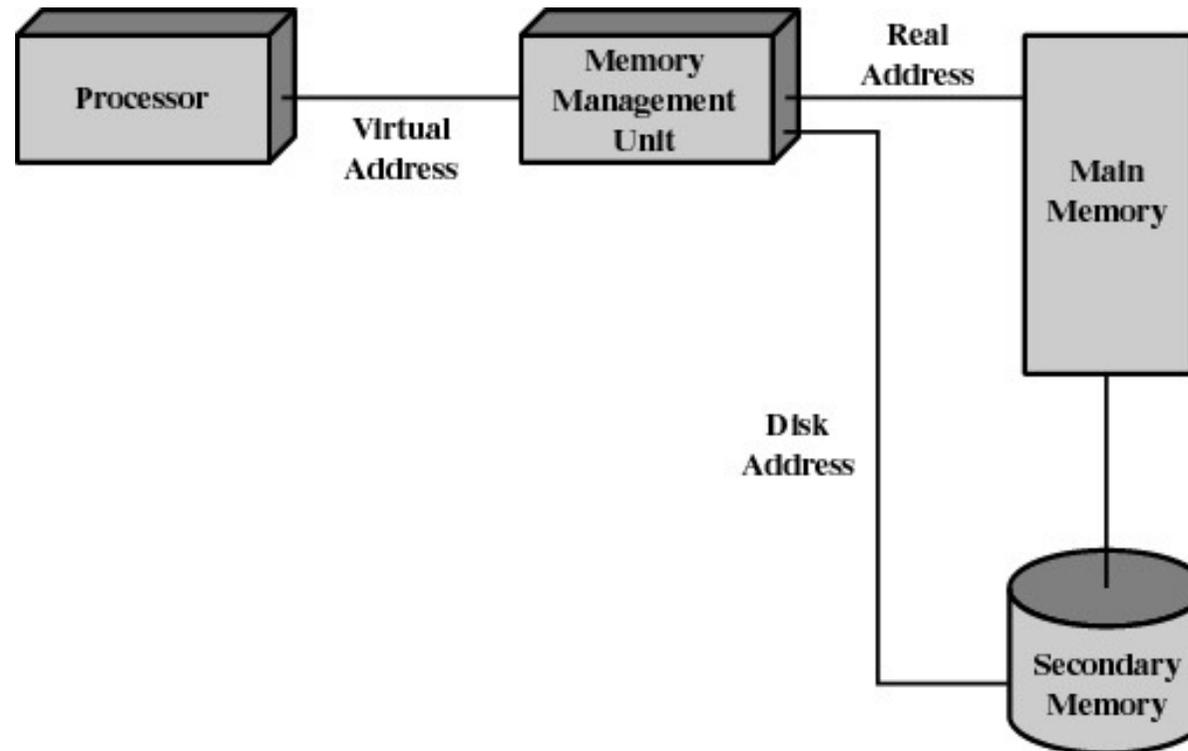


Figure 2.10 Virtual Memory Addressing